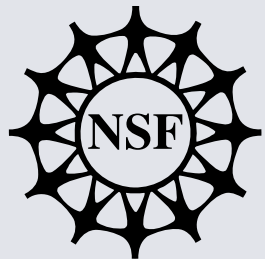


# Model-Based Design

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Chess Review  
October 4, 2006  
Alexandria, VA



# Model-Based Design



Model-based design focuses on the *formal representation, composition, and manipulation of models* during the design process.



# System Composition Approaches



## Component Behavior

### **Modeled on different levels of abstraction:**

- Generalized transition systems (FSM, Time Automata, Cont. Dynamics, Hybrid), fundamental role of time models
- Precise relationship among abstraction levels
- Research: dynamic/adaptive behavior

## Interaction

### **Expressed as a system topology :**

- Module Interconnection (Nodes, Ports, Connections)
- Hierarchy
- Research: dynamic topology

### **Describes interaction patterns among components:**

- Set of well-defined Models of Computations (MoC) (SR, SDF, DE,...)
- Heterogeneous, precisely defined interactions
- Research: interface theory (time, resources,..)

## Scheduling/ Resource Mapping

### **Mapping/deploying components on platforms:**

- Dynamic Priority
- Behavior guarantees
- Research: composition of schedulers

# Tool Composition Approaches



**Domain-Specific Tools;  
Design Environments**

**Domain-Specific Design Flows and  
Tool Chains:**

- ECSL - Automotive
- ESML - Avionics
- SPML - Signal Processing
- CAPE/eLMS

**Metaprogrammable  
Tools, Integration  
Frameworks**

**MIC Metaprogrammable Tool Suite:  
(mature or in maturation program)**

- Metamodeling languages
- Modeling Tools
- Model Transformations
- Model Management
- Design Space Construction and Exploration
- Tool Integration Framework

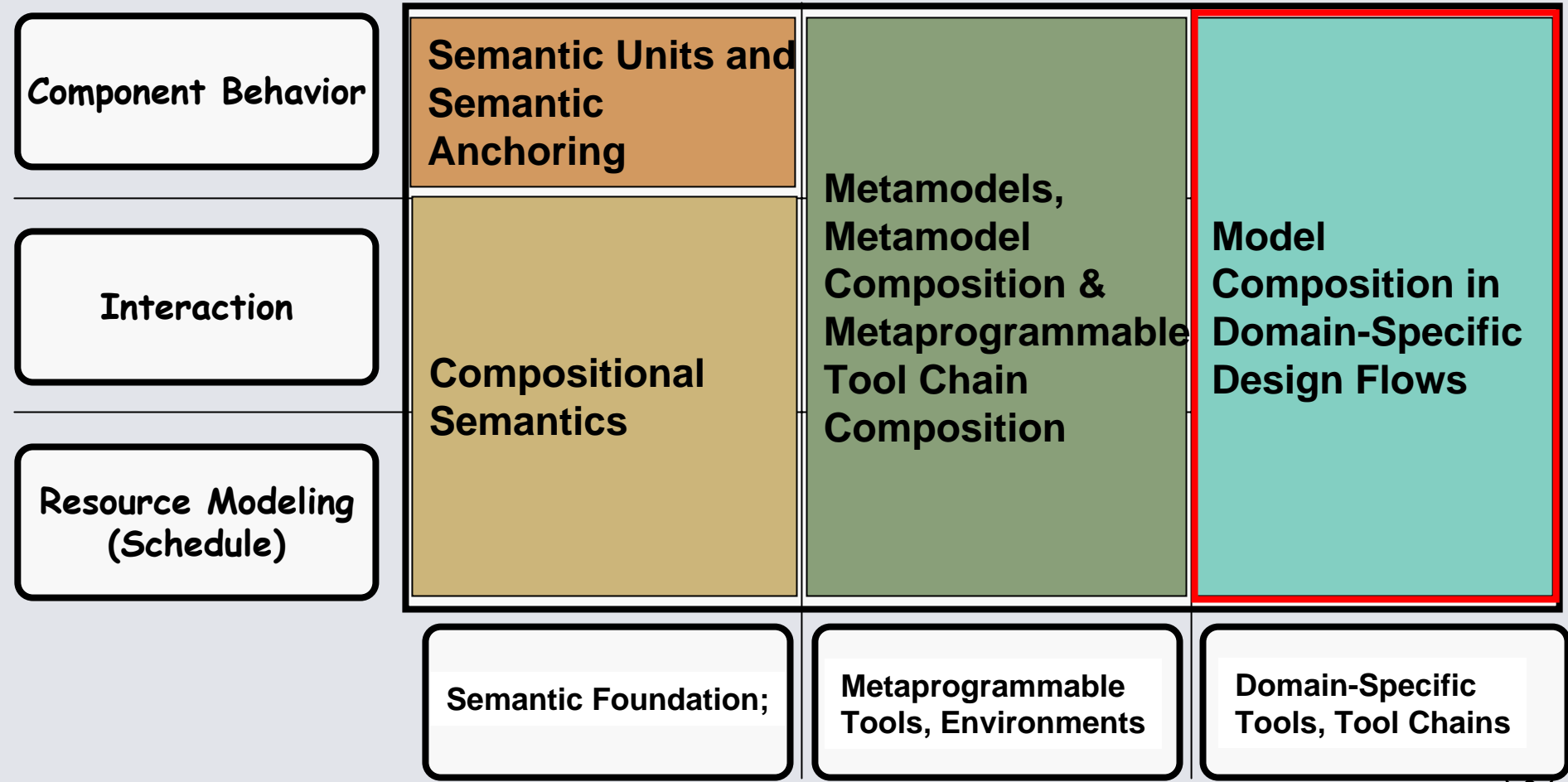
**Semantic Foundation**

**Semantic Foundations (work in progress):**

- Semantic Anchoring Environment (SAE)
- Verification
- Semantic Integration



# Intersection of System and Tool Composition Dimensions



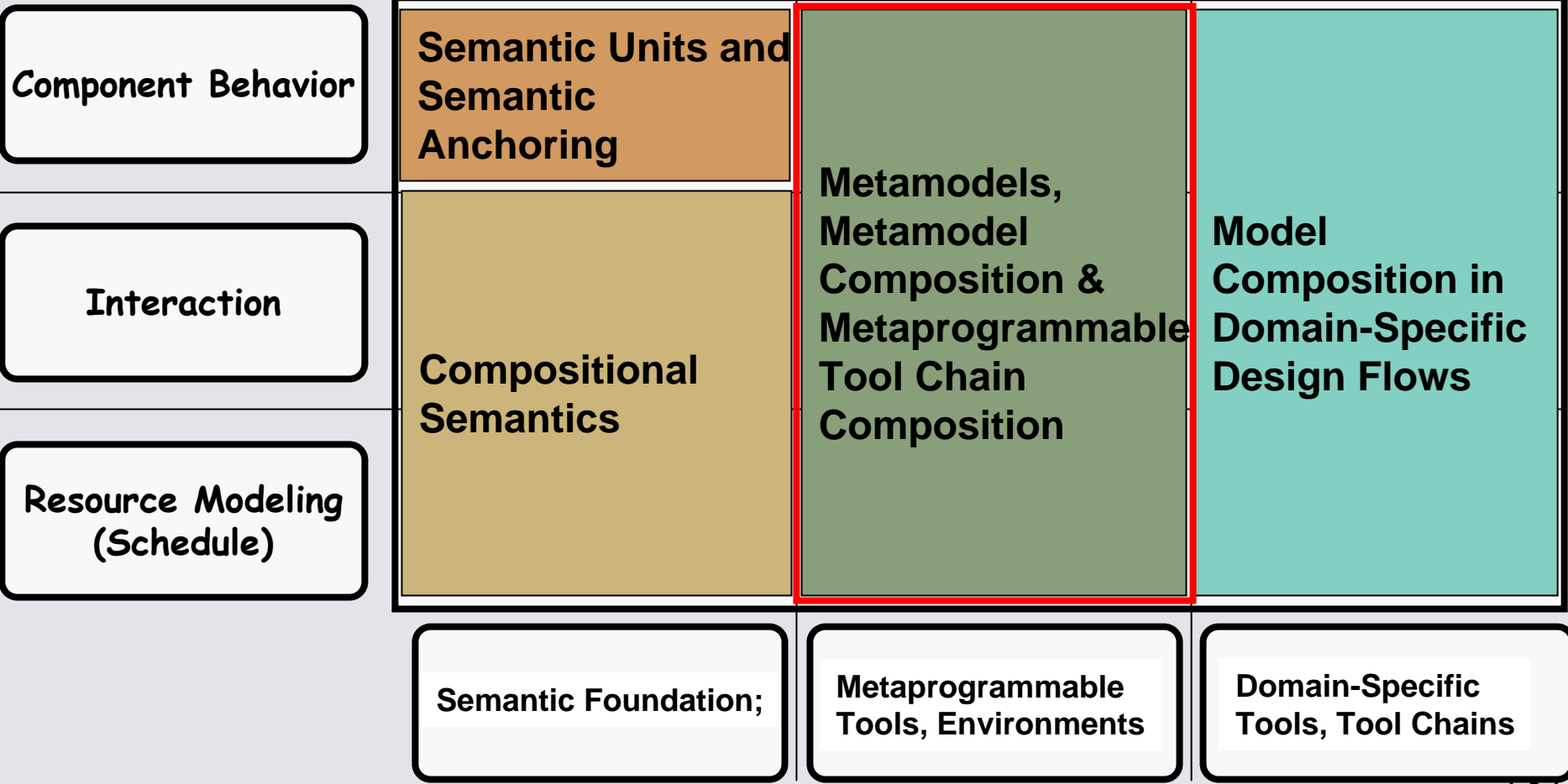
# Domain Specific Design Flows and Tool Chains



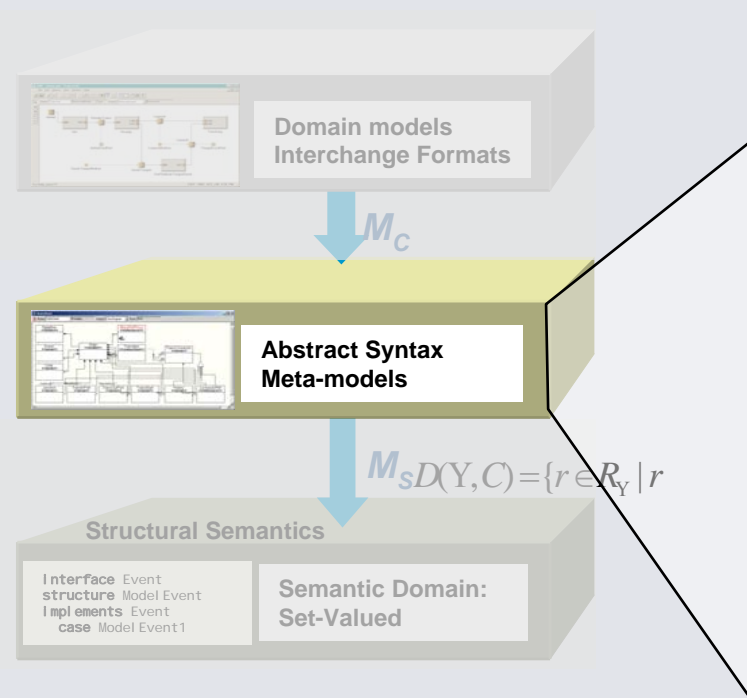
- Integration of tools into tool chains
  - ECSL - Control
  - ESML - Avionics
  - SPP - Signal Processing
  - FCS - Networked Embedded Systems
  - SCA - Software Defined Radio
- Integration among tool frameworks:  
Metropolis, Ptolemy II, MIC,  
Simulink/Stateflow, ARIES, CheckMate,...
- [www.escherinstitute.org](http://www.escherinstitute.org)



# Intersection of System and Tool Composition Dimensions



# Syntactic Layer

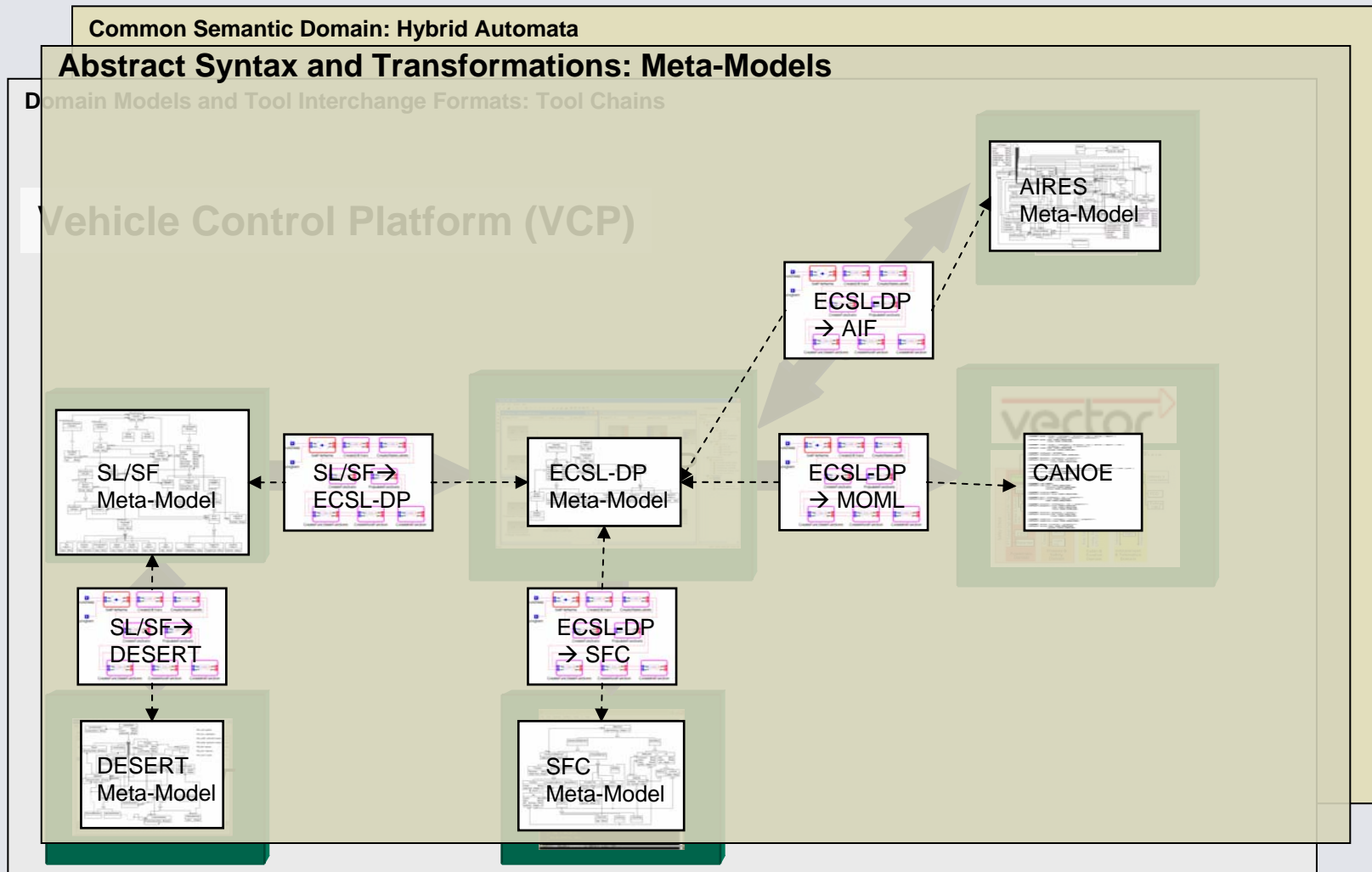


- Modeling & Metamodeling
- Model Data Management
- Model Transformation
- Tool Integration
- Design-Space Exploration

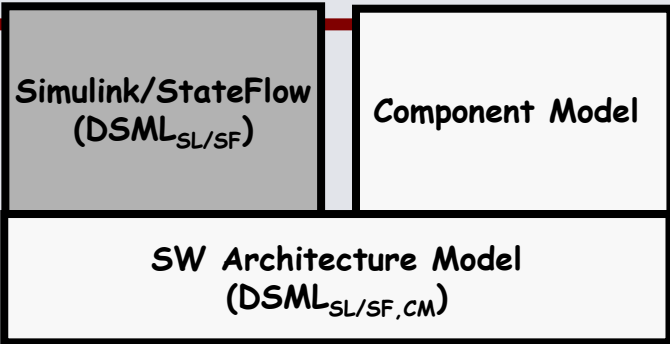




# Metamodeling View of a Tool Chain



# Need for Metamodel Composition:

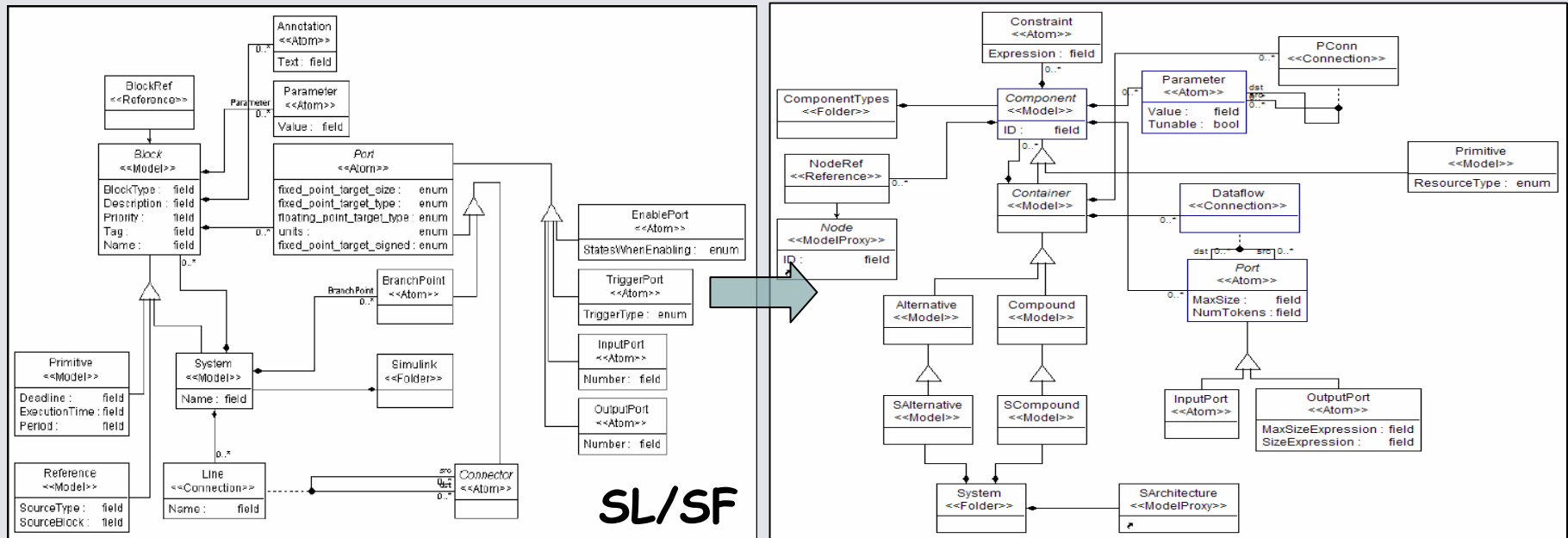


**Objective:** Optimize the SW architecture by selecting a component model and by allocating functions to components.

**Platform:** Heterogeneous Dataflow Component Model

**Tools:**

GME, GReAT, C Compiler, WCET Analyzer



## Functional blocks - SW Component Mapping



# Solutions for Compositional Metamodeling



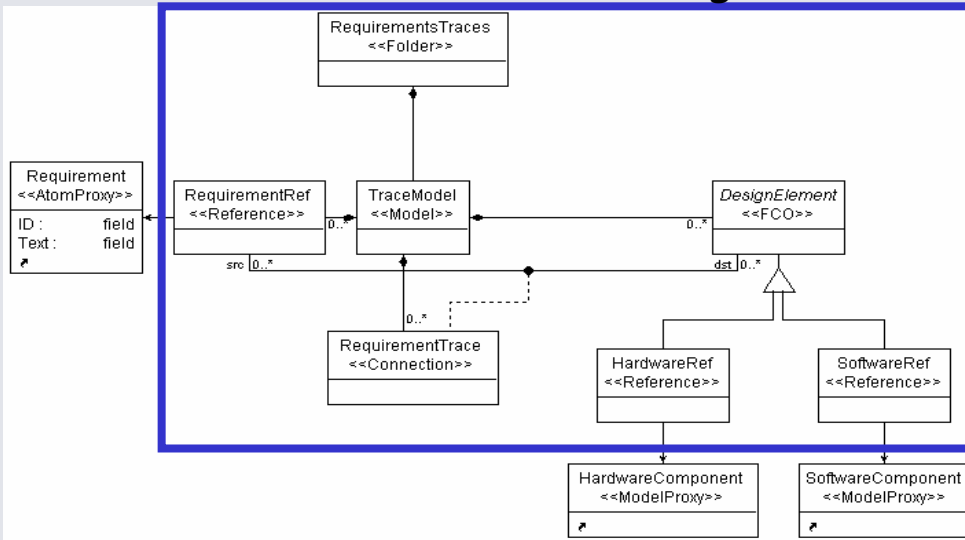
- Goal: Composing modeling languages (not models)
- Metamodel composition methods in the Generic Modeling Environment (GME):
  - Class Merge
  - Metamodel Interfacing
  - Class Refinement
  - Template Instantiation
  - Metamodel Transformations



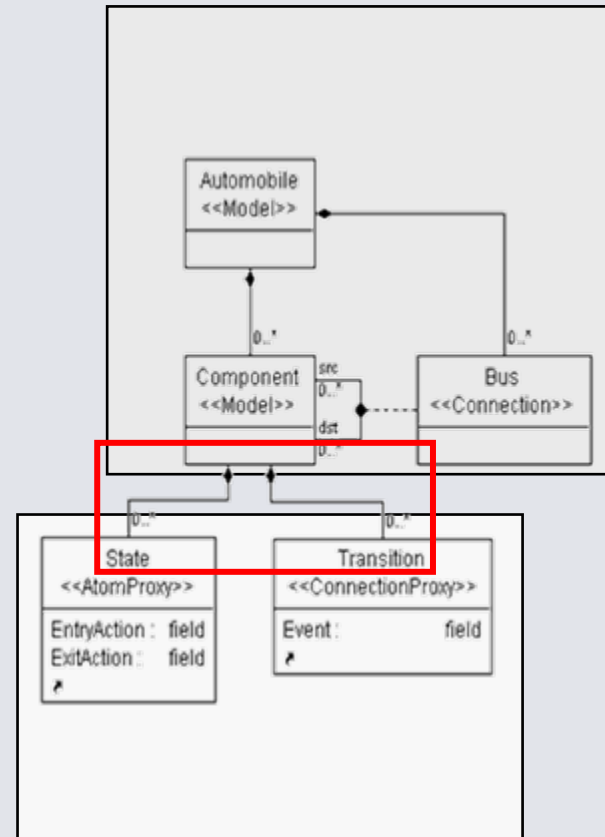
# Metamodel Composition Methods



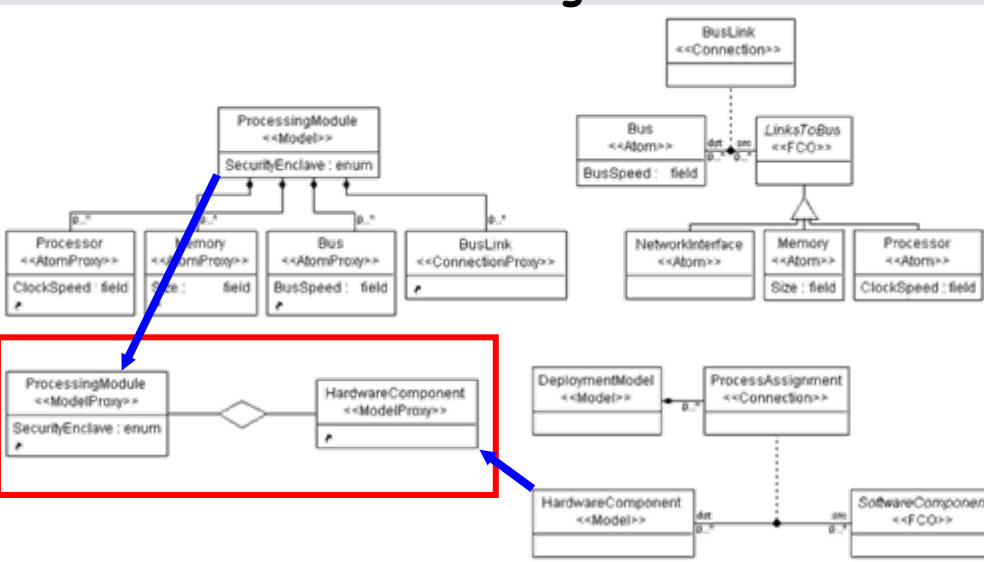
## Metamodel Interfacing



## Class Refinement



## Class Merge



# Summary of Progress in Model Transformations



Complex model transformations can be formally specified in the form of executable graph transformation rules

G/T semantics is very powerful but the implementation needs to be tailored for efficiency

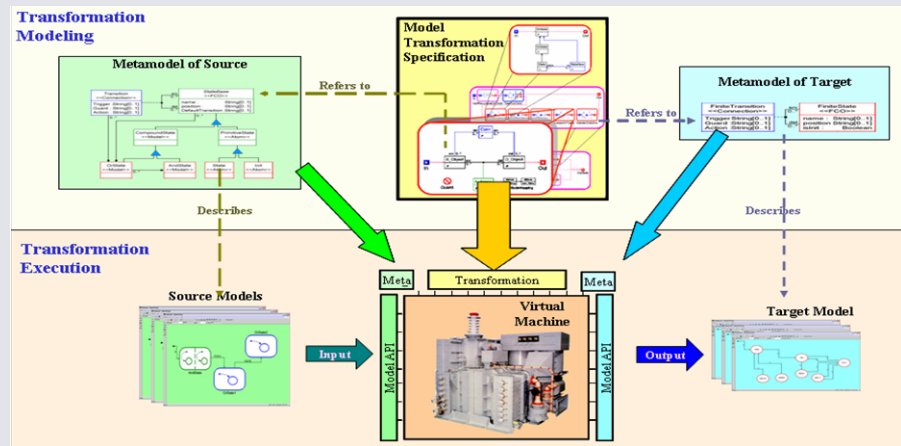
GReAT is an open source, metamodel-based model transformation language supported by tools: modeling tool, rewriting engine, code generator and debugger. It is based on attributed/typed graph matching, multi-domain rewriting rules, and explicitly sequenced rewriting operators.

Highlights of GReAT extensions: shared spaces, sorting of match results, cross-products of matches, higher-order operators (groups)

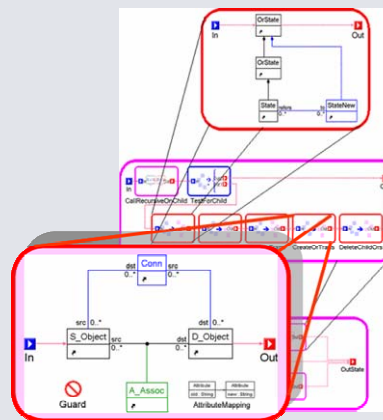
Applications of GReAT:

- Simulink/Stateflow verifying code generator
- Several model transformation tools in embedded system toolchains
- Semantic anchoring of domain-specific modeling languages

## Concept: Metamodel-based Transformations



## Language: Graph Transforms



## Toolsuite: GReAT

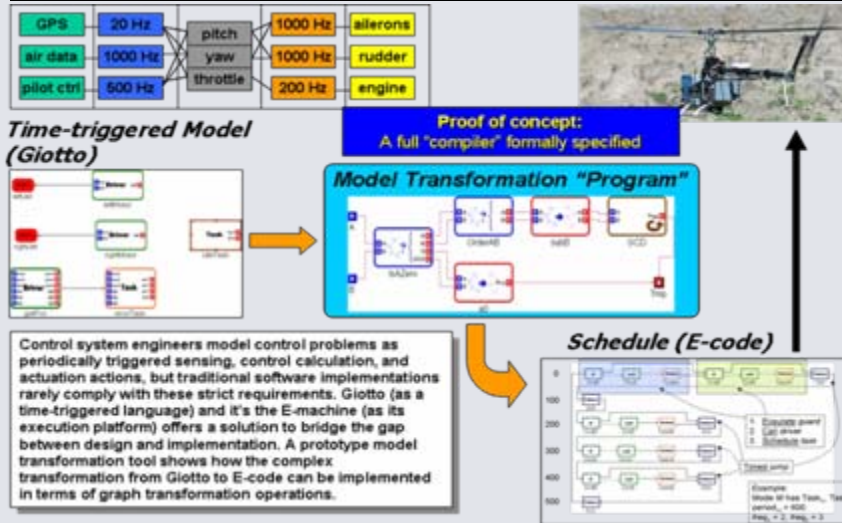


(Karsai et al, 2005-2006)

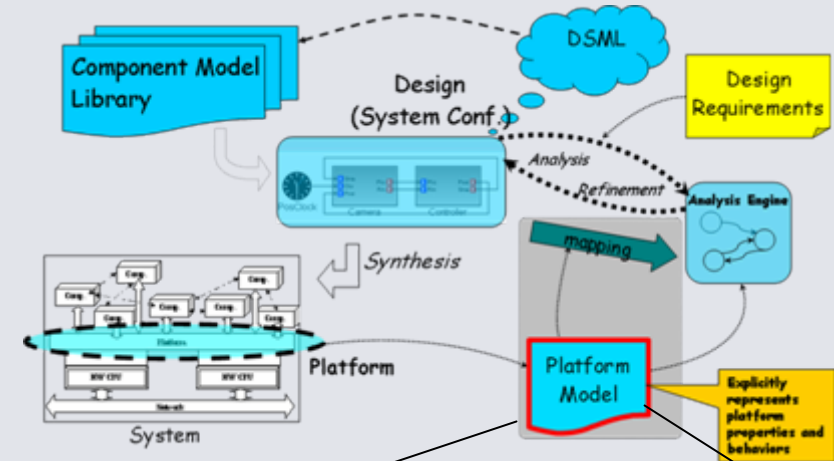
# Major Applications of Model Transformations



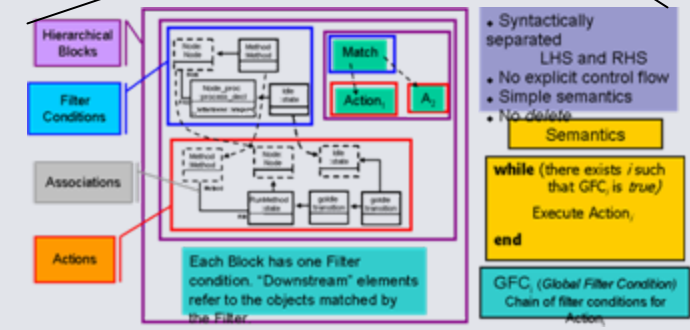
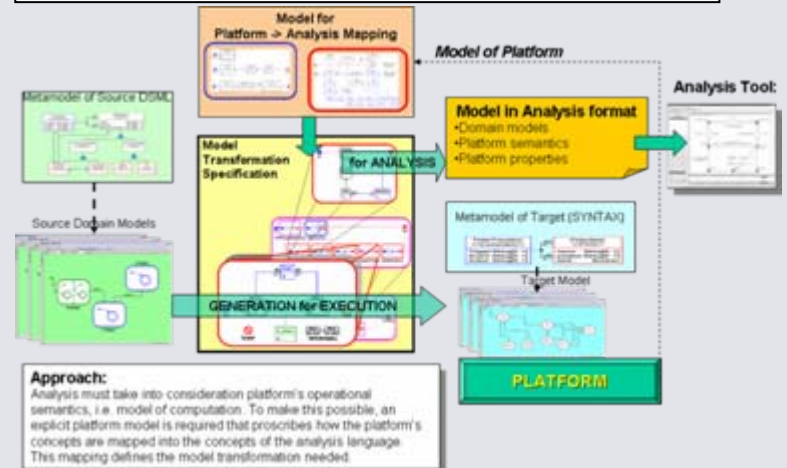
## Model Transformations for Schedule Generation:



## Explicit Platform Modeling Language:



## Implicit Platform Modeling for Analysis:



(Karsai et al, 2005-2006)

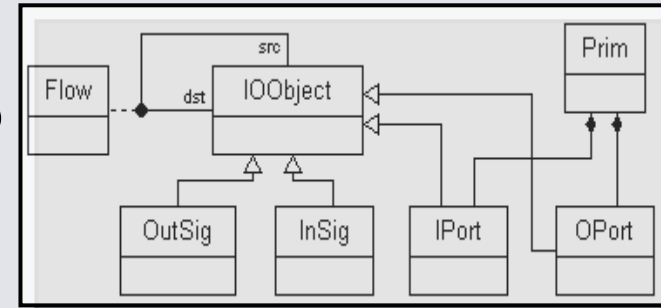




# Structural Semantics of Models and Metamodels



We followed a formal logic approach to structural semantics. A metamodel is mapped to a set of n-ary function symbols and constraints over an associated Herbrand Universe.



## Primitives of DSP Domain

$$\Upsilon = \begin{cases} \text{insig}(X) : X \text{ is an input signal} \\ \text{outsig}(X) : X \text{ is an output signal} \\ \text{prim}(X) : X \text{ is a basic DSP operation} \\ \text{iport}(X, Y) : X \text{ has an input port } Y \\ \text{oport}(X, Y) : X \text{ has an output port } Y \\ \text{inst}(X, Y) : X \text{ is the DSP operation } Y \\ \text{flow}(X_1, Y_1, X_2, Y_2) : \text{Data goes from opor} \\ \text{t } Y_1 \text{ on } X_1 \text{ to ipor} \\ \text{t } Y_2 \text{ on } X_2 \end{cases}$$

## Some Constraints with NAF

- 1 Instances must use primitives that are defined:  
 $\text{inst}(x, \text{prim}(y)) \wedge \neg \text{prim}(y) \Rightarrow \text{malform}(x),$
- 2 Ports are placed on defined primitives:  
 $\text{iport}(\text{prim}(x), y) \wedge \neg \text{prim}(x) \Rightarrow \text{malform}(y),$
- 3 Dataflow connections must start on defined ports:  
 $\text{flow}(x, \text{oport}(\text{prim}(y), z), u, w) \wedge \neg \text{oport}(\text{prim}(y), z) \Rightarrow \text{malform}(z).$

*These are the function symbols and some constraints for the example metamodel*

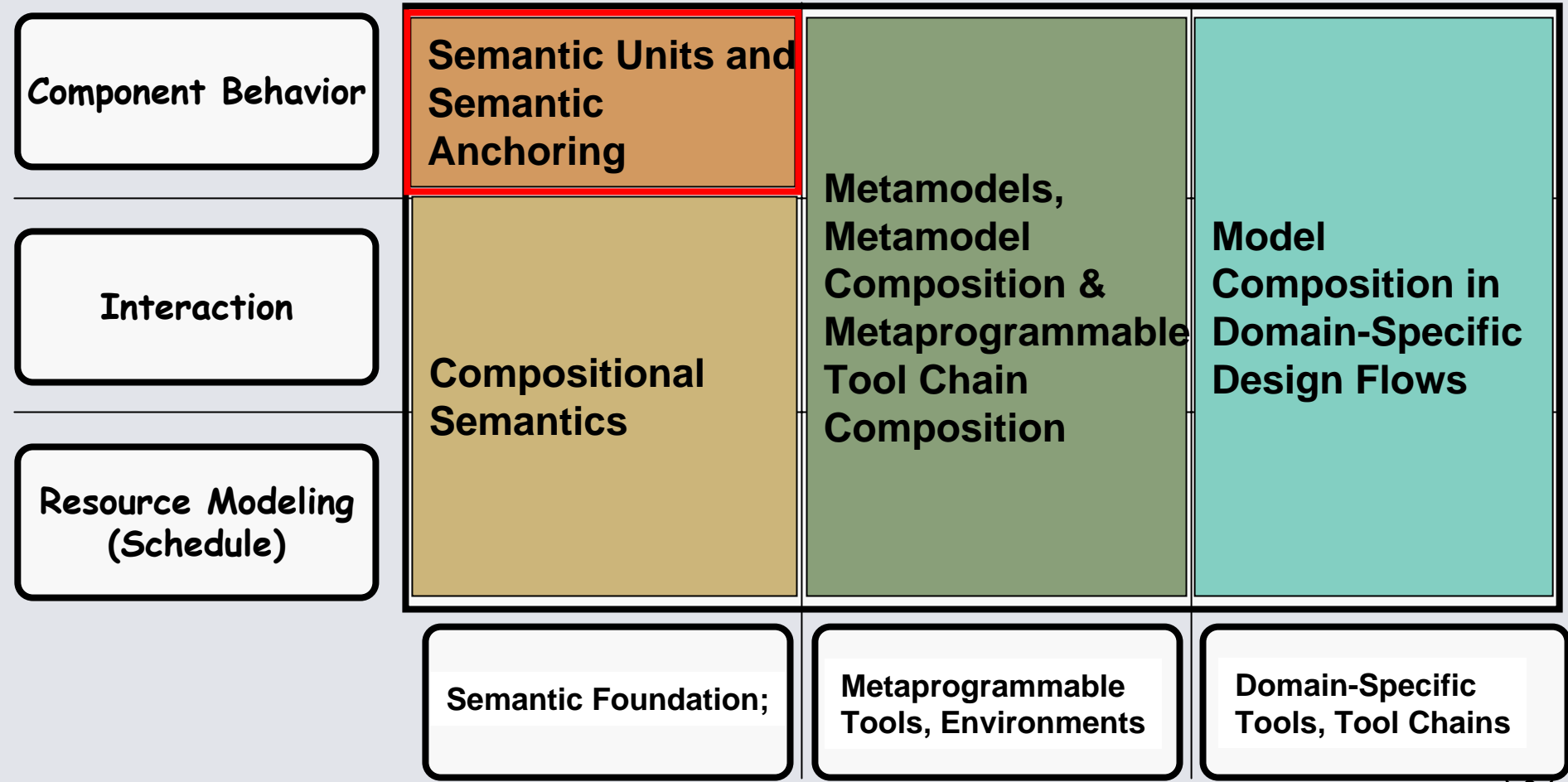
We use an **inference procedure** to prove well-formedness or malformedness. This inference mechanism is well-defined and tool independent.

We have constructed an **automatic theorem prover** that answers questions about structural semantics (see poster).

(Jackson, Sztipanovits 2006)

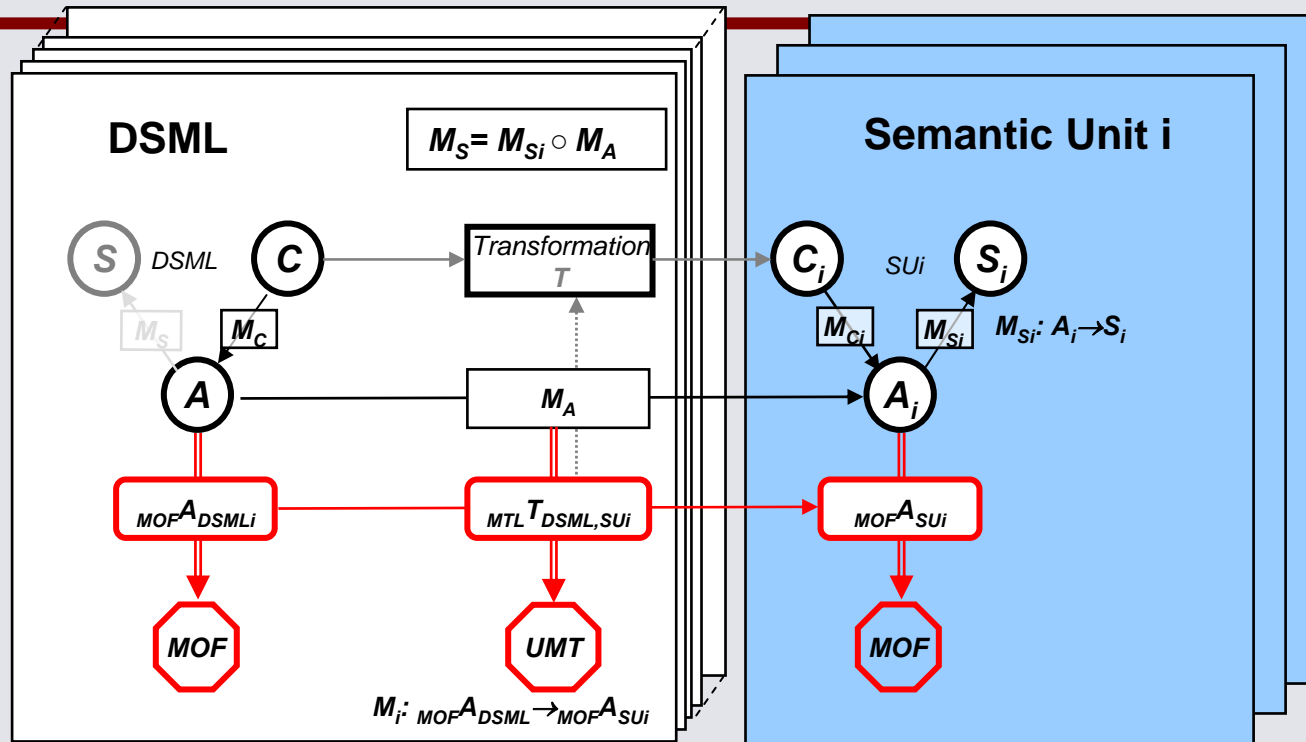


# Intersection of System and Tool Composition Dimensions





# Semantic Anchoring of DSML-s



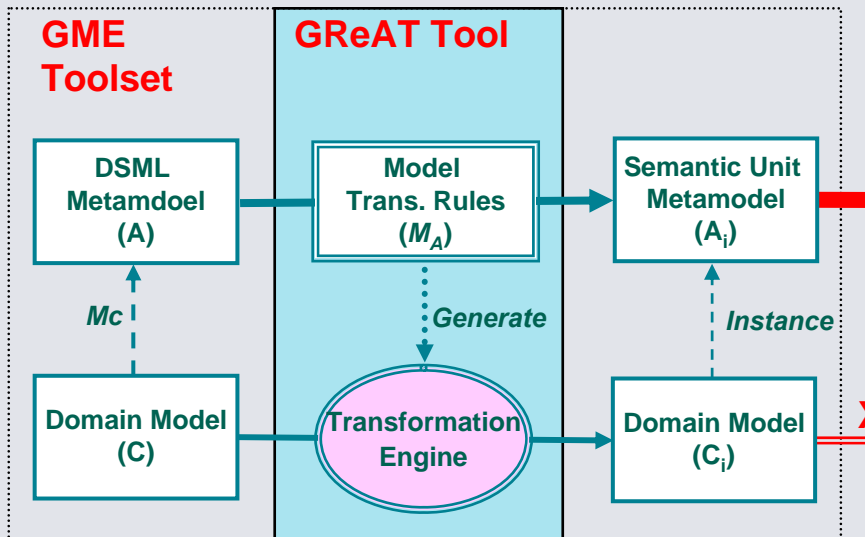
- Step 1
  - Specify the DSML  $\langle A, C, M_C \rangle$  by using MOF-based metamodels.
- Step 2
  - Select appropriate semantic units  $L = \langle A_i, C_i, M_{C_i}, S_i, M_{S_i} \rangle$  for the behavioral aspects of the DSML.
- Step 3
  - Specify the semantic anchoring  $M_A = A \rightarrow A_i$  by using UMT.



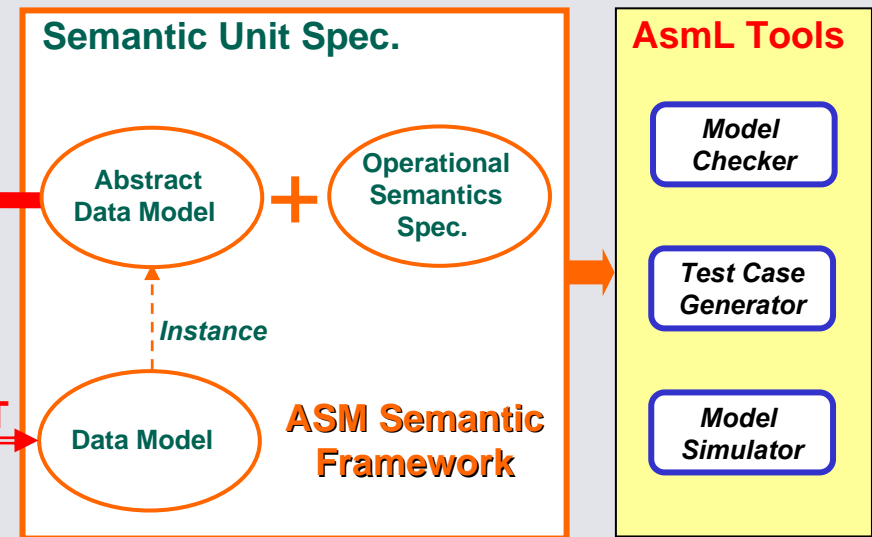
# Experimental Tool Suite for Semantic Anchoring



## Metamodeling and Model Transformation Tools



## Formal Framework for Semantic Units Specification



### Metamodeling and Model Transformation Tools

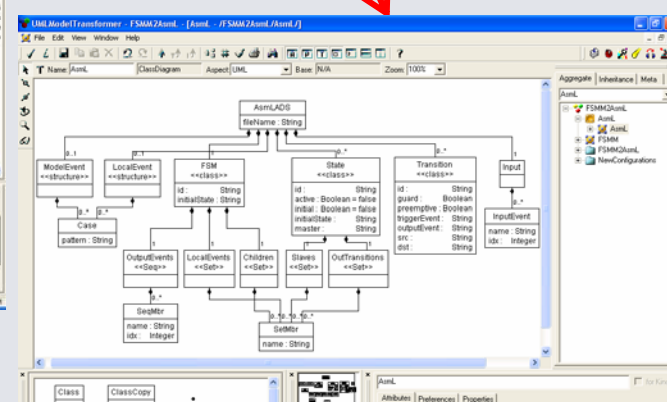
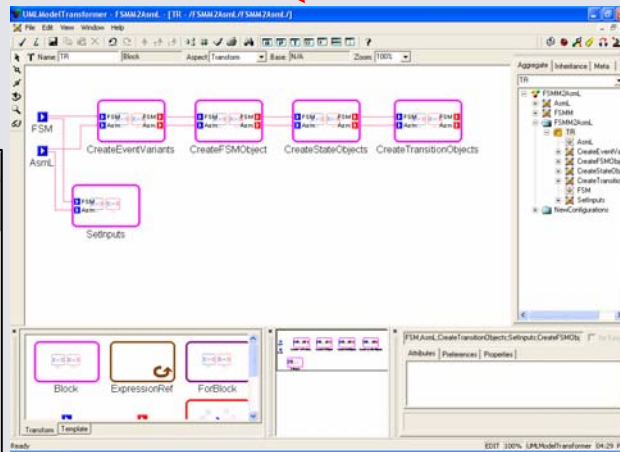
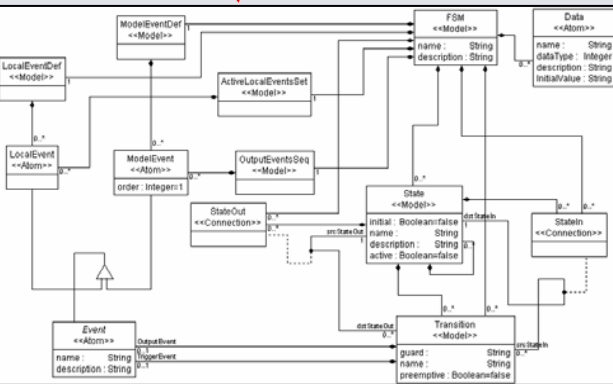
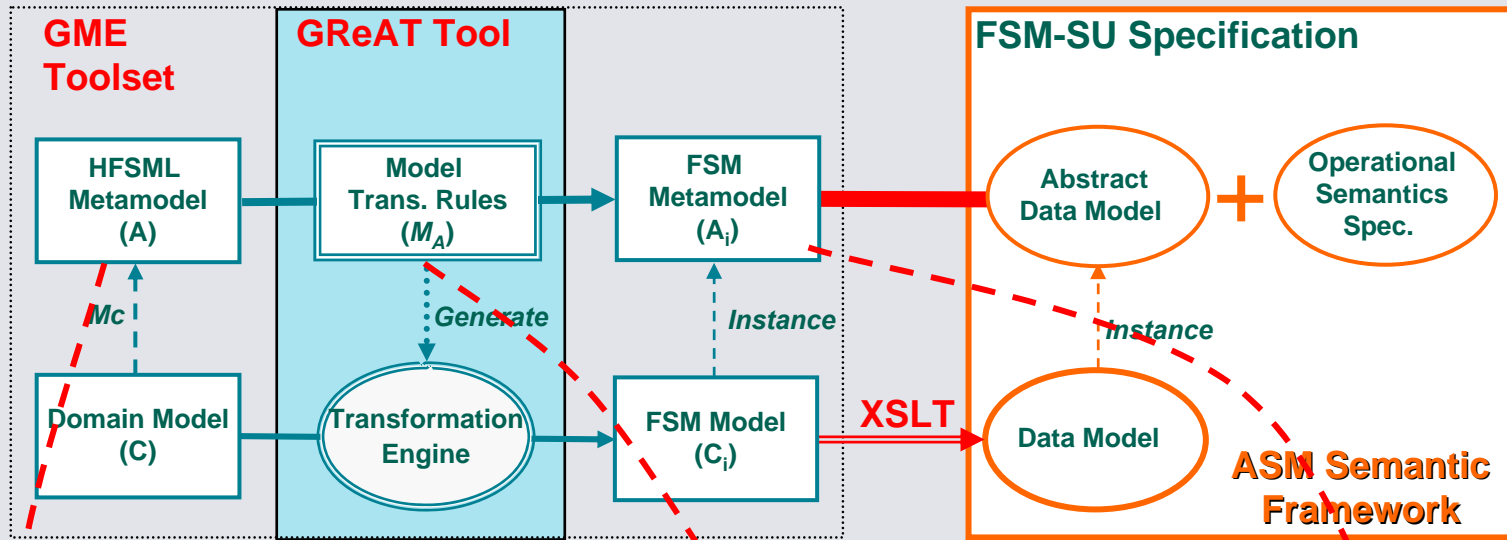
- GME: Provide a MOF-based metamodeling and modeling environment.
- GReAT: Build on GME for metamodel to metamodel transformation.

### Tools for Semantic Unit Specification

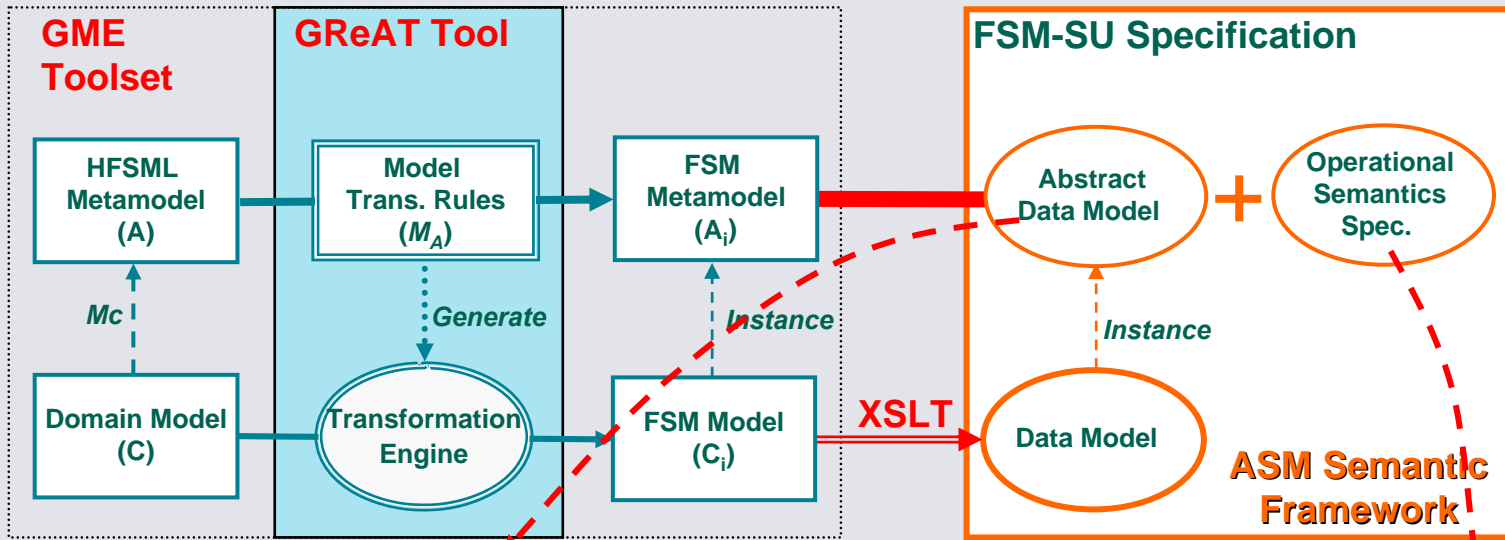
- ASM: A particular kind of mathematical machine, like the Turing machine. (Yuri Gurevich)
- AsmL: A formal specification language based on ASM. (Microsoft Research)



# Example: HFSML -> FSM-SU; 1/3



# Example: HFSML -> FSM-SU; 2/3



```

structure Event
  eventType as String

class State
  id as String
  initial as Boolean
  var active as Boolean = false

class Transition
  id as String

abstract class FSM
  id as String

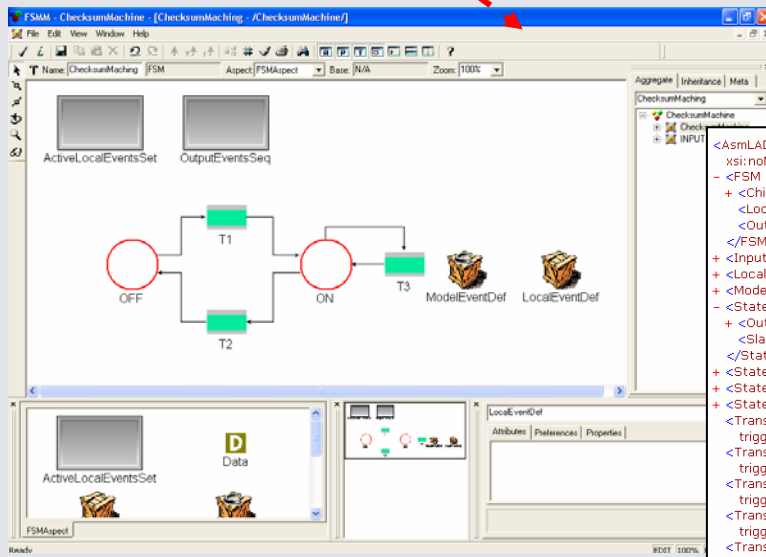
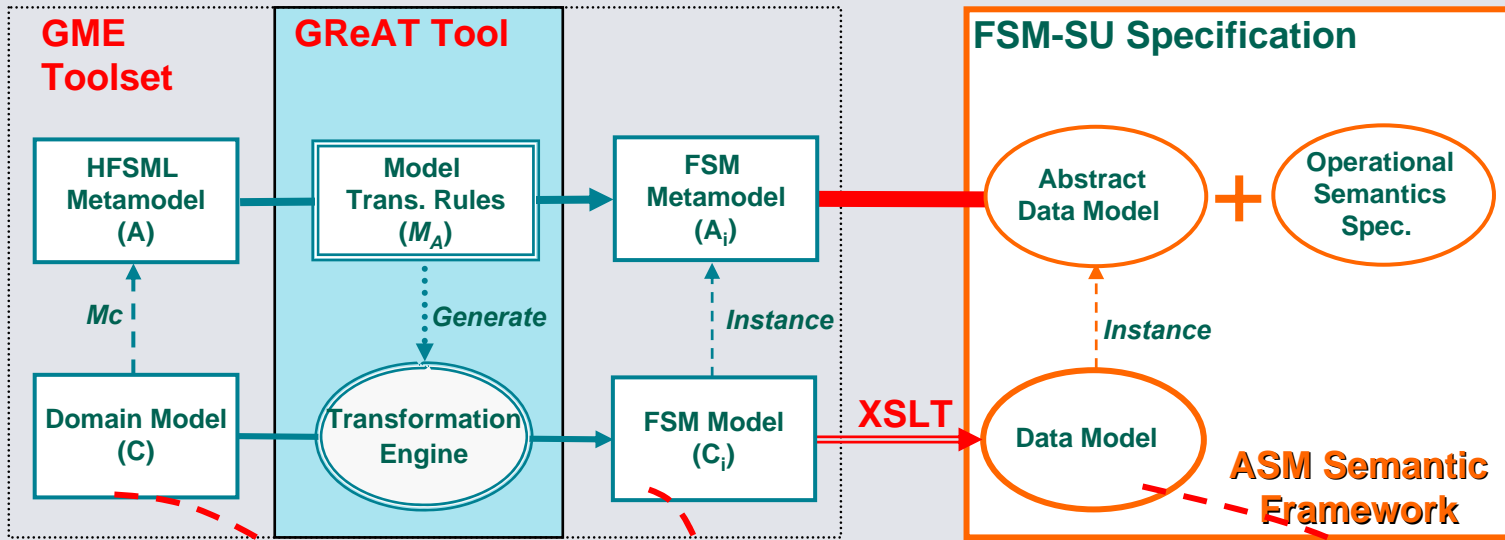
  abstract property states as Set of State
  get
  abstract property transitions as Set of Transition
  get
  abstract property outTransitions as Map of <State, Set of Transition>
  get
  abstract property dstState as Map of <Transition, State>
  get
  abstract property triggerEventType as Map of <Transition, String>
  get
  abstract property outputEventType as Map of <Transition, String>
  
```

```

React (e as Event) as Event?
step
  let CS as State = GetCurrentState ()
step
  let enabledTs as Set of Transition = {t | t in outTransitions (CS) where
  e.eventType = triggerEventType(t)}
step
  if Size (enabledTs) = 1 then
    choose t in enabledTs
      step
        // WriteLine ("Execute transition: " + t.id)
        CS.active := false
      step
        dstState(t).active := true
      step
        if t in me.outputEventType then
          return Event(outputEventType(t))
        else
          return null
    else
      if Size(enabledTs) > 1 then
        error ("NON-DETERMINISM ERROR!")
      else
        return null
  
```



# Example: HFSML -> FSM-SU; 3/3

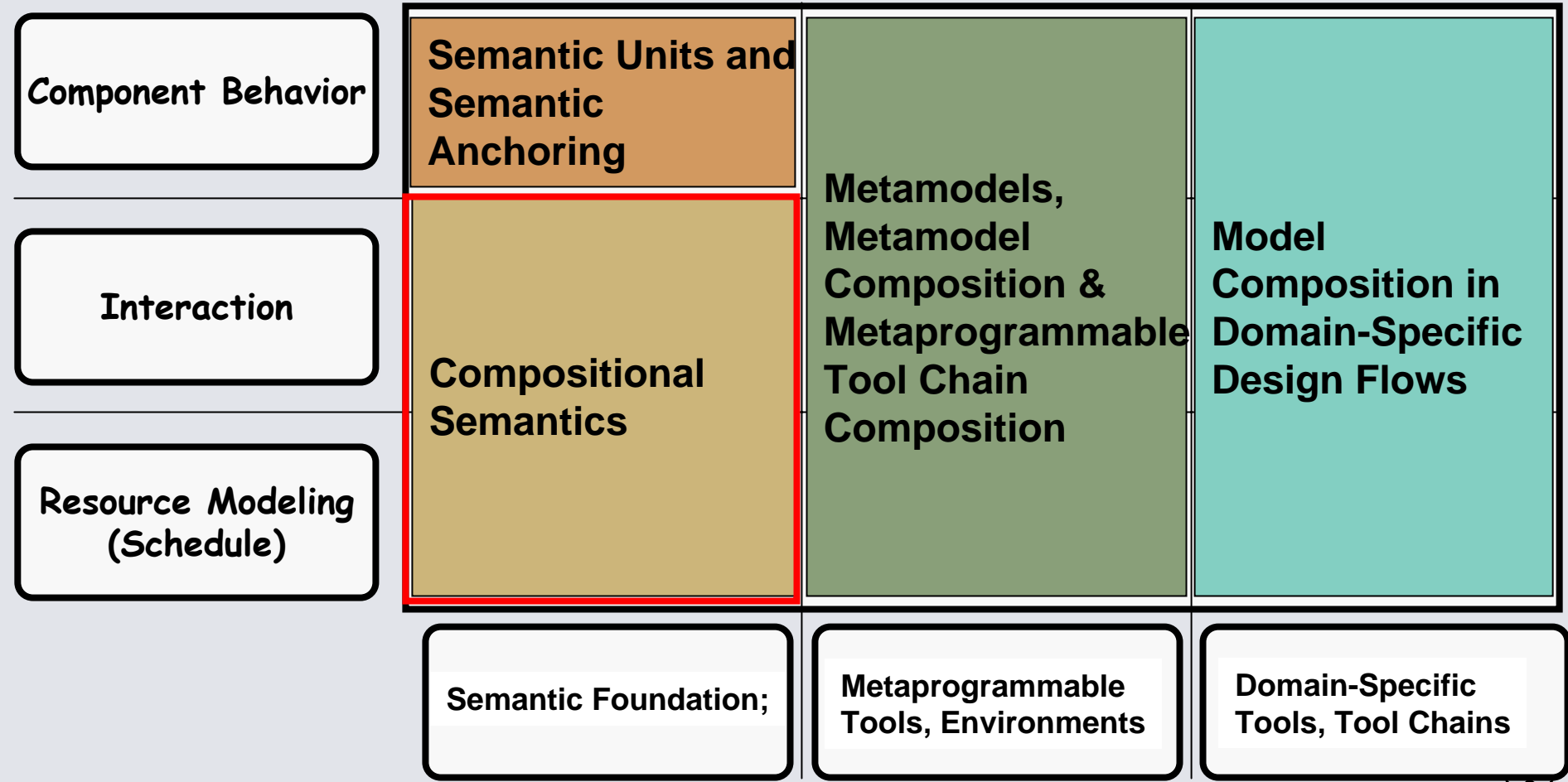


```
<AsmLADS_id="id988" fileName="" xmlns:xsi="http://www.w3.org/
xsi:noNamespaceSchemaLocation="UDM\AsmL.xsd">
- <FSM id="ChecksumMaching" _id="id9d5" initialState="OFF">
+ <Children _id="id9f4">
  <LocalEvents _id="id9e5" />
  <OutputEvents _id="id9e0" />
</FSM>
+ <Input _id="id98b">
+ <LocalEvent _id="id9bf">
+ <ModelEvent _id="id9ad">
- <State id="OFF" _id="ida17" active="false" master="" initial="true"
+ <State id="ON" _id="ida18" active="false" master="" initial="false"
+ <State id="ZERO" _id="ida74" active="false" master="ON" initial="false"
+ <State id="ONE" _id="ida75" active="false" master="ON" initial="false"
- <Transition id="T11" _id="idada" dst="ONE" src="ZERO" guard="true" preemptive="false"
triggerEvent="LocalEvent.one" />
- <Transition id="T12" _id="idadf" dst="ZERO" src="ONE" guard="true" preemptive="false"
triggerEvent="LocalEvent.one" />
- <Transition id="T13" _id="idae0" dst="ZERO" src="ZERO" guard="true" preemptive="false"
triggerEvent="LocalEvent.zero" />
- <Transition id="T14" _id="idae1" dst="ONE" src="ONE" guard="true" preemptive="false"
triggerEvent="LocalEvent.zero" />
- <Transition id="T1" _id="idb0e" dst="ON" src="OFF" guard="true" preemptive="false"
triggerEvent="ModelEvent.start"
triggerEvent="" />
- <Transition id="T2" _id="idb0f" dst="OFF" src="ON" guard="true" preemptive="false"
triggerEvent="ModelEvent.stop" />
- <Transition id="T3" _id="idb10" dst="ON" src="ON" guard="true" preemptive="false"
triggerEvent="ModelEvent.reset" />
</AsmLADS>
```

```
initStateAutomaton() as StateAutomaton
let S1 = State("S1", true)
let S2 = State("S2", false)
let T1 = Transition("T1")
let e1 = Event("e1")
let S = {S1, S2}
let T = {T1}
let E = {e1}
let Connections = {T1 -> (S1, S2)}
let TriggerEvent = {T1 -> e1}
let OutputEvent = {T1 -> e1}
let InitialState = S1
return new StateAutomaton(S, T, E, Connections, TriggerEvent, OutputEvent,
InitialState)
```



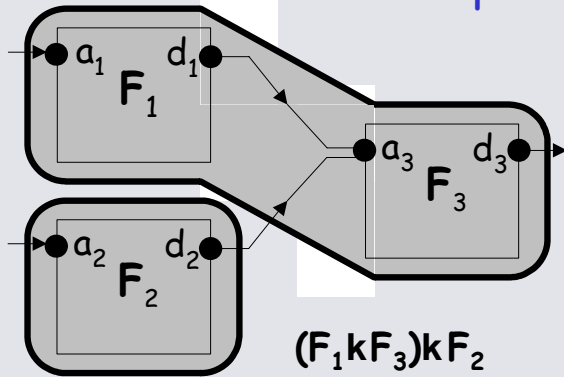
# Intersection of System and Tool Composition Dimensions



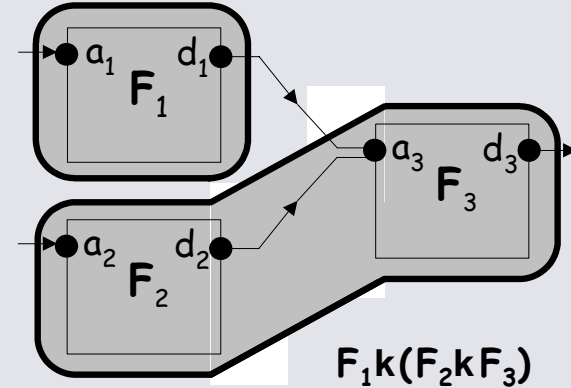
# Component-based Analysis



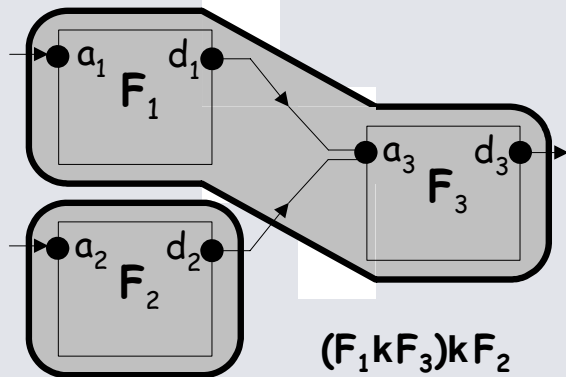
- Incremental design
  - Associative composition



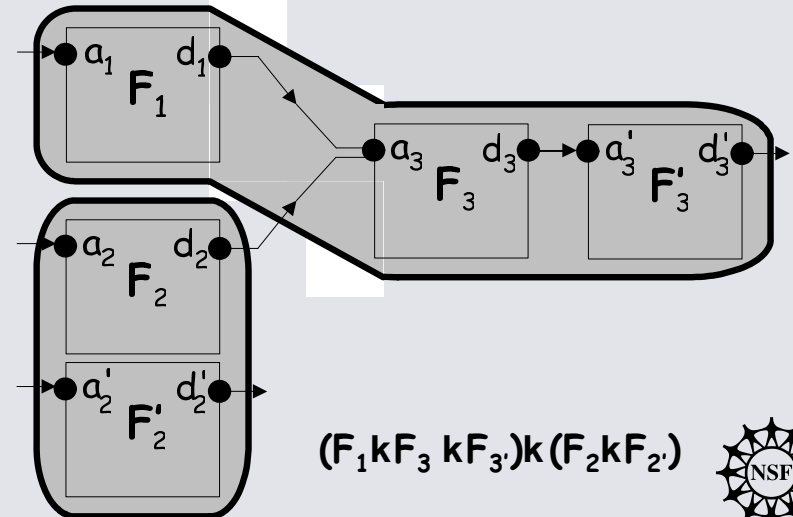
?



- Independent implementability
  - No global checks



0



(Matic and Henzinger, 2006)



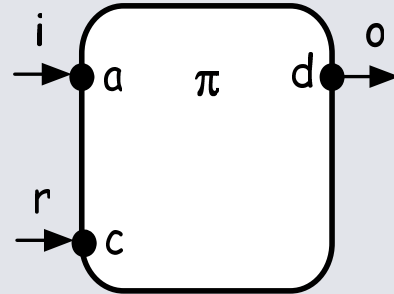
# Real-time Interface



## Assumption

requests bounded by  $a$

capacity larger than  $c$

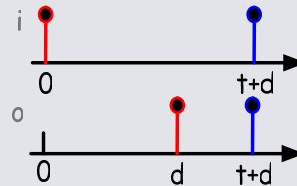


## Guarantee

output latency bounded by  $d$

Output rate function  $i^d$

$$i^d(t) = i(t+d)$$



Interface predicate

input

$$\phi^I = (r \geq c \wedge i \leq a)$$

output

$$\phi^O = o \leq i^d$$

(Matic and Henzinger, 2006)

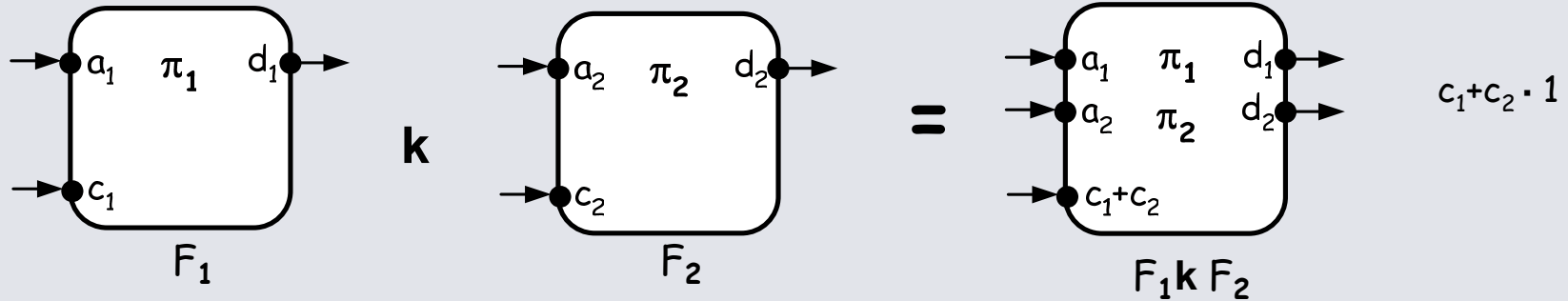




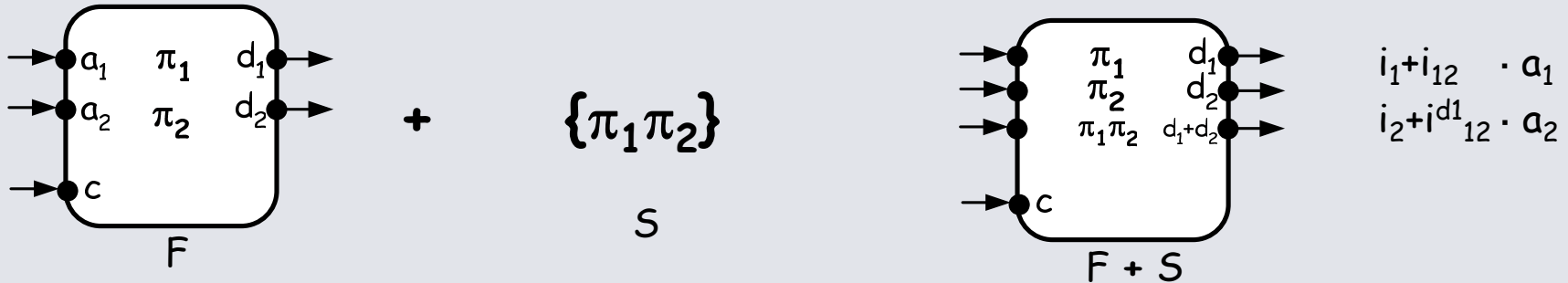
# Interface Algebra



- Composition operation  $\parallel$



- Connection operation  $+$



- Refinement relation  $\cdot$

$F'$  refines  $F$  if

- $S_F \mu S_{F'}$
- for each port valuation of  $F$  there exists a valuation of  $F'$  :

$$\phi_F^I \Rightarrow \phi_{F'}^I$$

$$\phi_{F'}^O \Rightarrow \phi_F^O$$



# Algebra Properties



➤ Incremental design

-  $(FkG)kH$  is defined

)  $Fk(GkH)$  is def.  $\text{AE } (FkG)kH = Fk(GkH)$

-  $(FkG)©S$  is defined

)  $(F©S)kG$  is def.  $\text{AE } (FkG)©S = (F©S)kG$

➤ Independent refinement

-  $FkG$  is defined  $\text{AE } F' \supseteq F$

)  $F'kG$  is def.  $\text{AE } F'kG \supseteq FkG$

-  $F©S$  is defined  $\text{AE } F' \supseteq F$

)  $F'©S$  is def.  $\text{AE } F'©S \supseteq F©S$

for all  $j=1, \dots, n$ :  $F'_j \supseteq F_j$

$E(F'_1, \dots, F'_n) \supseteq E(F_1, \dots, F_n)$

