

Hybrid Systems: From Models to Code

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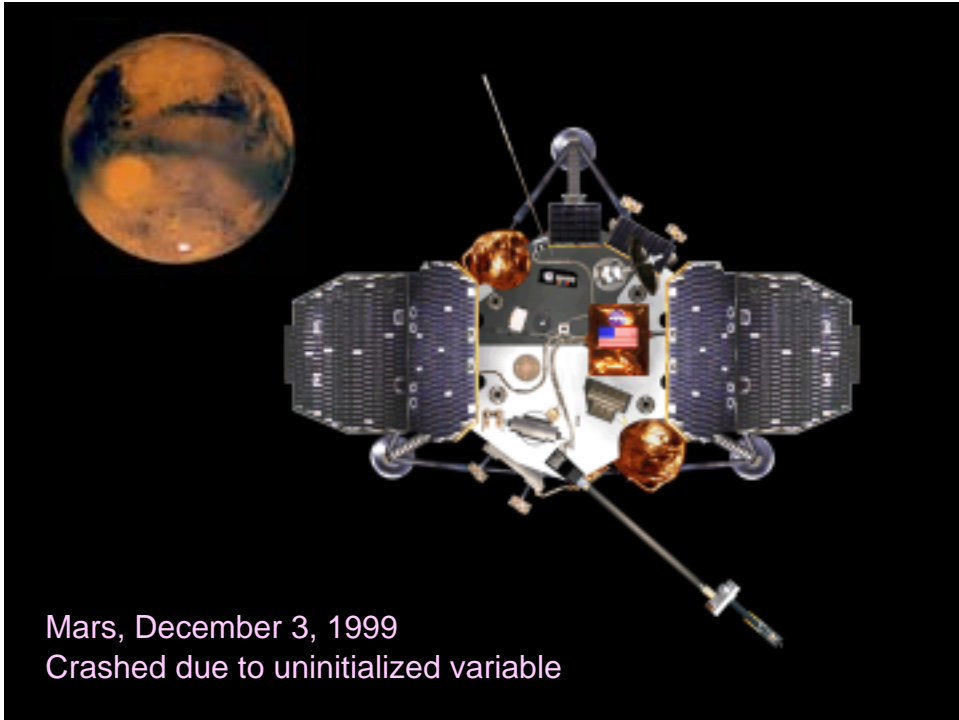


UC Berkeley: Chess
Vanderbilt University: ISIS
University of Memphis: MSI

Foundations of Hybrid and Embedded Software Systems

French Guyana, June 4, 1996
\$800 million embedded software failure





Mars, December 3, 1999
Crashed due to uninitialized variable



\$4 billion development effort
40-50% system integration & validation cost

Sources of Complexity

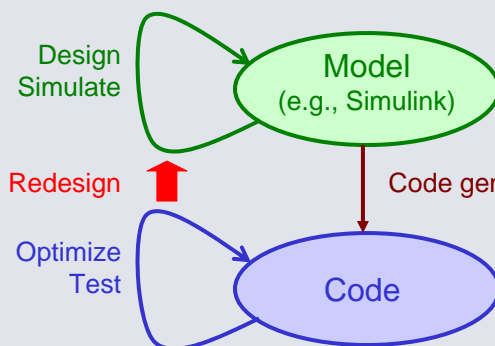


- concurrency
- real time
- heterogeneity

A hybrid system consists of multiple continuous (physical) and discrete (computational) components that interact with each other in real time.

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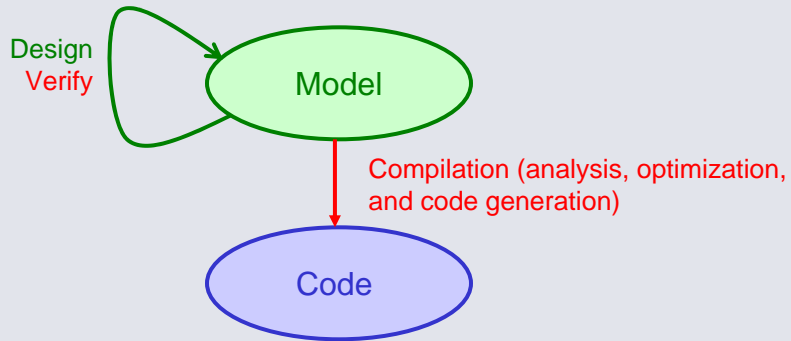
Embedded Software Design: Current State



No formal connection between requirements, model, and resources:
expensive development cycle iterates all stages

No exact correspondence between model and code:
-difficult to upgrade code
-difficult to reuse code

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The FRESCO Project (Formal Real-Time Software Components)

Hybrid System Model

MASACCIO:
correctness by formal verification
against requirements

Time-Safe Code

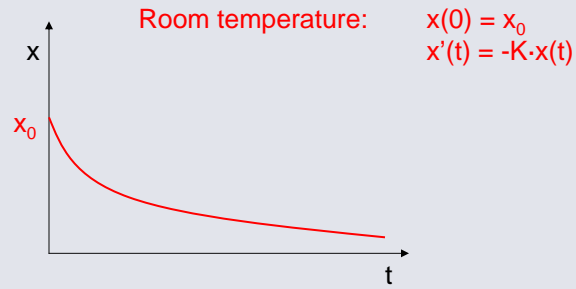
GIOTTO:
correctness by schedulability
analysis against resources

Continuous (Euclidean) Systems



State space: \mathbb{R}^n

Dynamics: initial condition + differential equations



Analytic complexity.

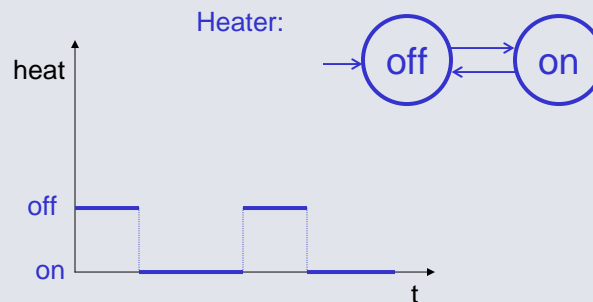
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Discrete (Boolean) Systems



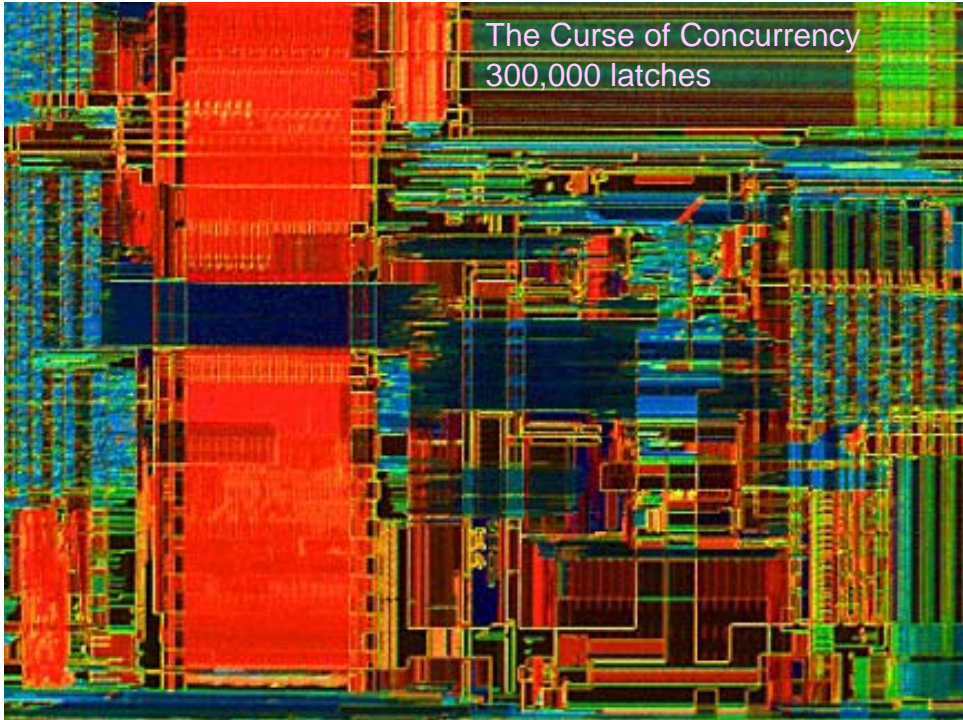
State space: \mathbb{B}^m

Dynamics: initial condition + transition relation



Combinatorial complexity.

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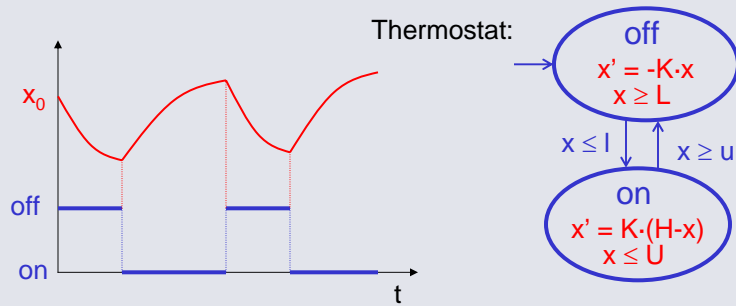


Hybrid Systems



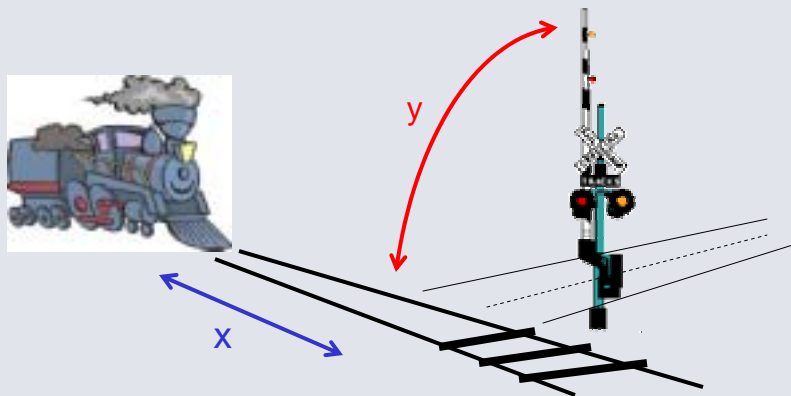
State space: $B^m \times R^n$

Dynamics: initial condition + transition relation
+ differential equations



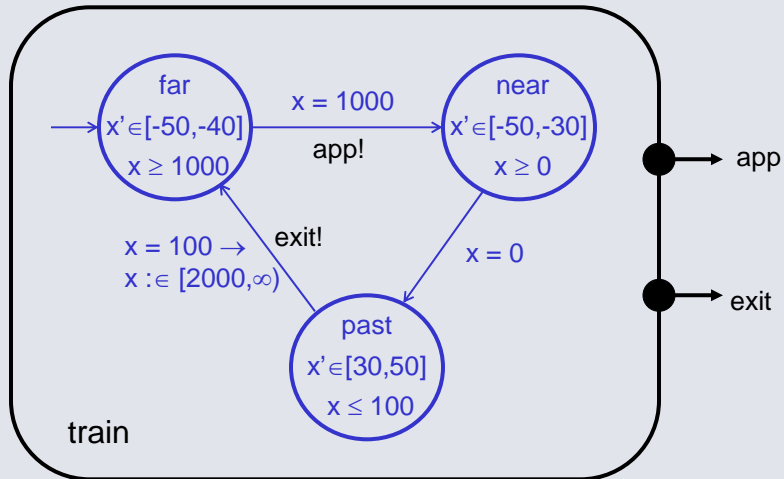
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Hybrid Automata

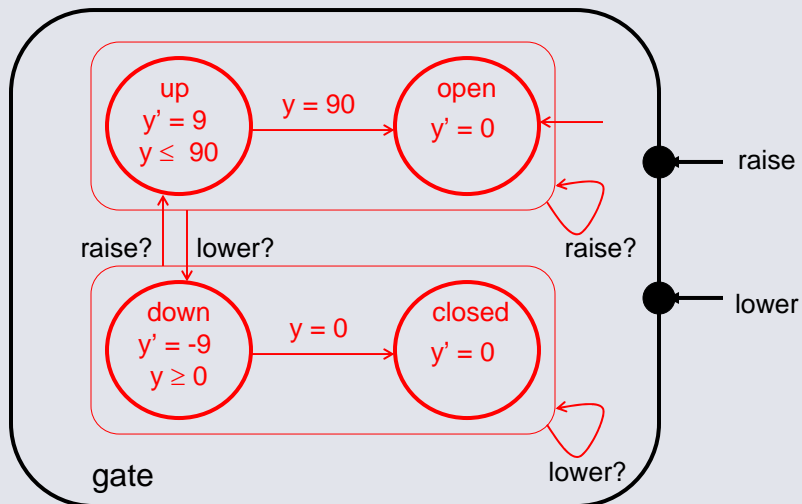


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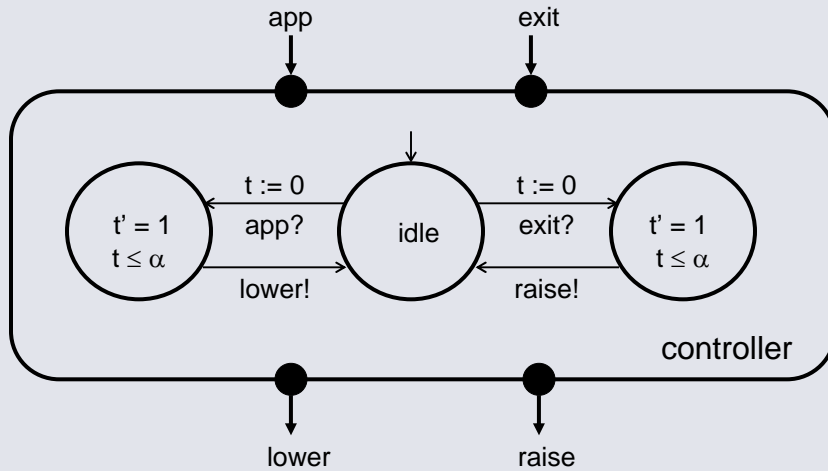
Hybrid Automata



Hybrid Automata



Hybrid Automata



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Requirements



Safety: $\forall \square (x \leq 10 \Rightarrow loc[gate] = closed)$

Liveness: $\forall \square \forall \diamond (loc[gate] = open)$

Real time: $\forall \square z := 0. (z' = 1 \Rightarrow \forall \diamond (loc[gate] = open \wedge z \leq 60))$

Verification and failure analysis by model checking (e.g., HyTech).

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Two Problems with Hybrid Automata



1. Scalability

Possible solutions:

- hierarchy (MASACCIO)
- assume-guarantee decomposition (interfaces)

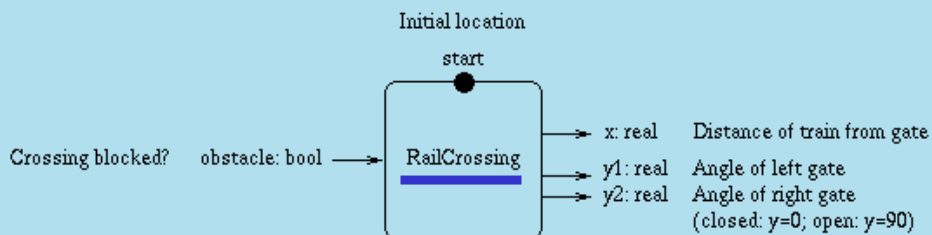
2. Robustness

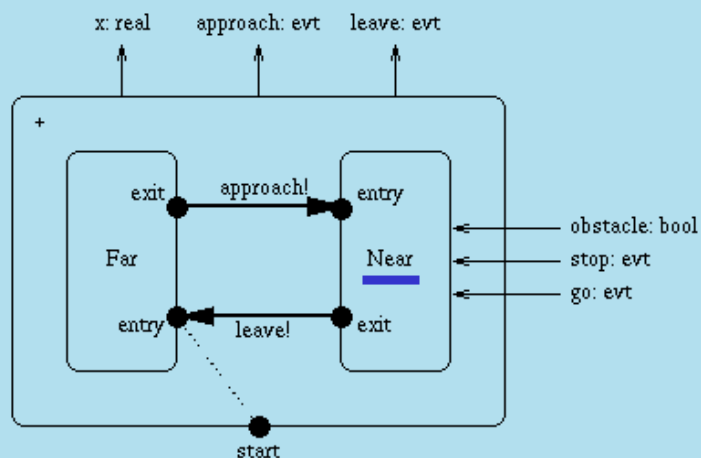
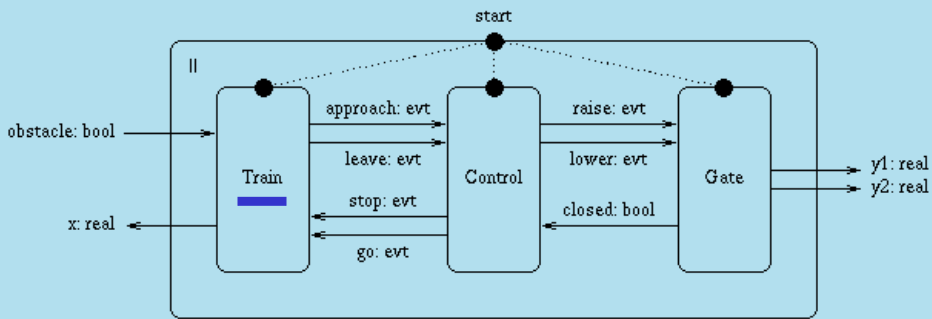
Possible solutions:

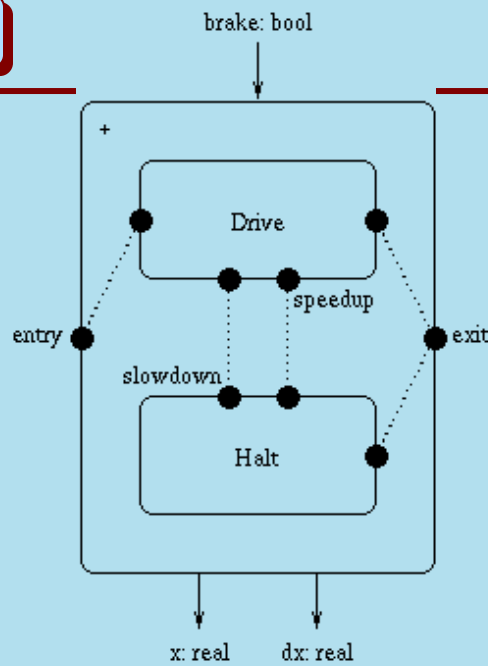
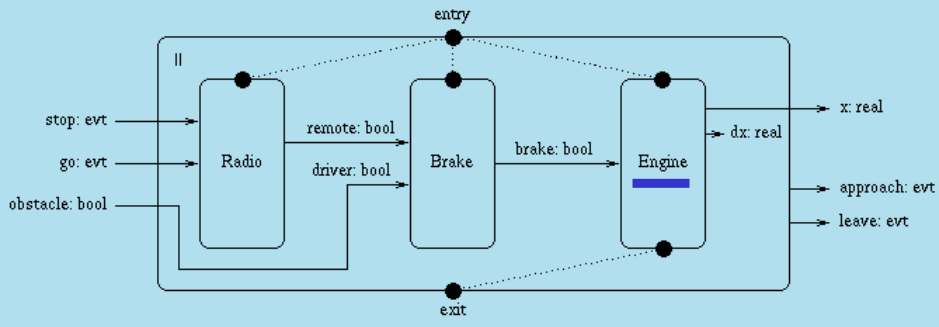
- ϵ -variability
- discounted future

MASACCIO

Hierarchical Hybrid Automata







Two Problems with Hybrid Automata



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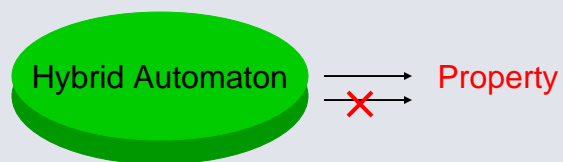
2. Robustness

Possible solutions:

- ϵ -variability
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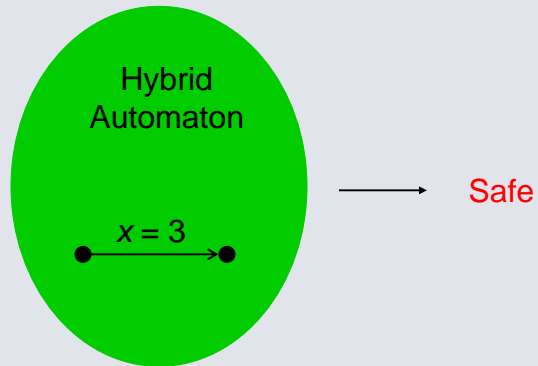
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The Robustness Problem



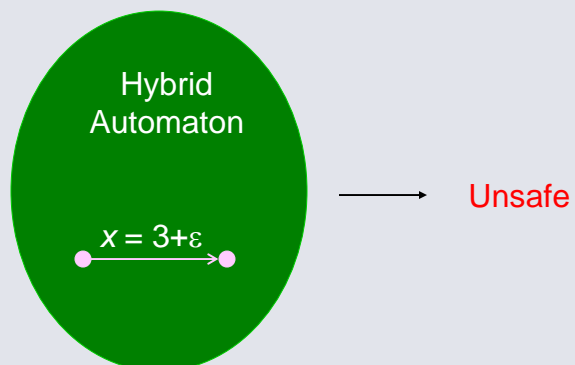
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The Robustness Problem



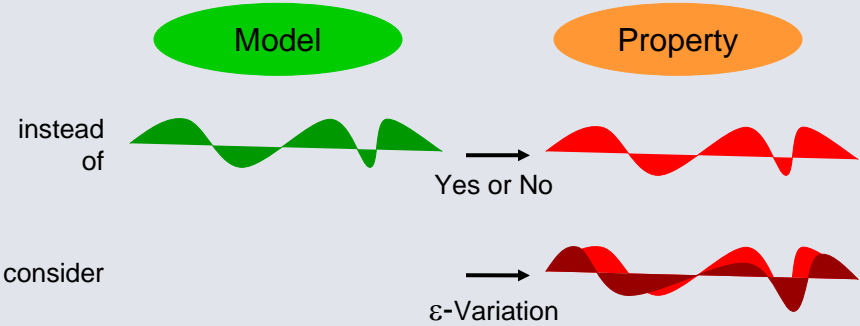
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The Robustness Problem



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A Possible Solution of the Robustness Problem: Metrics on Traces



A More Radical Solution of the Robustness Problem: Discounting the Future



$\text{value}(\text{Model}, \text{Property}): \text{States} \rightarrow \{\text{Yes}, \text{No}\}$



$\text{value}(\text{Model}, \text{Property}): \text{States} \rightarrow \mathbb{R}$

A More Radical Solution of the Robustness Problem: Discounting the Future



$\text{value}(\text{Model}, \text{Property}): \text{States} \rightarrow \{\text{Yes}, \text{No}\}$

$$\text{value}(m, \diamond T) = \mu X. (T \vee \text{pre}(X))$$



$\text{discountedValue}(\text{Model}, \text{Property}): \text{States} \rightarrow \mathbb{R}$

$$\text{discountedValue}(m, \diamond T) = \mu X. \max(T, \lambda \cdot \text{pre}(X))$$

discount factor $0 < \lambda < 1$

A More Radical Solution of the Robustness Problem: Discounting the Future



Robustness Theorem:

If $\text{discountedBisimilarity}(m_1, m_2) > 1 - \varepsilon$,
then $|\text{discountedValue}(m_1, p) - \text{discountedValue}(m_2, p)| < f(\varepsilon)$.

Further Advantages of Discounting:

- approximability** because of geometric convergence (avoids non-termination of verification algorithms)
- applies also to **probabilistic** systems and to **games** (enables reasoning under uncertainty and control)

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Hybrid System Model

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The History of Computer Science: Lifting the Level of Abstraction



High-level languages:
Programming to the application

Requirements
focused code

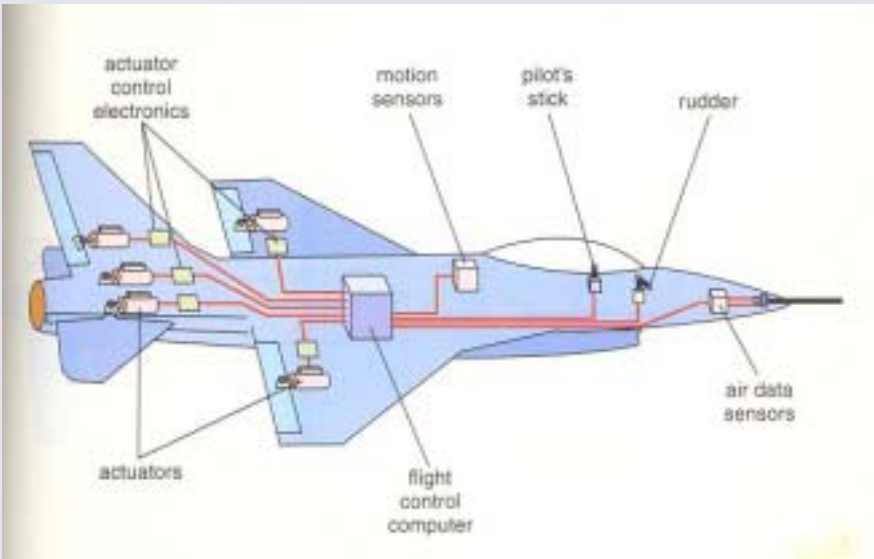


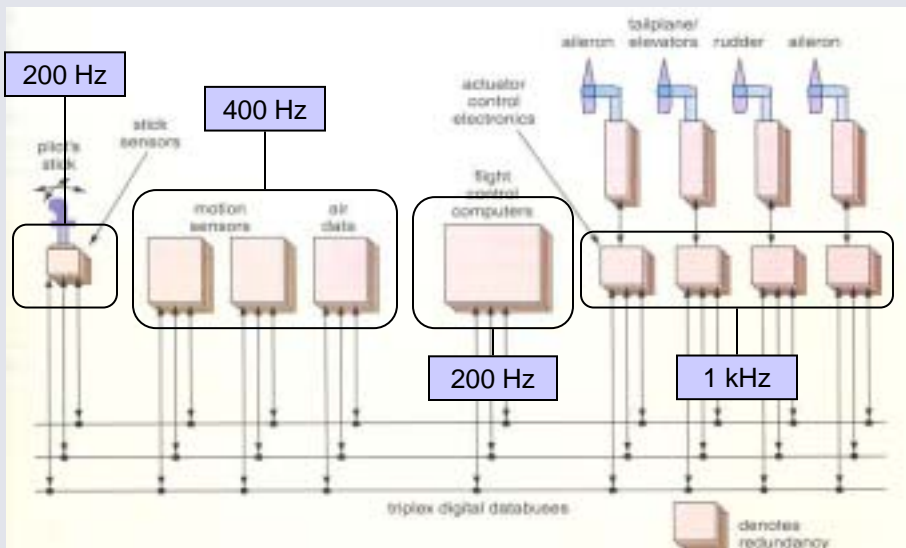
Compilation

The "assembly age":
Programming to the platform

Resource
focused code

- Traditional high-level languages abstract time.
- This abstraction is unsuitable for real-time applications, which are still programmed in terms of **platform time** ("priority tweaking").
- GIOTTO: Real-time programming in terms of **application time**.



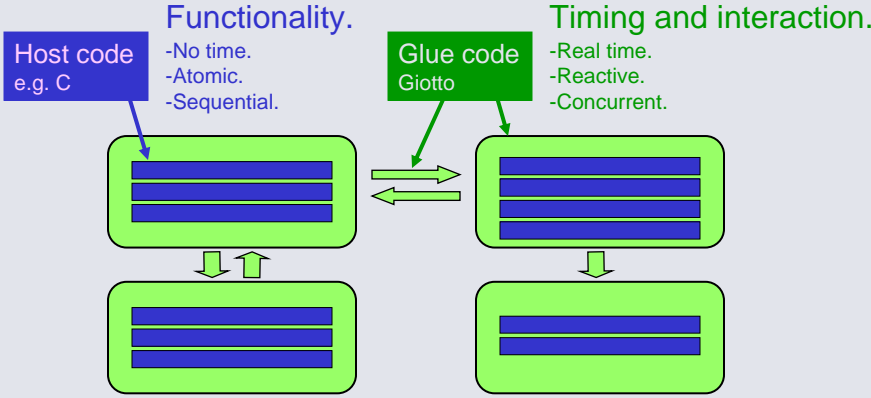
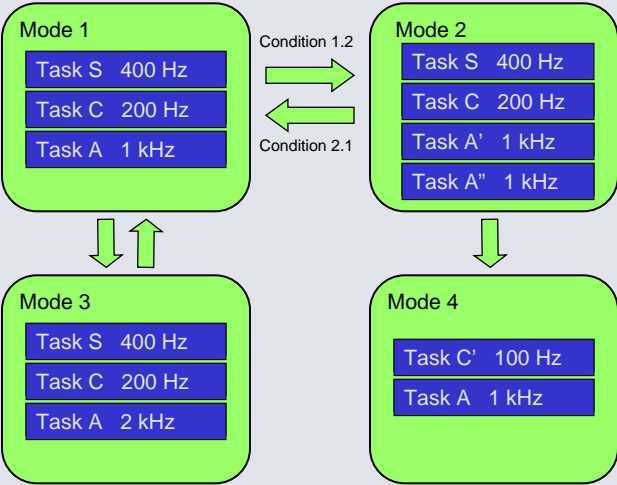


1. Concurrent Periodic Tasks:

- sensing
- control law computation
- actuating

2. Multiple Modes of Operation:

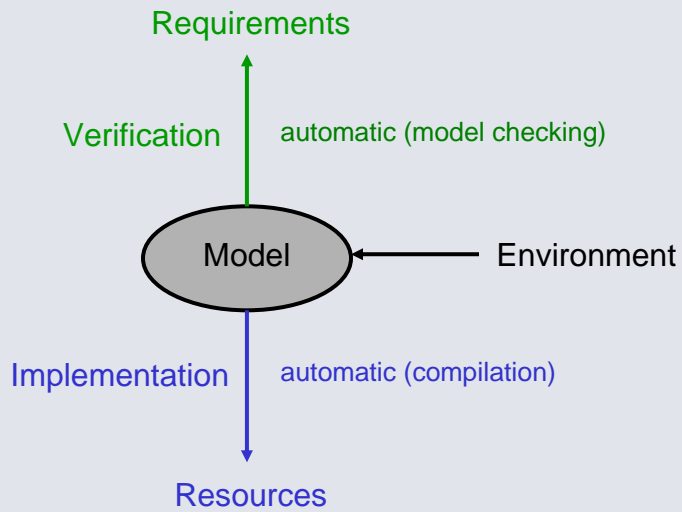
- navigational modes (autopilot, manual, etc.)
- maneuver modes (taxi, takeoff, cruise, etc.)
- degraded modes (sensor, actuator, CPU failures)



This kind of software is reasonably well understood.

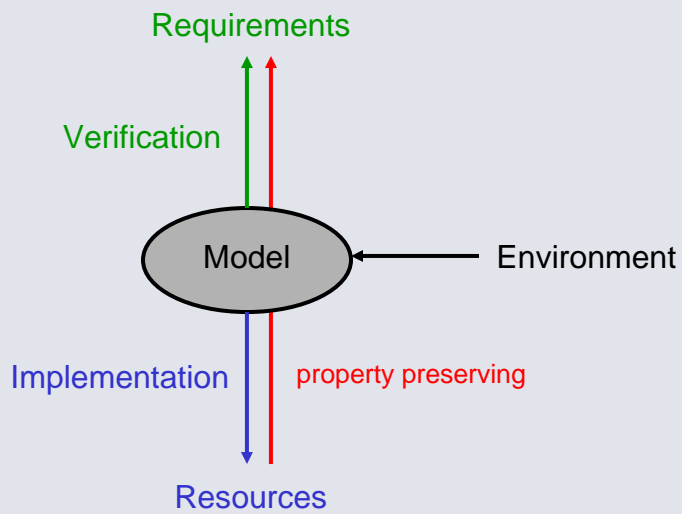
The software complexity lies in the glue code.

Two Opposing Forces



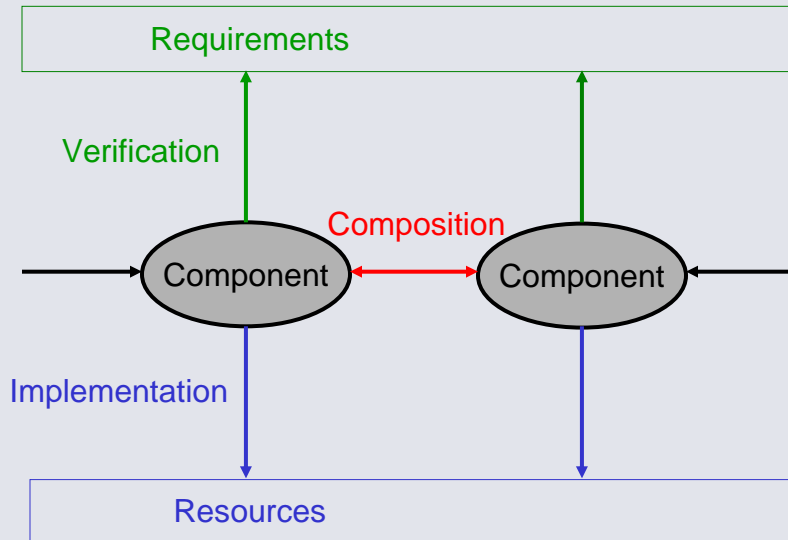
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Two Opposing Forces



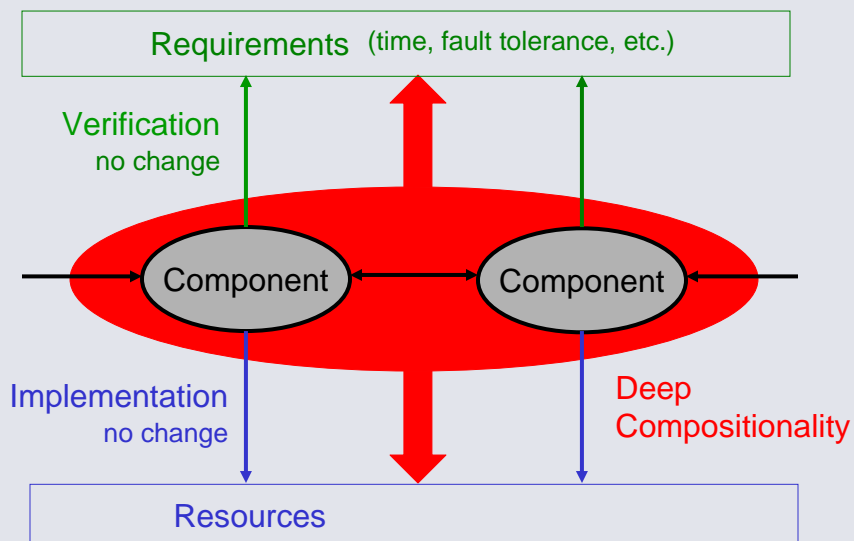
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Two Opposing Forces



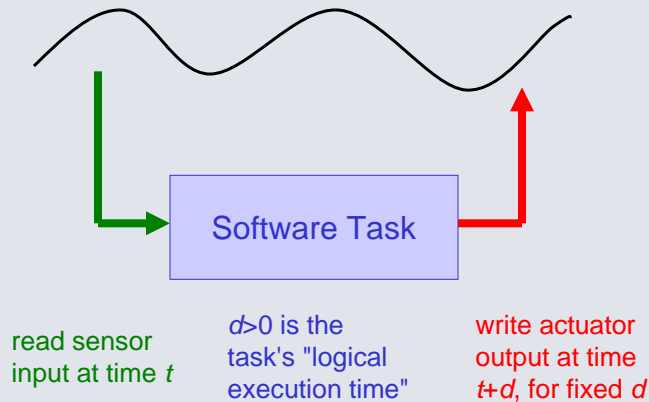
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Two Opposing Forces



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Achieving Verifiability and Compositionality in GIOTTO: The FLET (Fixed Logical Execution Time) Assumption



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Embedded Programming in GIOTTO

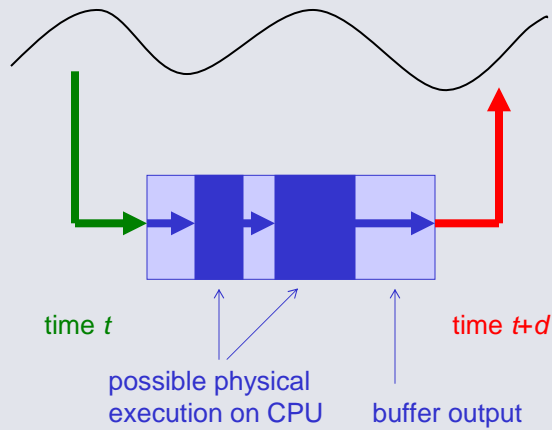


The **programmer** specifies sample rate d and jitter j to solve the control problem at hand.

The **compiler** ensures that d and j are met on a given platform (hardware resources and performance); otherwise it rejects the program.

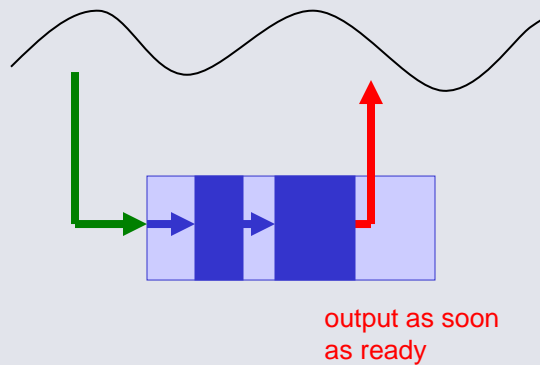
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Implementing the FLET Assumption



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Contrast the FLET with Standard Practice



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Advantages of the FLET and GIOTTO



- predictable** timing and value behavior (no internal race conditions, minimal jitter)
- portable, composable** code (as long as the platform offers sufficient performance)

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Research Agenda



From Hybrid Models

- robust hybrid models (tube topologies, discounting)
- model checking for hierarchical and stochastic hybrid models
- multi-aspect assume-guarantee decomposition of hybrid models (interface theories for time, resources, fault tolerance)

To Embedded Code

- distributed schedulability analysis and code generation
- on-line code modification and fault tolerance

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Credits



Scalable and Robust Hybrid Systems: Luca de Alfaro, Arkadeb Ghosal, Marius Minea, Vinayak Prabhu, Marcin Jurdzinski, Rupak Majumdar

GIOTTO: Ben Horowitz, Christoph Kirsch, Rupak Majumdar, Slobodan Matic, Marco Sanvido

Collaborators of the FRESCO Project



- Alex Aiken on time-safety analysis of embedded code
- Karl Hedrick on Giotto implementation of electronic throttle control
- Edward Lee on Giotto modeling and code generation in Ptolemy
- Edward Lee on rich interface theories as type theories for component interaction
- George Necula on model checking device drivers
- George Necula on scheduler-carrying embedded code
- Alberto Sangiovanni-Vincentelli on synthesis of protocol converters from interfaces
- Alberto Sangiovanni-Vincentelli and Shankar Sastry on platform-based design of a helicopter flight control system using Giotto
- Shankar Sastry on hybrid automata