



Functional DIF

William Plishker, Nimish Sane, Mary Kiemb, Kapil Anand, Shuvra S. Bhattacharyya University of Maryland, College Park (Dept. of Electrical and Computer Engineering)



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Outline

- Introduction
- Overview of the Dataflow Interchange Format
- Functional DIF
- Preliminary Results
- Demo







Introduction

- **Motivation**: dataflow tools can reduce the time to a functional prototype
- **Problem**: going from heterogeneous dataflow to implementation is time consuming and error prone
- Our Solution: Extend popular dataflow language with inline functional simulation semantics













DIF Package











Dataflow Interchange Format

- Describe DF graphs in text
- Simple DIF file:











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More features of DIF

• Ports









More features of DIF

• Production and consumption
 production {
 e1 = 4096;
 e10 = 1024;
 ...

```
consumption {
    e1 = 4096;
    e10 = 64;
```

- Computation keyword
- User defined attributes



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The DIF Language Syntax

```
dataflowModel graphID {
 basedon { graphID; }
 topology {
  nodes = nodelD, ...;
  edges = edgeID (srcNodeID,
snkNodeID), ...; }
 interface {
  inputs = portID [:nodeID], ...;
  outputs = portID [:nodeID], ...; }
 parameter {
  paramID [:dataType];
  paramID [:dataType] = value;
  paramID [:dataType] : range; }
 refinement {
  subgraphID = supernodeID;
  subPortID : edgeID;
  subParamID = paramID; }
```

builtInAttr {
 [elementID] = value;
 [elementID] = id;
 [elementID] = id1, id2, ...; }
attribute usrDefAttr{
 [elementID] = value;
 [elementID] = id1, id2, ...; }
actor nodeID {
 computation = stringValue;
 attrID [:attrType] [:dataType] = value;
 attrID [:attrType] [:dataType] = id;
 attrID [:attrType] [:dataType] = id1, ...; }
}









Need: Functional Sim in DIF Package:











Functional DIF

- Natural actor description
- Semantic foundation for simulating deterministic dataflow applications
- Scheduler/simulator for heterogeneous designs
- All with a focus to get to a correct implementation faster







Related Work

- Actor Description Languages
 - Xilinx's CAL, Mathwork's S-Functions
- Heterogeneous Semantic Formalisms
 - Stream Based Functions (SBF)
 - Metropolis
 - Ptolemy
 - Transactors
- Dataflow Design Environments
 - Autocoding Toolset from MCCI
 - CoCentric System Studio from Synopsis
 - Compaan from Leiden University
 - Gedae from Gedae Inc.
 - Grape from K. U. Leuven
 - LabVIEW from National Instruments
 - PeaCE from Seoul National University
 - Ptolemy II from U. C. Berkeley
 - StreamIt from MIT.











Writing actors in Functional DIF

- Divide actors into a set of *modes*
 - Each mode has a fixed consumption and production behavior
- Write the enabling conditions for each mode
- Write the computation associated with each mode
 - Including next mode to enable and then invoke
- For example, consider a standard Switch:

Switch Actor



Production & consumption behavior of switch modes

mode	consumes		produces	
	Control	Data	True	False
Control	1	0	0	0
True	0	1	1	0
False	0	1	0	1

Mode transition diagram between switch modes









Semantic Foundation

- Enable-Invoke Dataflow (EIDF)
 - Enabling Function, ε , for an actor, a:

 $\varepsilon_a:(T_a\times M_a){\,\rightarrow\,} B\,,$

- T_a , the number of input tokens on each edge
- M_a , the set of modes associated with actor a
- *B* is the Boolean set of *{true, false}*
- Invoking function, κ (Non-Deterministic)

$$\kappa_a: (I_a \times M_a) \to (O_a \times Pow(M_a)) ,$$

- I_{a} , O_{a} , input and output tokens consumed by this mode
- *Pow(M_a)*, set of valid next modes
- Core Functional Dataflow (CFDF)
 - Modify invoke to be deterministic by making one unique next mode:

 $\kappa_a^* : (I_a \times M_a) \rightarrow (O_a \times M_a)$









ACS



- DIF Graph -capture basic dataflow features (nodes, edges, tokens, etc)
- EIDF Functional, nondeterministic dataflow
- CFDF Deterministic dataflow
- Many popular forms of dataflow are **directly** supported by CFDF
 - SDF needs only one mode
 - CSDF phases correspond to modes
- Functional DIF integrates CFDF into the DIF package





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Generalized Schedule Trees

- Represents a schedule as a tree with internal nodes representing iteration counts
- Leaves represent an actor invocation
- Execution may be guarded using the enabling function of CFDF











Functional DIF Application Design Features

- Supports Heterogeneous Composition
- Unit testing to determine models of actors
- Heterogeneous Simulator
 - Simulating is as simple as walking the schedule tree

Example design using CSDF and BDF











Results: Polynomial Evaluation Accelerator (PEA)

• Polynomial Evaluation is a commonly used primitive in communication area.

$$P_i(x) = \sum_{k=0}^{n_i} C_k \times x^k$$

- The degree of P and the coefficients change in run time.
- There are four types of instructions.
 - Reset (RST), Store Polynomial (STP), Evaluate
 Polynomial (EVP), Evaluate Block (EVB)









Dataflow Modeling of PEA Testbench











The modes of PEA

Mode	behavior	Consumption		Production	
		Control	Data	Result	Status
Normal	Wait for an instruction	1	0	0	0
RST	Reset all of the coefficients	0	0	0	0
STP	Store coefficients	0	1	0	1
EVP	Evaluate the value of P	0	1	1	1
EVB	Evaluate block	0	1	1	1









Enable method pseudocode of PEA

```
Bool A.enable(x1, x2, ..., xn, mode){
   if (mode = Normal) then
          if (there is 1 token in Control buffer) then return true;
                     return false:
          else
   else if (mode = RST) then return true;
   else if (mode = STP) then
         if (there is 1 token in Data buffer) then return true;
          else
                     return false:
   else if (mode = EVP) then
          if (there is 1 token in Data buffer ) then return true;
          else
                     return false;
   else if (mode = EVB) then
         if (there is 1 token in Data buffer) then return true;
                     return false;
          else
   end if;
```









Invoke method pseudocode of PEA

```
nextMode A.invoke(x1, x2, ..., xn, mode){
   if (mode = Normal) then
          Decode instruction token
          return next mode based on instruction
   else if (mode = RST) then
          Reset all of the coefficients;
   else if (mode = STP) then
          Store coefficient for P_i(x);
          if( not done storing coefficients ) then return STP mode;
   else if (mode = EVP) then
          Evaluate P_i(x);
   else if (mode = EVB) then
          Evaluate P_i(x);
          if( not done with block) then return EVB mode;
   end if;
   return Normal mode:
```









PEA Results

Simulation times of Verilog and Functional DIF for two different sets of instructions

Instruction set	Verilog Simulation Time (ms)	Functional DIF Simulation Time (ms)	Speedup
Case 1	250	55	4.6x
Case 2	170	33	5.1x
Average	210	44	4.9x

























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Dual PEA results

Simulation times and max buffer sizes using different schedules

Application		Simulation	Max observed buffer
style	Schedule	Time (s)	size (tokens)
BDF Strict	Canonical	6.88	2,327,733
BDF Strict	Single appearance	1.72	1,729
BDF Strict	Multiple appearance	1.59	1,722
CFDF	Canonical	3.57	1,018,047
CFDF	Single appearance	0.95	1,791
CFDF	Multiple appearance	0.99	1,800















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Current status: Software Architecture

- DIF packages made up of a set of Jars
 - MoCGraph graph package for models of computation
 - MAPSS Core DIF package
 - DIF2C Software synthesis plug-in
- Unit Testing Infrastructure
- Reliance on Ptolemy
 - Typing package
 - Kernel exceptions
 - Make scripts









Summary

- Extend DIF with functional simulation by
 - Actor design considerations
 - Semantic foundation for execution
 - Supporting simulation and scheduling in the DIF package
- Simulation speeds better than Verilog
- Future Work
 - More heterogeneous applications
 - Parameterization of CFDF









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