



Formal Analysis of Timing Effects on Closed-loop Properties of Cyber Physical Systems

Arne Hamann, Corporate Research, Robert Bosch GmbH

Joint work with: Matthias Wöhrle (Bosch), Goran Frehse (Université Joseph Fourier Grenoble),

Sophie Quinton (INRIA Grenoble)

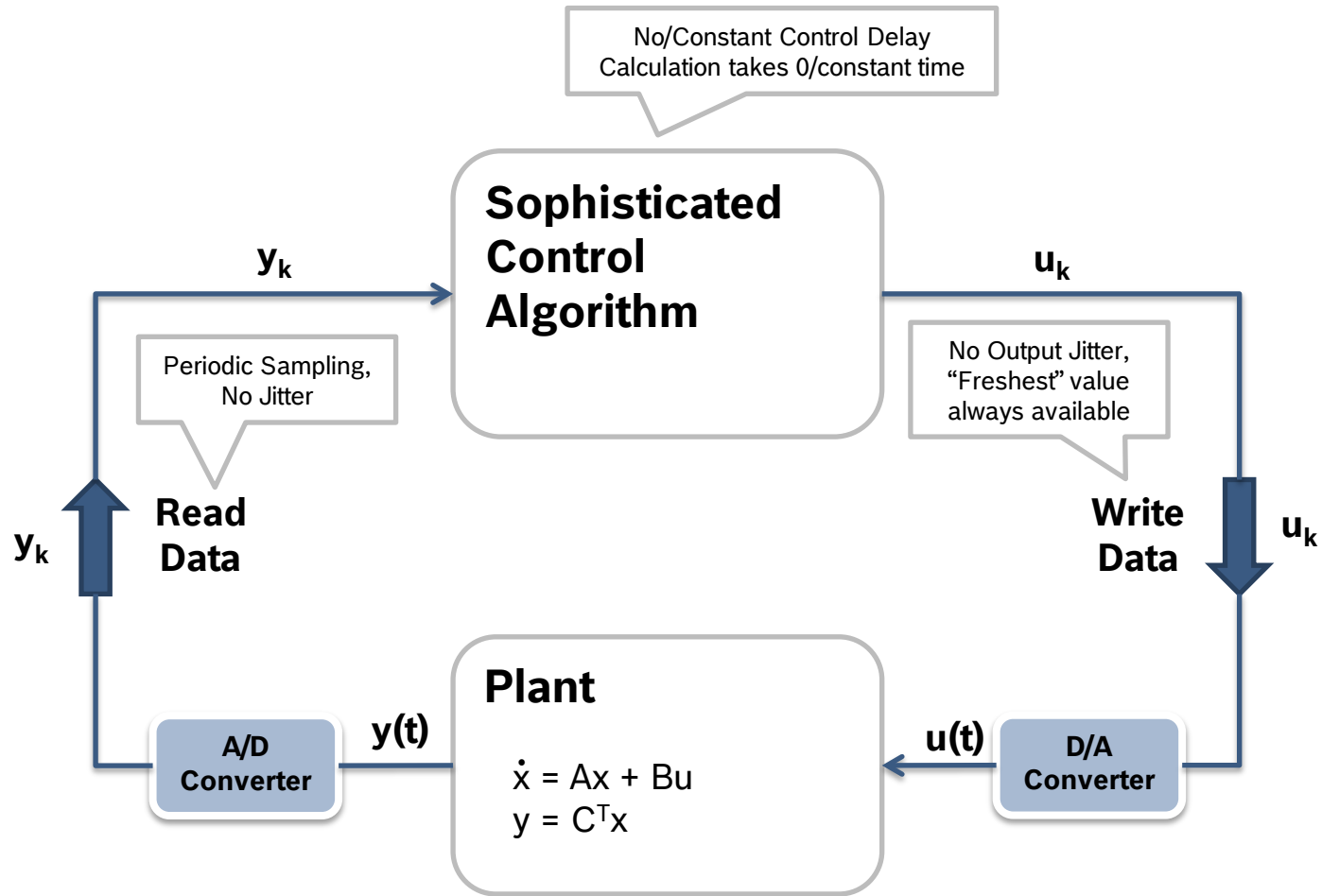


BOSCH

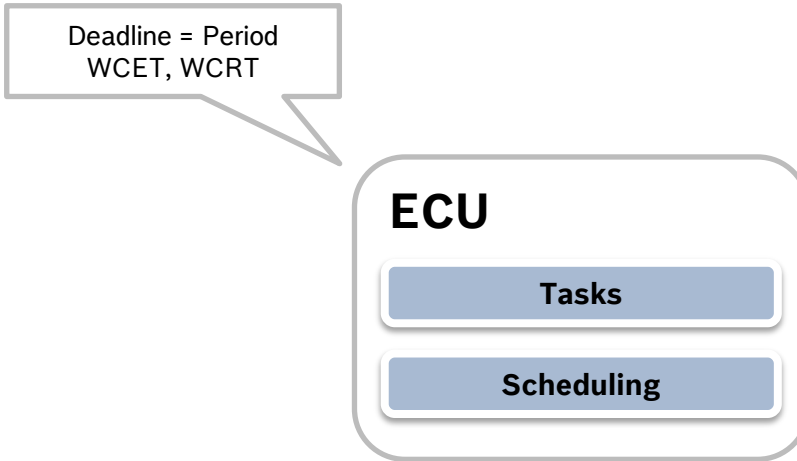
Outline

- Problem statement & goals
- Interaction model for co-engineering between control and real-time engineering
- Electro Mechanic Braking System (EMB)
- Formal analysis of EMB system using hybrid automata and reachability analysis
- Conclusion

System as seen by the control engineer



System as seen by the real-time engineer



$$\sum_{i=1}^n \frac{C_i}{T_i} \leq n \cdot (\sqrt[n]{2} - 1)$$

$\ln 2 \approx 69,3\%$

$$R_i = C_i + \sum_{j \in \text{hp}(i)} C_j \quad \left\lceil \frac{R_i}{T_j} \right\rceil \leq D_i = T_i$$

Problem Statement - Shortcomings

Control engineering

- Theory:
 - Equidistant sampling
 - Zero input-output latencies
- Reality:
 - Varying execution and response times due to preemption, blocking, data-dependencies, ...
 - Sampling interval jitter
 - Non negligible response times



Real-time system engineering

- Theory:
 - Timing models and requirements that are motivated by the runtime system rather than functionality (e.g. deadline = period)
- Reality:
 - Timing requirements do not exist per se and must be derived from functional requirements



Result:

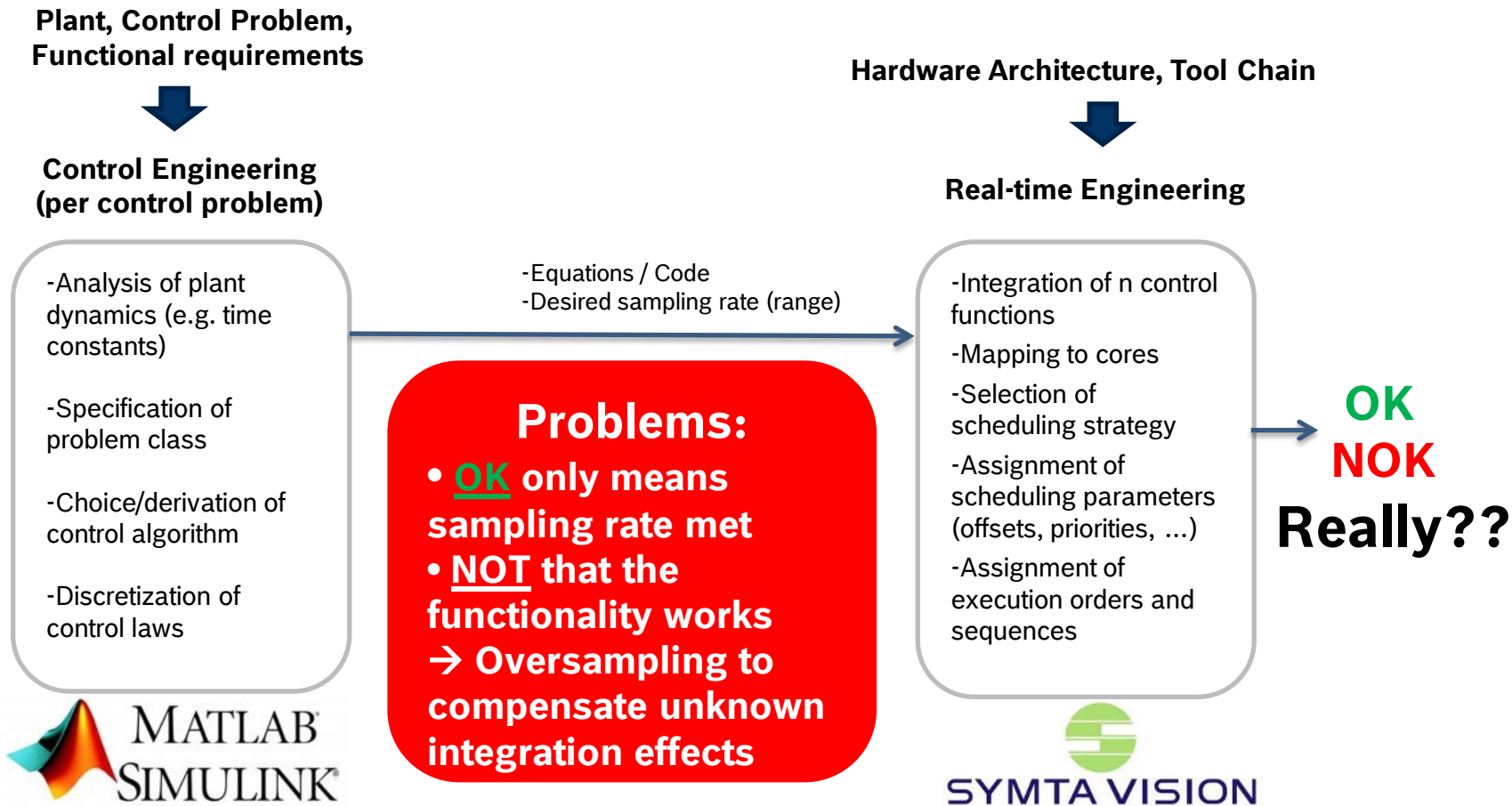
- Functional integration effects due to timing are unpredictable
- Severe migration problems in case of platform modifications



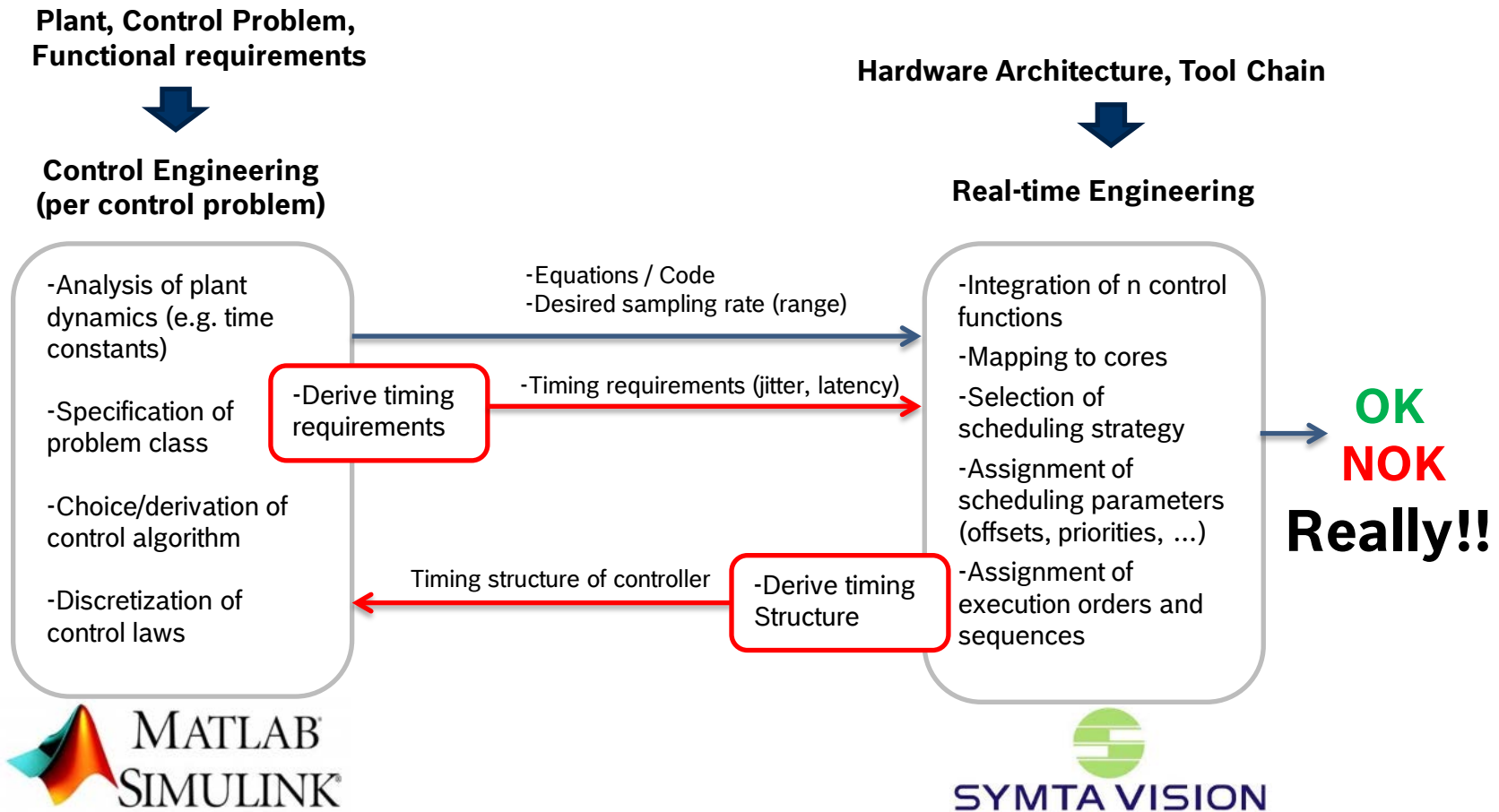
Goals

- **Co-engineering** between real-time and control engineering
- Assessment of functional behavior under the influence of resource sharing **during design time on PC**
- Systematically **derive timing requirements** that are necessary to fulfill functional requirements
- Use these timing requirements for **system synthesis** using adequate platform mechanisms

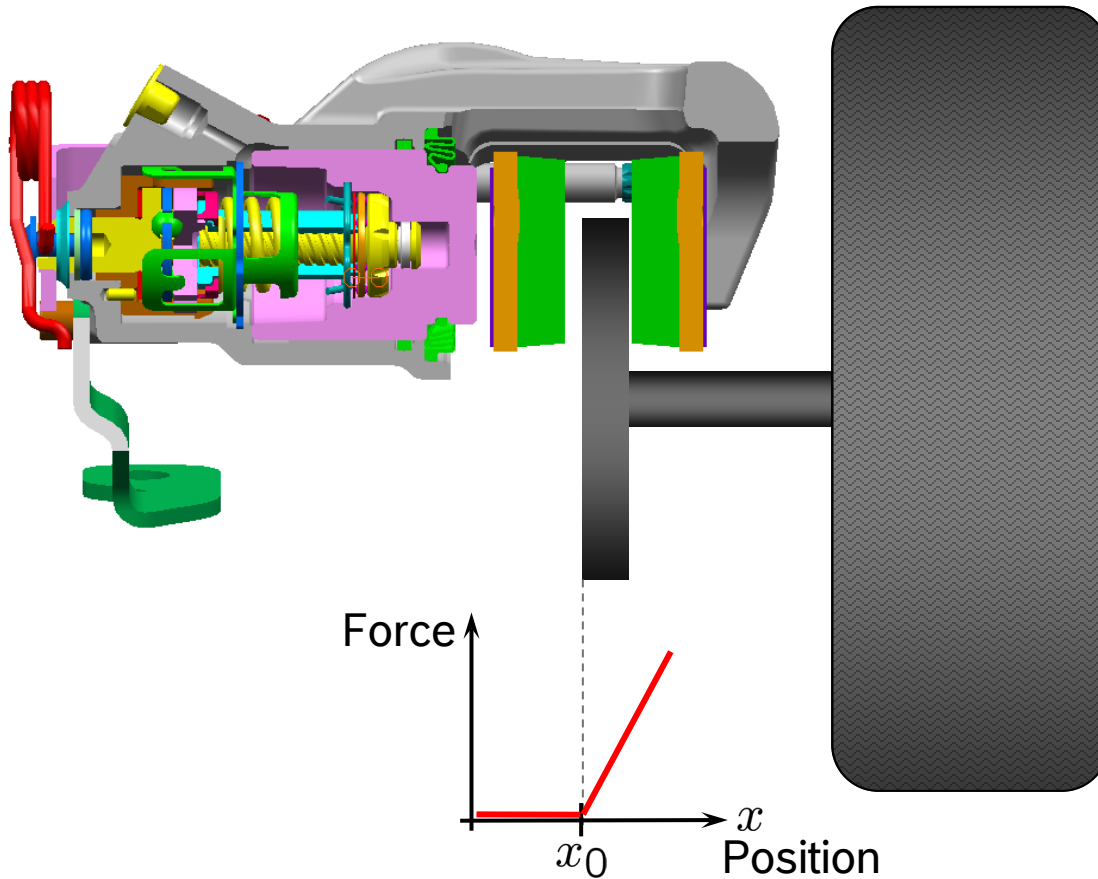
Current Interaction Model



Co-engineering Interaction Model

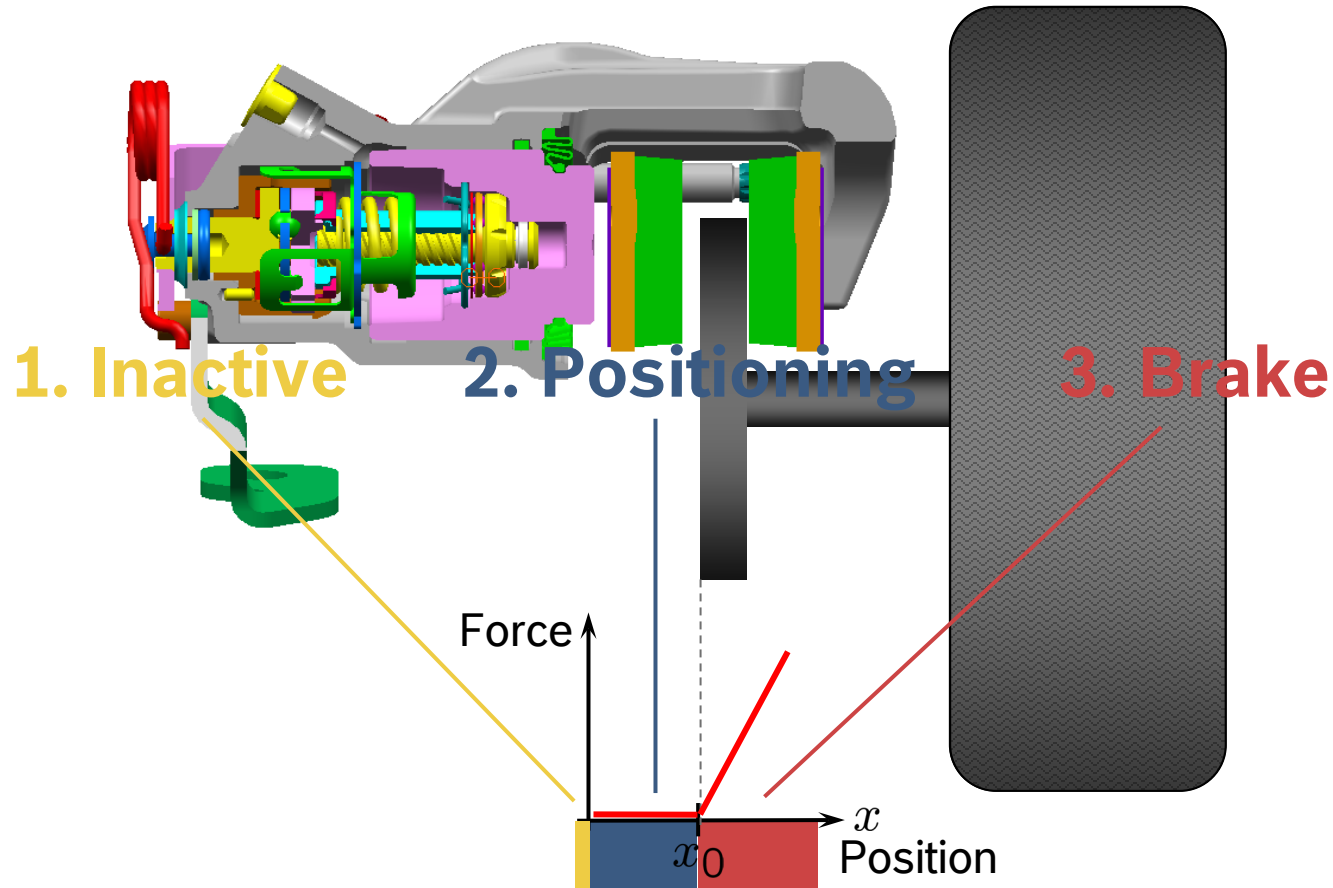


Electro Mechanic Braking System



BOSCH

Electro Mechanic Braking System

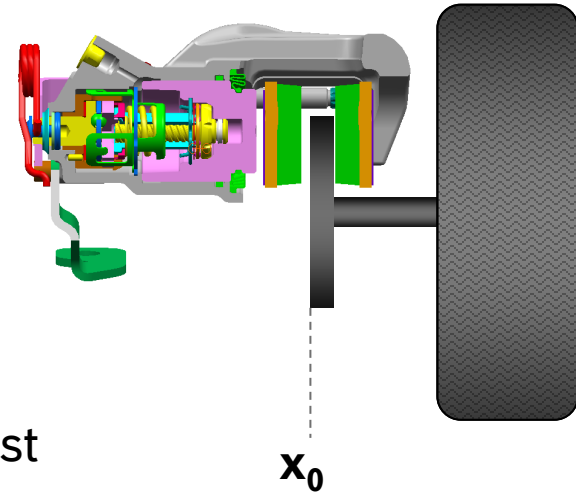


BOSCH

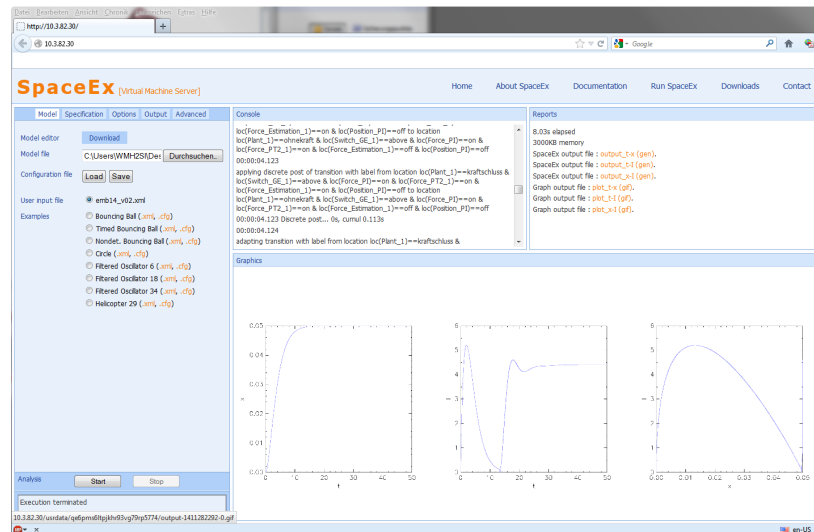


Functional Requirements

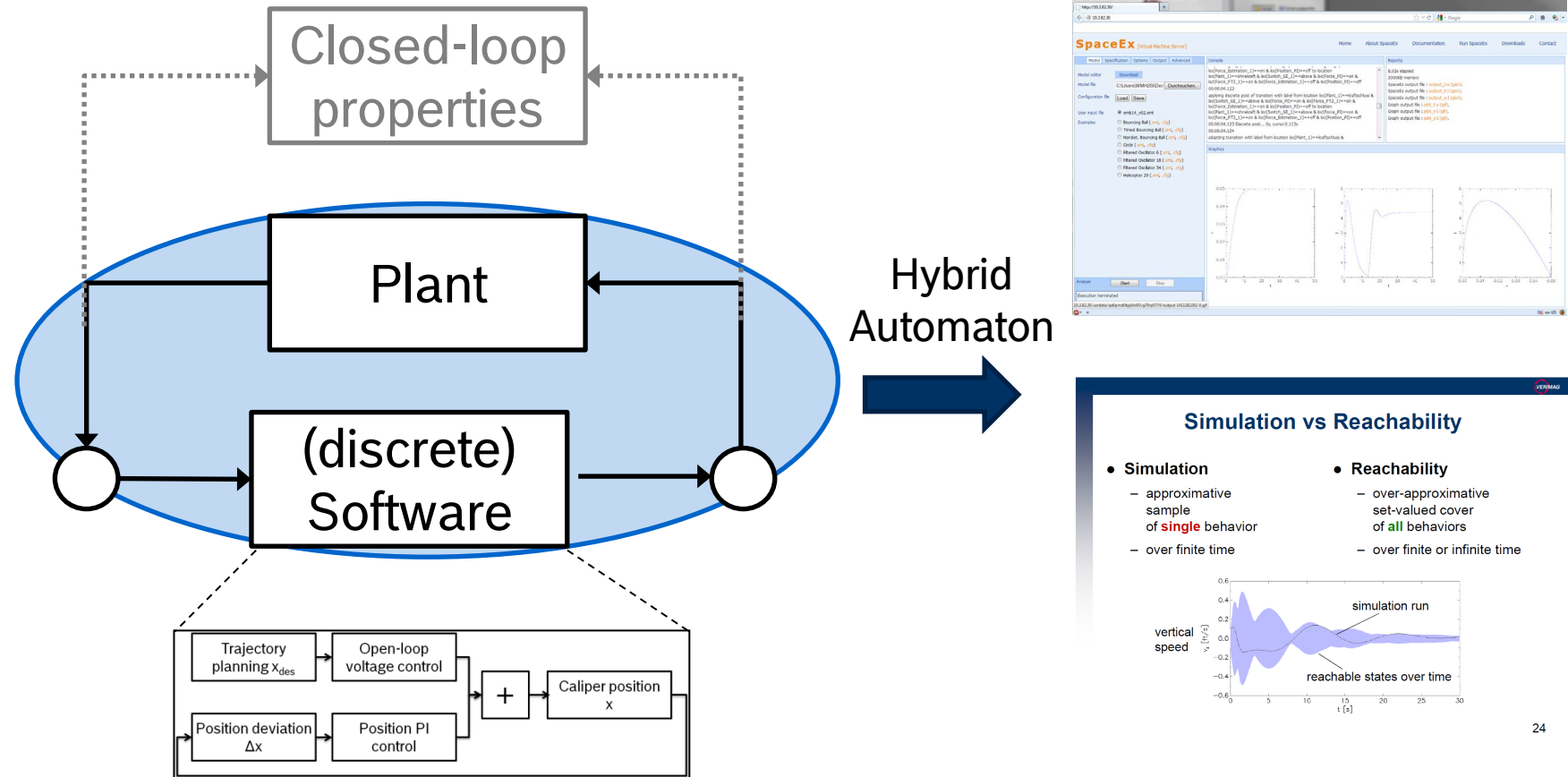
- “Ready-to-brake” position $x_0 = 5 \text{ mm}$
 - Preparation of braking system for applying brake force, no force closure
- **Req. 1: Short response time**
 - Reactiveness of the system
 - Caliper must be at x_0 after the braking request is issued within 20ms with a precision of 4%
- **Req. 2: Small impulse before braking**
 - Driver feels an abrupt deceleration
 - The caliper speed at contact must be below 2mm/s
 - Might be acceptable for braking, but not in other scenarios , e.g. disk wiping



Formal analysis using hybrid automata and reachability analysis



Functional Verification with ZET* Assumption

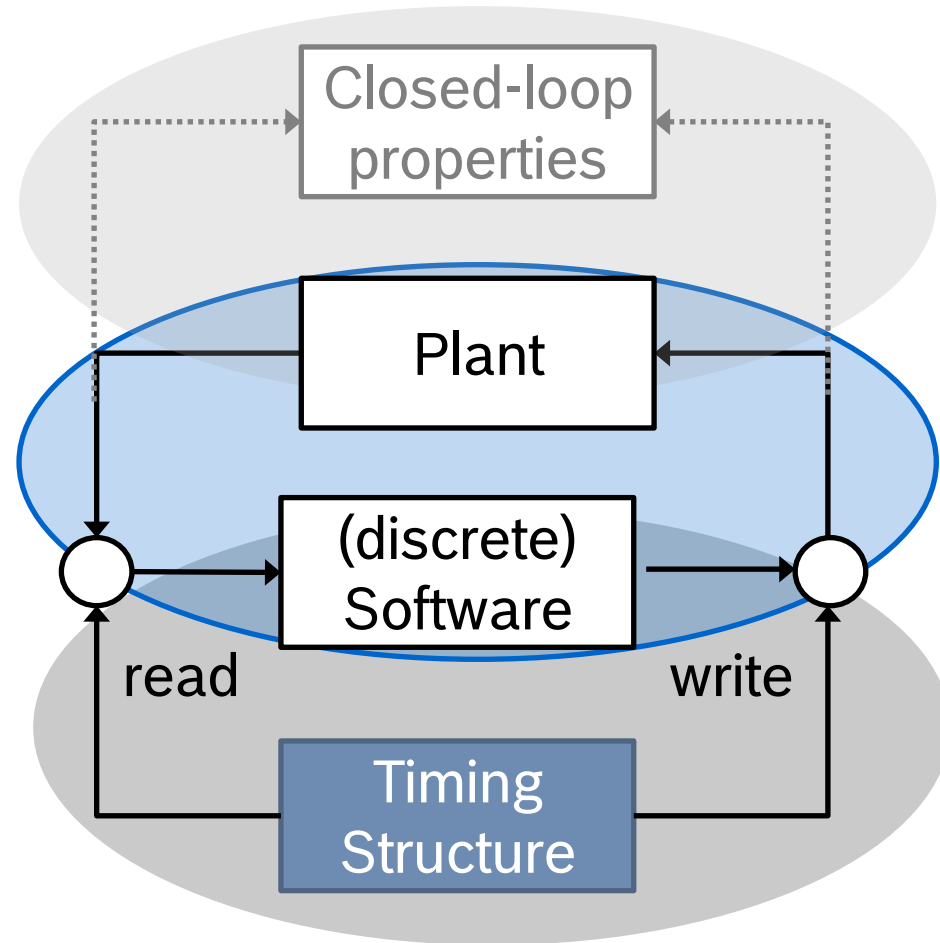


*ZET = Zero Execution Time



BOSCH

Functional Verification considering Timing

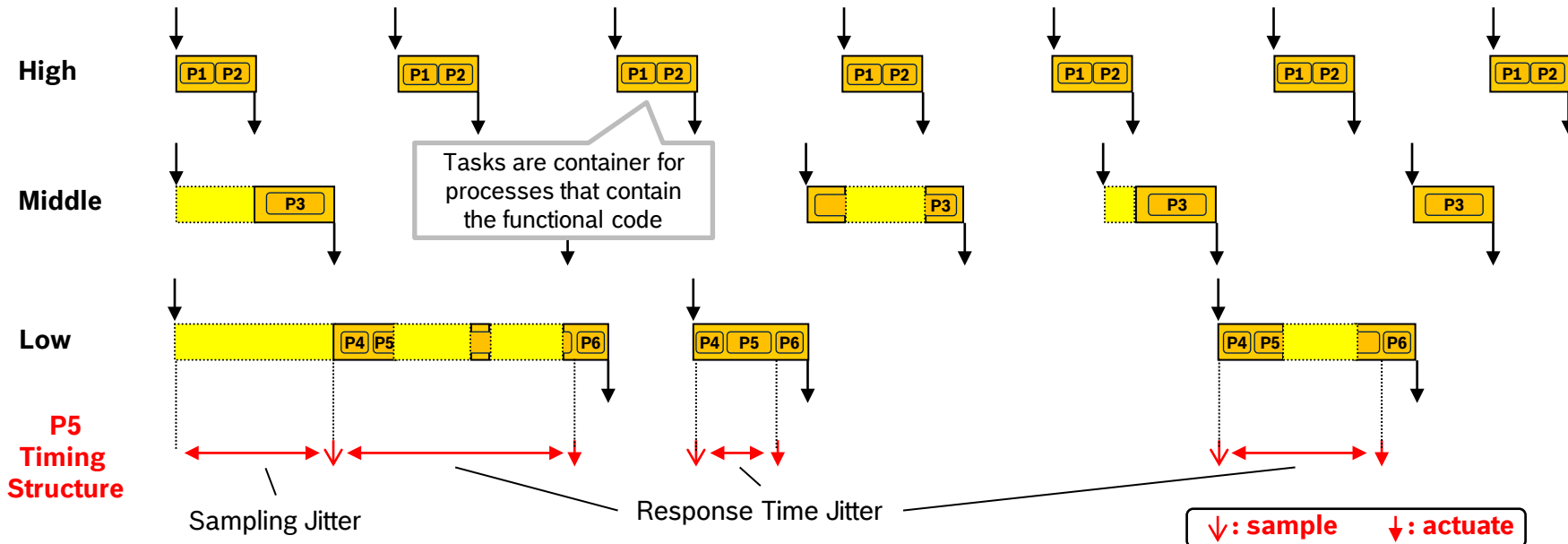
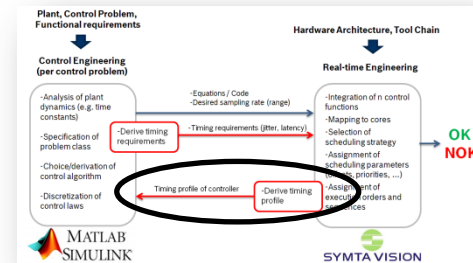


- Model Timing in Hybrid Automaton
 - When is data written / read
 - Non-deterministic model
- Possible models
 - Logical Execution Times
 - Arrival Curves
 - Typical Worst-Case Models
 - ...
- Drivers for choosing a model
 - Generality / analysis trade-off
 - Decision to simplify design for verifiability
 - Functional requirements



Timing Structure – OSEK Systems

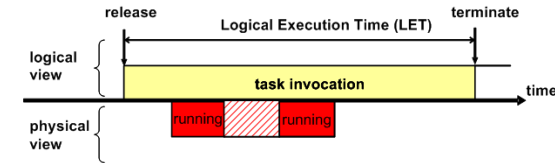
- Description of points in time where the plant is sampled, and where the actuation takes place
- Assumption: functionality implemented by a single process
- Example: Bosch Engine Management
 - Copy-in of required data at task release
 - Copy-out of produced data at process completion



Which Timing Model to choose?

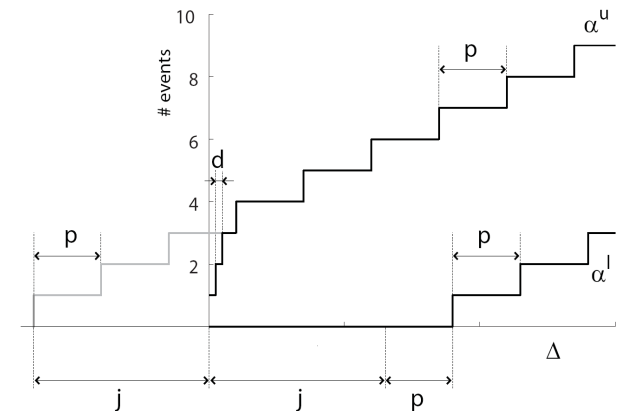
→ LET ?

- Trade Jitter against Latency → Determinism
- Great simplification of verification task
- Ok for “robust” control tasks based on exact models and little external disturbances



→ Arrival Curves?

- Precise model of possible system timing behavior
- Large space of possible timings
- Closed-loop verification very difficult



→ Typical Worst-Case Model !

- Allows for trade-off between both models



Typical Worst-Case Analysis

→ Principle

- Identify typical bounds for the behavior of a system and how often the system may leave these bounds

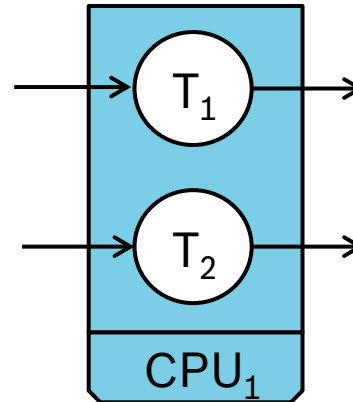
→ Output for each task

- a “safe” bound on its response times: SWCRT
- a typical bound: TWCRT
- a function err such that out of every k consecutive executions, at most $\text{err}(k)$ response times may be larger than TWCRT

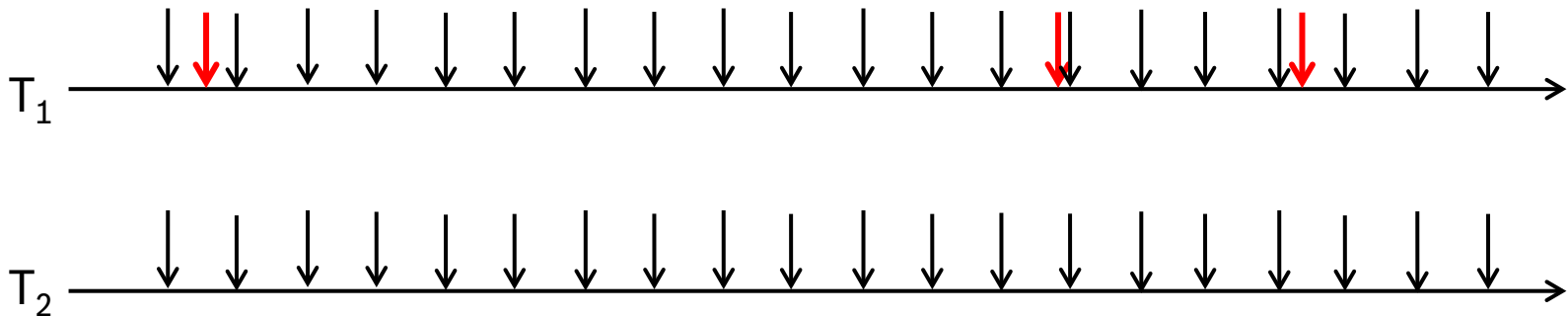
→ Advantages

- Approach is computationally very efficient
- **m-out-of-k** constraints are easy to understand
- No assumptions w.r.t. dependencies

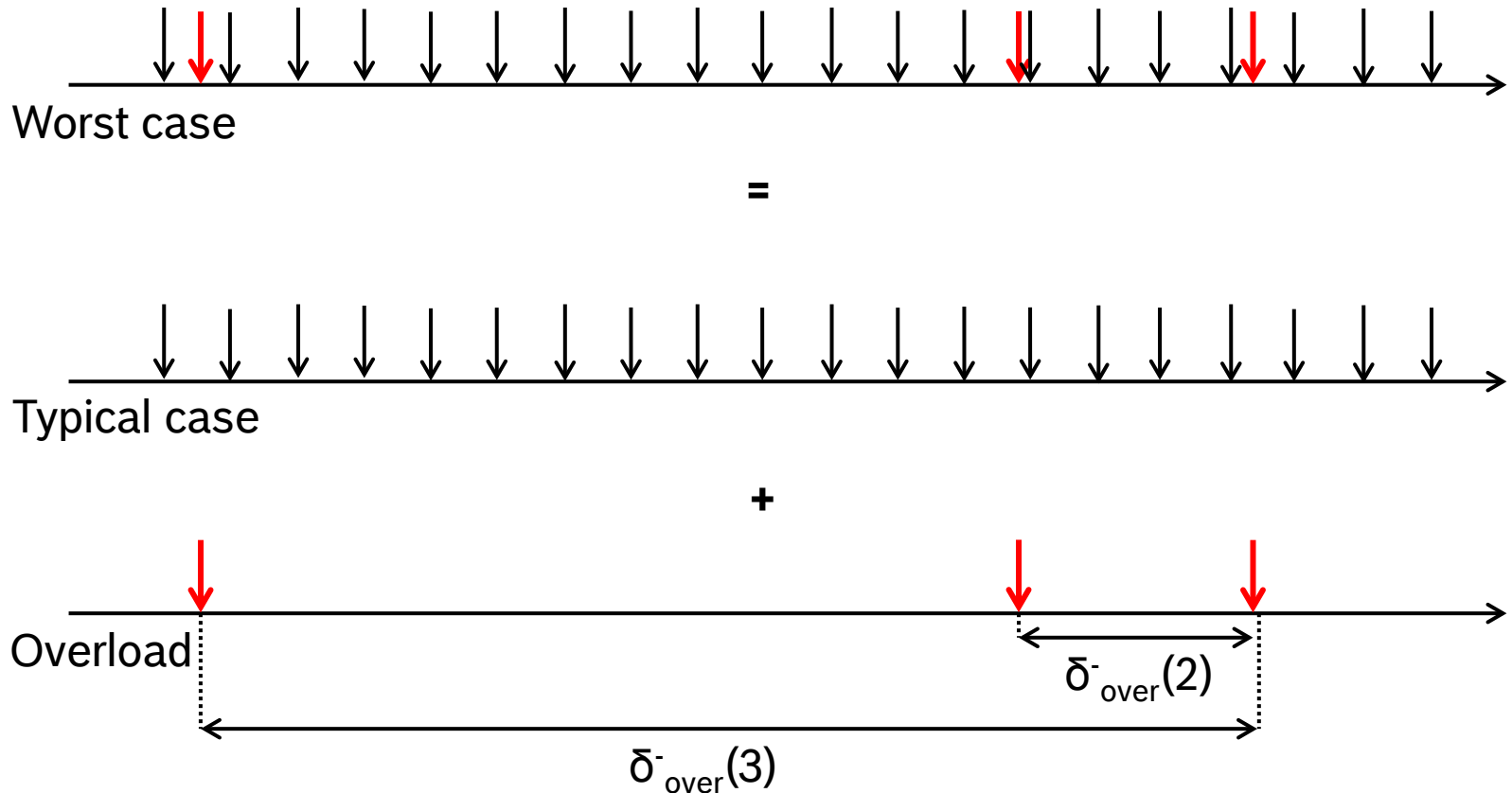
Formal Analysis of Sporadic Overload



Scheduling policy: SPP
(Static Priority Preemptive)
 $T_1 > T_2$



Modeling Sporadic Overload



Formal Analysis of Sporadic Overload

→ Input:

1. a worst-case model of the system
2. a typical model ignoring the overload
3. a model of the overload

→ Analysis (for each task):

1. a busy window analysis of the worst-case model
→ **Safe Worst-Case Response Time (SWCRT)**
2. a busy window analysis of the typical-case model
→ **Typical-Case Response Time (TWCRT)**
3. a computation of the error model based on the result of 1. and the overload model
→ **function err** such that out of every k consecutive executions, at most $\text{err}(k)$ response times may be larger than TWCRT



Using TWCRT Model for Closed-loop Functional Verification

- Idea: Data is written to plant deterministically at $\text{TWCRT} \ll \text{WCRT}$ (using LET)
 - Trade-off between determinism & functional requirements
- TWCRT misses are bounded by error function
 - Scalable “discrete” timing model

*Data for EMB
example*

Period = 1 ms

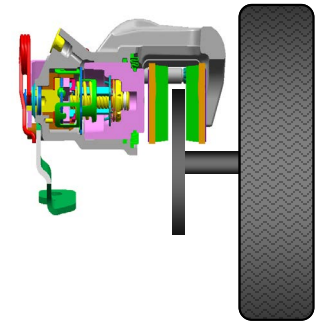
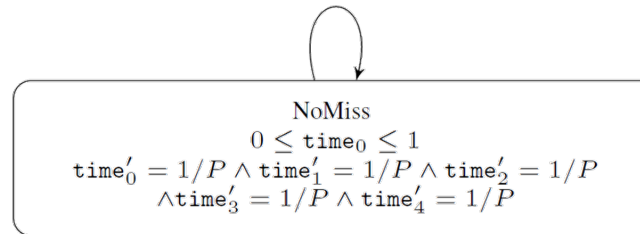
WCRT = 0.8 ms

TWCRT = 0.4 ms

# deadline misses	consecutive executions
2	2
3	18
4	20
5	56



deadline_miss
 $\text{time}_0 \geq 1 \wedge \text{time}_1 \geq \text{miss}(2) \wedge \text{time}_2 \geq \text{miss}(3)$
 $\wedge \text{time}_3 \geq \text{miss}(4) \wedge \text{time}_4 \geq \text{miss}(5)$
 $\text{time}_4 := \text{time}_3 \wedge \text{time}_3 := \text{time}_2 \wedge \text{time}_2 := \text{time}_1$
 $\text{time}_1 := \text{time}_0 \wedge \text{time}_0 := 0$

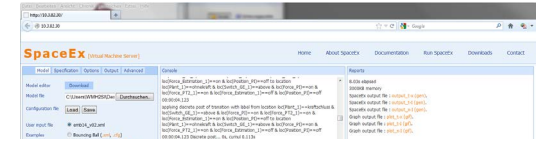
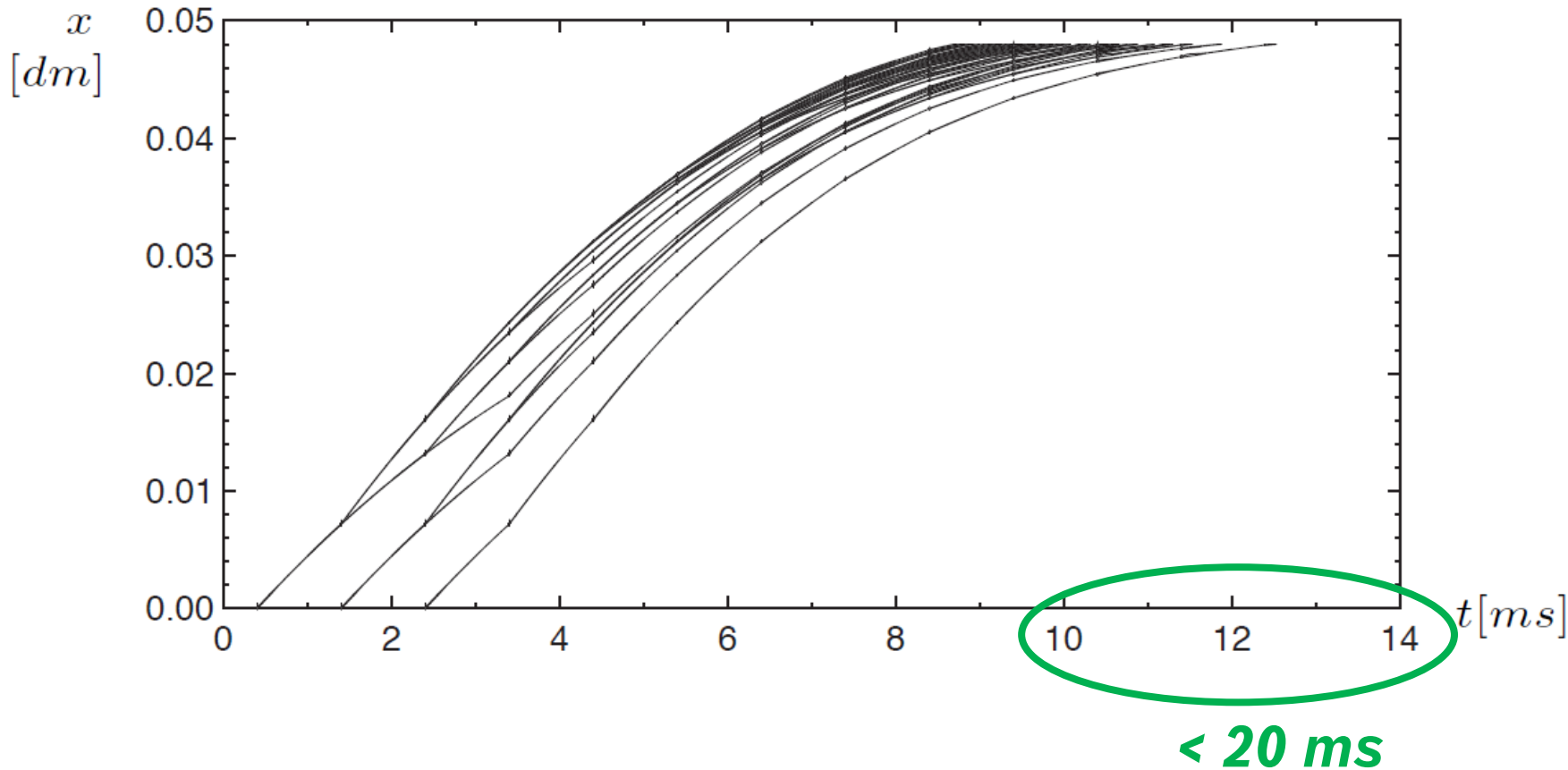


(to be published RTSS 2014)

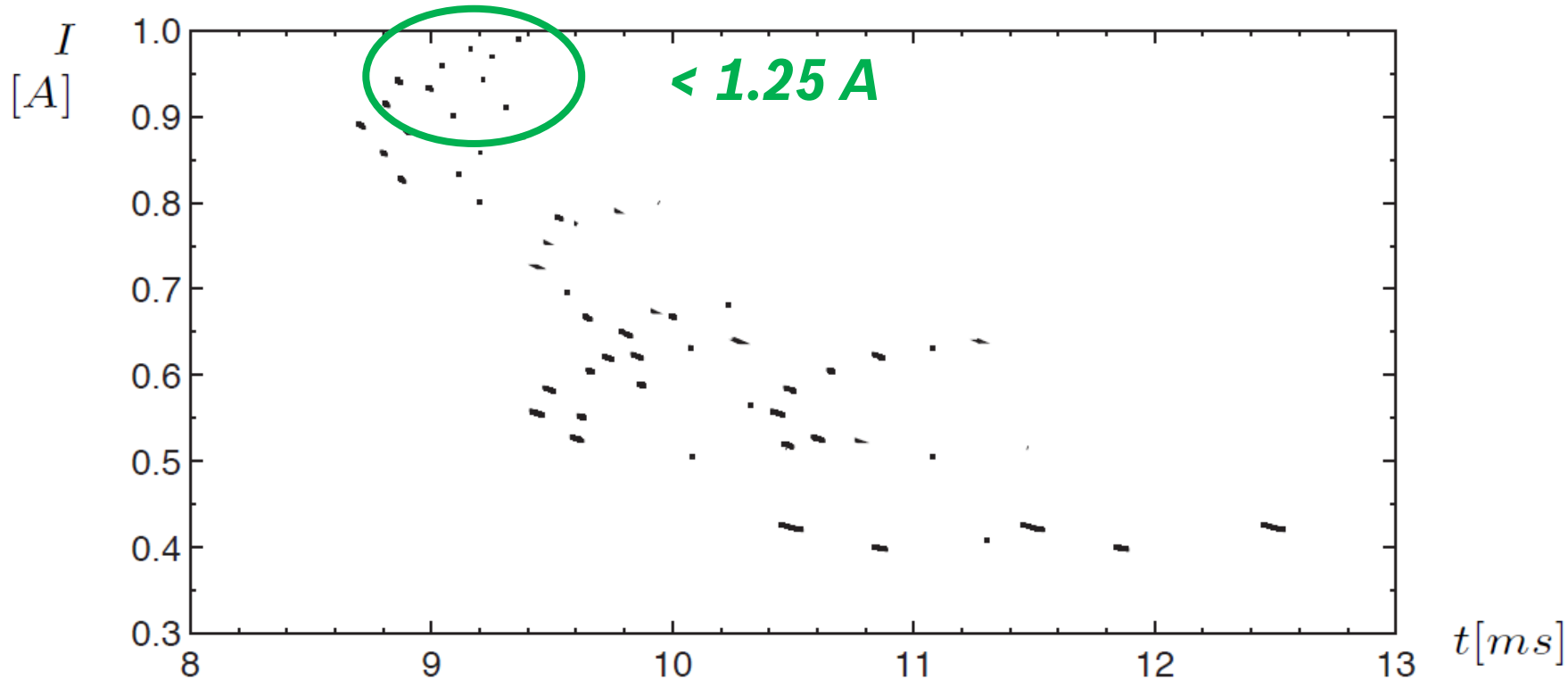


BOSCH

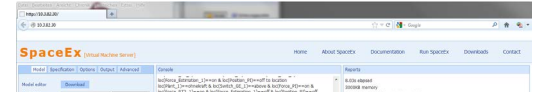
Requirement 1: Response time



Requirement 2: Small impulse



- Current I proportional to the caliper velocity
- Intersection reachable states with the plane of contact
- Bounds $[0.38, 0.99]$ satisfies the requirement 2.



BOSCH

Conclusion

- Both control and real-time engineers have idealized system models for physical systems
- Functional integration effects are not considered by both disciplines
 - Integration effects are anticipated with overdesign
 - ...but even then, functional correctness cannot be guaranteed
- Reachability analysis for hybrid automata is an adequate tool to verify closed loop properties under timing influences
 - Recent advances allow analysis of industrial strength applications
- One promising approach to close the gap between control and real-time system engineering
 - Verify correctness and performance of control software
 - Derive timing requirements for system synthesis



Questions ???

Formal Analysis of Timing Effects on Closed-loop Properties of Cyber Physical Systems

Arne Hamann, Corporate Research, Robert Bosch GmbH

Joint work with: Matthias Wöhrle (Bosch), Goran Frehse (Université Joseph Fourier Grenoble),

Sophie Quinton (INRIA Grenoble)



BOSCH