



The Gigascale Silicon Research Center

The GSRC Semantics Project

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What is GSRC?

The MARCO/DARPA Gigascale Silicon Research Center

- keep the fabs full
- close the productivity gap
- rebuild the RTL foundation
- enable scalable, heterogeneous, component-based design

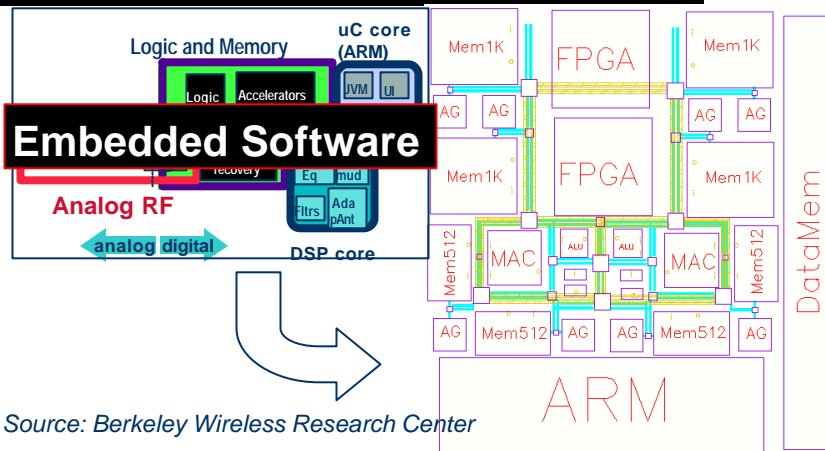
Participants:

- UC Berkeley
- CMU
- Stanford
- Princeton
- UCLA
- UC Santa Barbara
- UC San Diego
- Purdue
- Michigan
- UC Santa Cruz

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What is System Level?



Source: Berkeley Wireless Research Center

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Focus on Capabilities, not Languages

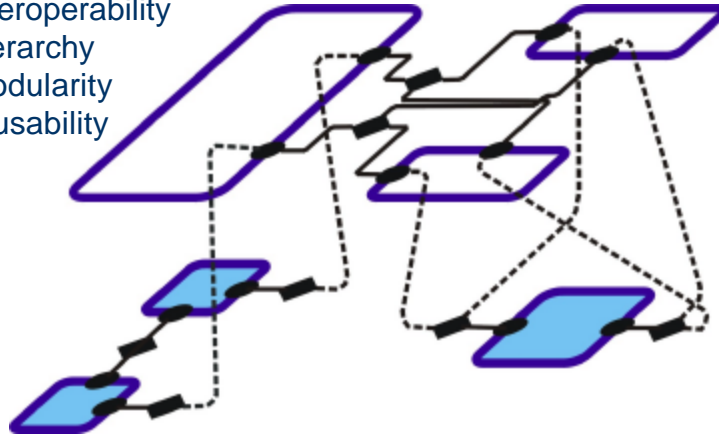
- Modeling
- Simulation
- Visualization
- Synthesis
- Verification
- Modularization

The problem we are here to address is *interoperability* and *design productivity*.
Not standardization.

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Component-Based Design

interoperability
hierarchy
modularity
reusability



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Interoperability Levels

- Code can be written to translate the data from one tool to be used by another.
- Tools can open each other's files and extract useful information (not necessarily *all* useful information).
- Tools can interoperate dynamically, exchanging information at run time.

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Principle: Orthogonalize Concerns in SDLs

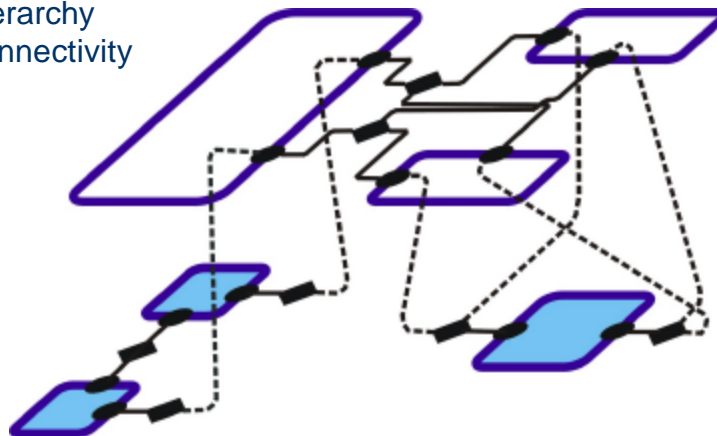
- Abstract Syntax
- Concrete Syntax
- Syntactic Transformations
- Type System
- Component Semantics
- Interaction Semantics

Do this first, since without it, we won't get anywhere

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Abstract Syntax

hierarchy
connectivity



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Not Abstract Syntax

- Semantics of component interactions
- Type system
- File format (a concrete syntax)
- API (another concrete syntax)

An abstract syntax is the logical structure of a design.
What are the pieces, and how are they related?

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Must Be Able to Specify

- Netlists
- Block diagrams
- Hierarchical state machines
- Object models
- Dataflow graphs
- Process networks

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Interfaces and Ports

- A **partially ordered set** *Interfaces*
- A **set** *Ports*
- A **function** *ports: Interfaces* \rightarrow $\wp(\text{Ports})$ s.t.
 - if $i < j$ then $\text{ports}(i) \subseteq \text{ports}(j)$

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Properties

- A **set** *Properties*
- A **function** *properties: Interfaces* \rightarrow $\wp(\text{Properties})$ s.t.
 - if $i < j$ then $\text{properties}(i) \subseteq \text{properties}(j)$

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Inheritance Hierarchy

- *Interfaces*
- a partial ordering relation “<”
- *Ports*
- *ports: Interfaces* \rightarrow $\wp(\text{Ports})$
- *Properties*
- *properties: Interfaces* \rightarrow $\wp(\text{Properties})$

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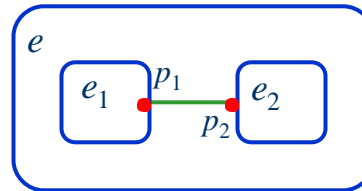
Entities and Containment Hierarchy

- A **set** *Entities*
- A **member** $root \in \text{Entities}$
- A **function** *interface: Entities* \rightarrow *Interfaces*
- A **function** *containedEntities: Entities* \rightarrow $\wp(\text{Entities})$

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Internal Links

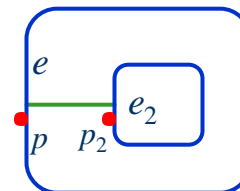
- A function *internalLinks*:
 $Entities \rightarrow \wp(Entities \times Ports \times Entities \times Ports)$
 - $(e_1, p_1, e_2, p_2) \in internalLinks(e) \Rightarrow$
 - $p_1 \in ports(interface(e_1))$
 - $p_2 \in ports(interface(e_2))$
 - $e_1 \in containedEntities(e)$
 - $e_2 \in containedEntities(e)$



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Interface Links

- A function *interfaceLinks*:
 $Entities \rightarrow \wp(Ports \times Entities \times Ports)$
 - $(p, e_2, p_2) \in interfaceLinks(e) \Rightarrow$
 - $p \in ports(interface(e))$
 - $p_2 \in ports(interface(e_2))$
 - $e_2 \in containedEntities(e)$



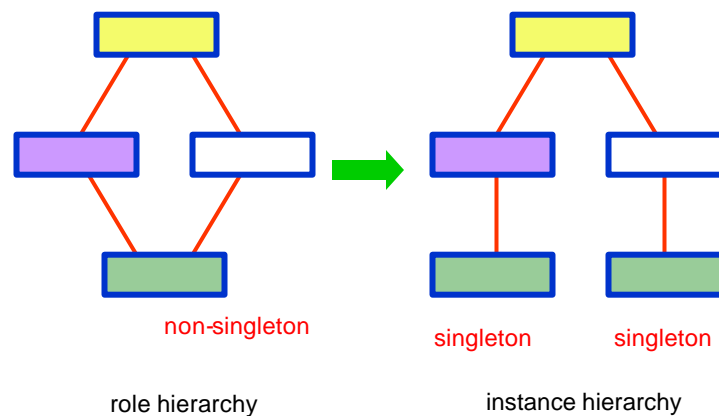
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Instance and Role Hierarchies

- A function *isSingleton*: *Entities* \rightarrow *Boolean*
 - *isSingleton* (*root*) = *true*
- An **instance hierarchy** is a containment hierarchy where
 - $\forall e \in \text{Entities}, \text{isSingleton}(e) = \text{true}$
- A **role hierarchy** is any other containment hierarchy
 - Every role hierarchy can be unrolled to a unique instance hierarchy.

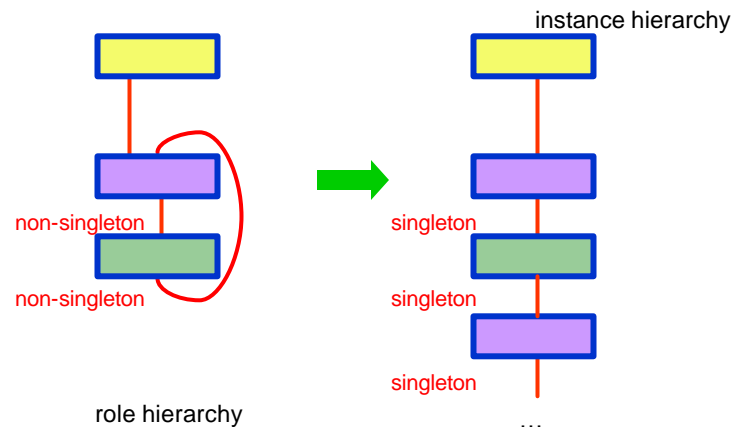
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Unrolling



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Recursive Containment



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The GSRC Abstract Syntax

- Models hierarchical connected components
 - block diagrams, object models, state machines, ...
 - abstraction and refinement
- Supports classes and instances
 - object models
 - inheritance
 - static and instance variables
- Supports multiple simultaneous hierarchies
 - structure and function
 - objects and concurrency

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Concrete Syntaxes

- Persistent file formats
- Close to the abstract syntax
- Make it extensible to capture other aspects
- Enable design data exchange
 - without customization of the tools

Most language discussions focus on concrete syntaxes, which are arguably the least important part of the design

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MoML – An XML Concrete Syntax

```
<?xml version="1.0" standalone="no"?>
<!DOCTYPE model PUBLIC "... " "http://...">
<model name="top" class="path name">
  <entity name="source" class="path name">
    <port name="output"/>
  </entity>
  <entity name="sink" class="path name">
    <port name="input"/>
  </entity>
  <relation name="r1" class="path name"/>
  <link port="source.output" relation="r1"/>
  <link port="sink.input" relation="r1"/>
</model>
```

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MoML DTD

Modeling Markup Language

```

<!ELEMENT link EMPTY>
<!--ATTLIST link port CDATA #REQUIRED
      relation CDATA #REQUIRED
      vertex CDATA #IMPLIED-->

```

Since this document type definition captures only the abstract syntax, it is very small and simple. Other information is embedded using distinct XML DTDs.

```

<!ELEMENT model (attribute | class | configure | doc | director| entity | import | link | relation)*
<!--ATTLIST model name CDATA #REQUIRED
      class CDATA #REQUIRED-->

<!ELEMENT attribute (doc | configure)*
<!--ATTLIST attribute class CDATA #IMPLIED
      name CDATA #REQUIRED
      value CDATA #IMPLIED-->

<!ELEMENT class (attribute | configure | director | doc | entity | link)*
<!--ATTLIST class name CDATA #REQUIRED
      name CDATA #REQUIRED-->

<!ELEMENT configure (PCDATA)
<!--ATTLIST configure source CDATA #IMPLIED-->

<!ELEMENT director (attribute | configure)*
<!--ATTLIST director name CDATA #REQUIRED
      class CDATA #REQUIRED-->

<!ELEMENT doc (PCDATA)

<!ELEMENT entity (attribute | class | configure | doc | director | entity | rendition | relation)*
<!--ATTLIST entity name CDATA #REQUIRED
      class CDATA #REQUIRED-->

<!ELEMENT import EMPTY
<!--ATTLIST import source CDATA #REQUIRED-->

<!ELEMENT link EMPTY
<!--ATTLIST link port CDATA #REQUIRED
      relation CDATA #REQUIRED
      vertex CDATA #IMPLIED-->

<!ELEMENT location EMPTY
<!--ATTLIST location x CDATA #REQUIRED
      y CDATA #IMPLIED
      z CDATA #IMPLIED-->

<!ELEMENT port (doc | configure)*
<!--ATTLIST port name CDATA #REQUIRED
      class CDATA #REQUIRED
      direction (input | output | both) "both"-->

<!ELEMENT relation (vertex)*
<!--ATTLIST relation name CDATA #REQUIRED
      class CDATA #REQUIRED-->

<!ELEMENT rendition (configure | location)*
<!--ATTLIST rendition class CDATA #REQUIRED-->

<!ELEMENT vertex (location)*
<!--ATTLIST vertex name CDATA #REQUIRED
      pathTo CDATA #REQUIRED-->

```

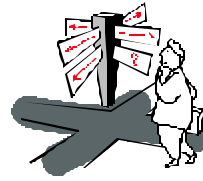
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Syntactic Transformations

- A set of operations on models
 - creation of ports, relations, links, and entities
 - mutation
- Applications
 - visual editors
 - higher-order functions
 - instantiation
 - unrolling recursion

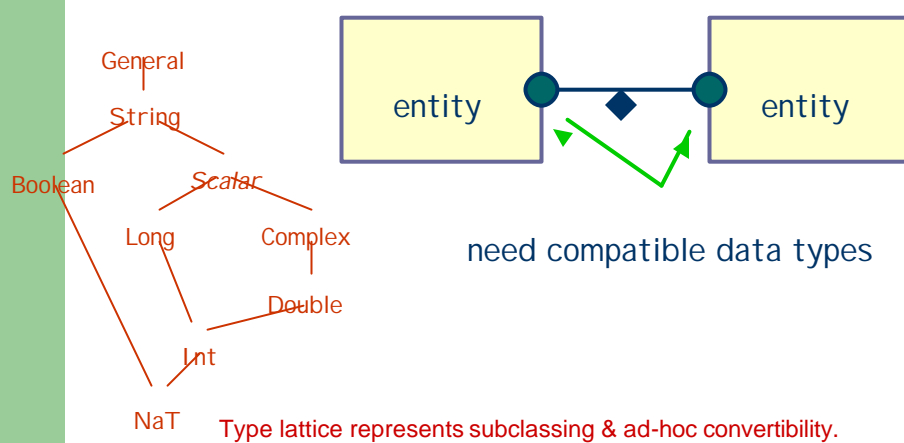
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Where We Are...

- Abstract Syntax ✓
 - Concrete Syntax ✓
 - Syntactic Transformations ✓
 - Type System
 - Component Semantics
 - Interaction Semantics
- } logical structure
- } meaning

Type Systems





Desirable Properties in a Type System

- Strong typing
- Polymorphism
- Propagation of type constraints
- Composite types (arrays, records)
- User-defined types
- Reflection
- Higher-order types
- Type inference
- Dependent types

We can have compatible type systems without compatible languages (witness CORBA)

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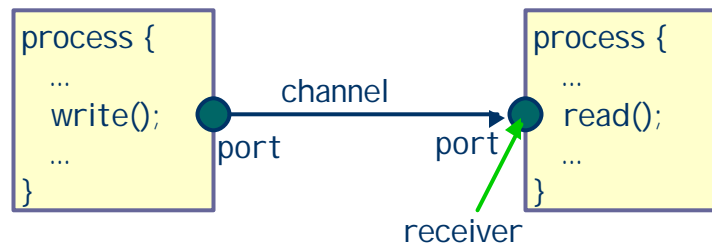
Component Semantics

Entities are:

- States?
- Processes?
- Threads?
- Differential equations?
- Constraints?
- Objects?

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One Class of Semantic Models: Producer / Consumer



- Are actors active? passive? reactive?
- Are communications timed? synchronized? buffered?

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Particular Consumer/Producer Frameworks (*Domains*)

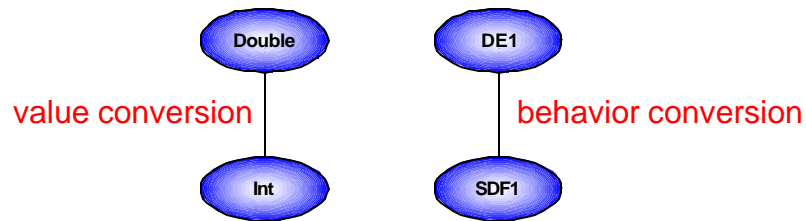
- CSP – concurrent threads with rendezvous
- CT – continuous-time modeling
- DE – discrete-event systems
- DT – discrete time (cycle driven)
- PN – process networks
- SDF – synchronous dataflow
- SR – synchronous/reactive

Each of these defines a component ontology and an interaction semantics between components. There are many more possibilities!

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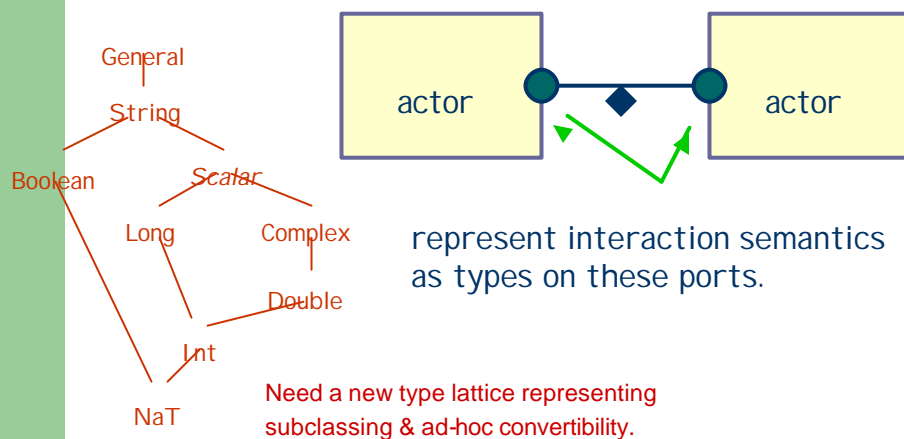
Interfaces

- Represent not just data types, but interaction types as well.



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GSRG Current Approach – System-Level Types



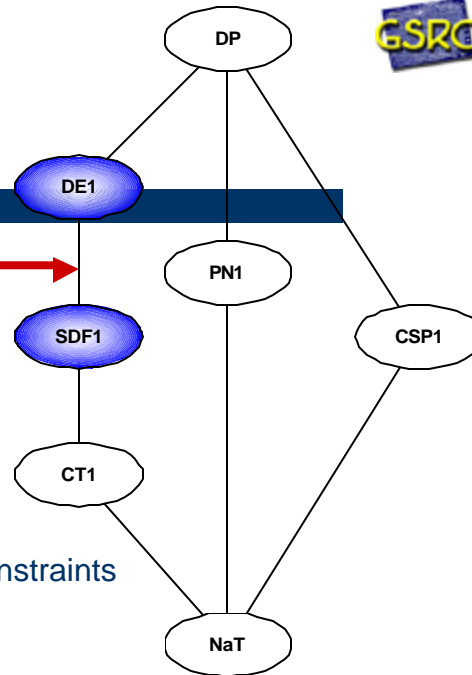
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Type Lattice

Simulation relation →

Achievable properties:

- Strong typing
- Polymorphism
- Propagation of type constraints
- User-defined types
- Reflection



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System-Level Types

- Declare dynamic properties of component interfaces
- Declare timing properties of component interfaces

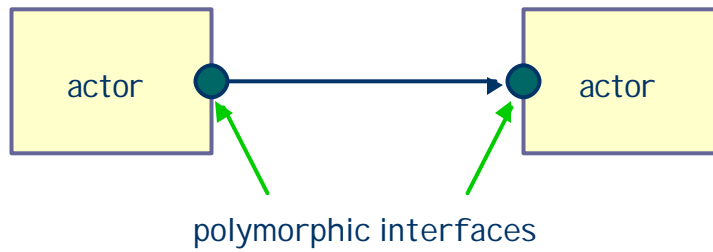
Benefits:

- Ensure component compatibility
- Clarify interfaces
- Provide the vocabulary for design patterns
- Detect errors sooner
- Promote modularity
- Promote polymorphic component design

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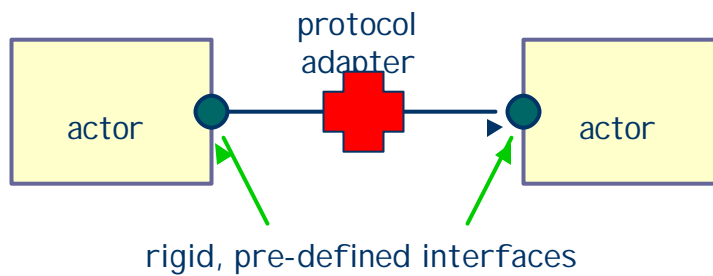


Our Hope – Polymorphic Interfaces

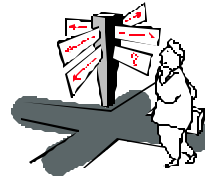


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Approach Used by Others – Interface Synthesis



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Where We Are...

- Abstract Syntax ✓
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- Component Semantics ✓
- Interaction Semantics ✓

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Benefits of Orthogonalization

- Modularity in language design
 - e.g. can build on existing abstract syntax
- Different levels of tool interoperability
 - e.g. visualization tool needs only the abstract syntax
- Terminology independent of concrete syntax
 - e.g. design patterns
- Focus on frameworks instead of languages
 - dealing with heterogeneity
- Issue-oriented not ASCII-oriented

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Ptolemy Project – Sanity Check



Ptolemy II –

- A reference implementation
- Testbed for abstract syntax
- Block diagram MoML editor
- Mutable models
- Extensible type system
- Testbed for system-level types

<http://ptolemy.eecs.berkeley.edu>