



Fundamental Algorithms for System Modeling, Analysis, and Optimization

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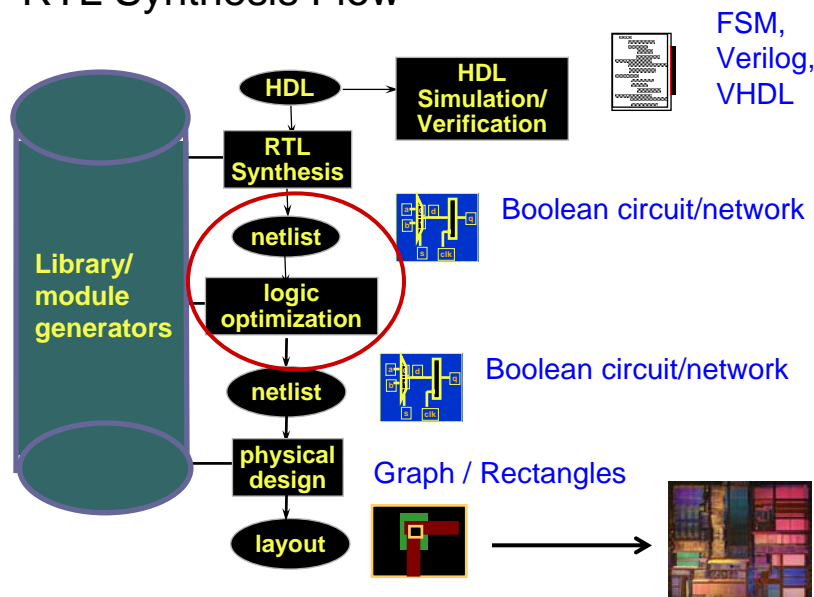
UC Berkeley
EECS 144/244
Fall 2011

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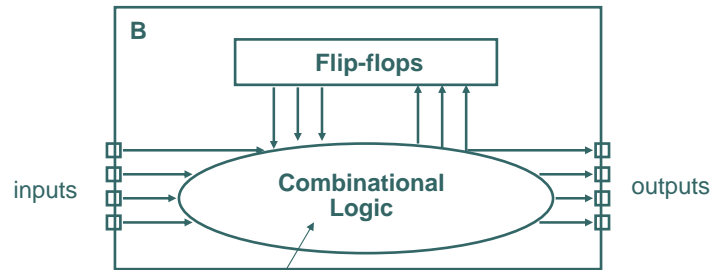
Lec 3: Boolean Algebra and Logic Optimization - 1

Thanks to S. Devadas, K. Keutzer, S. Malik, R. Rutenbar for several slides

RTL Synthesis Flow



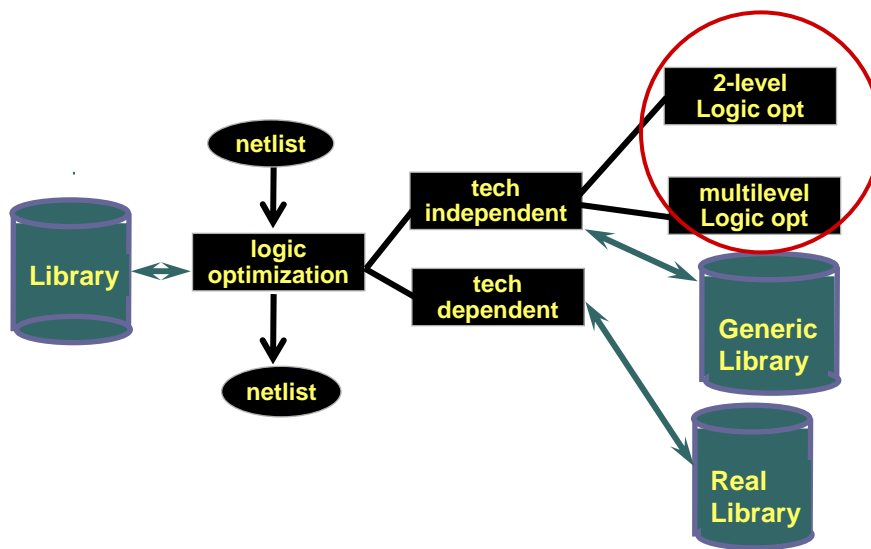
Reduce Sequential Ckt Optimization to Combinational Optimization



Optimize the size/delay/etc. of the combinational circuit (viewed as a Boolean network)

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Logic Optimization



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Outline of Topics

Basics of Boolean algebra

Two-level logic optimization

Multi-level logic optimization

Boolean function representation: BDDs

Definitions – 1: What is a Boolean function?

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Let $B = \{0, 1\}$ and $Y = \{0, 1\}$

Input variables: $X_1, X_2 \dots X_n$

Output variables: $Y_1, Y_2 \dots Y_m$

A logic function **ff** (or ‘Boolean’ function, switching function) in **n** inputs and **m** outputs is a map

$$\text{ff}: B^n \longrightarrow Y^m$$

Definition used in Logic Optimization

Let $B = \{0, 1\}$ and $Y = \{0, 1, 2\}$

Input variables: $X_1, X_2 \dots X_n$

Output variables: $Y_1, Y_2 \dots Y_m$

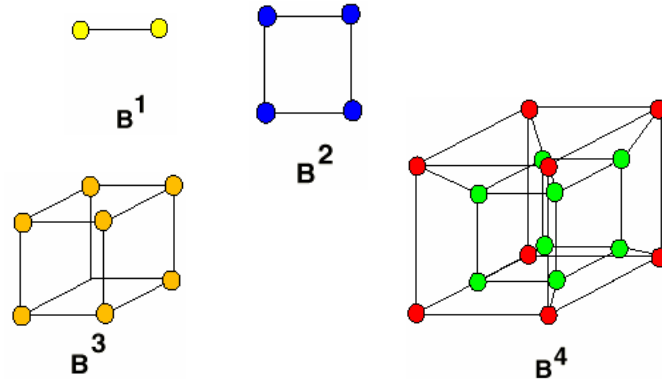
A logic function **ff** (or ‘Boolean’ function, switching function) in **n** inputs and **m** outputs is a map

$$\text{ff}: B^n \longrightarrow Y^m$$

don't care – aka “X”



The Boolean n-Cube, B^n



- $B = \{0, 1\}$
- $B^2 = \{0, 1\} \times \{0, 1\} = \{00, 01, 10, 11\}$

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Definitions – 2: ON/OFF/DC sets

If a logic function ff maps some input $b \in B^n$ to a 2 on some output i then function is incompletely specified, else completely specified

ON-SET _{i} $\subseteq B^n$, the set of all input values for which $ff_i(x) = 1$

OFF-SET _{i} $\subseteq B^n$, the set of all input values for which $ff_i(x) = 0$

DC-SET _{i} $\subseteq B^n$, the set of all input values for which $ff_i(x) = 2$

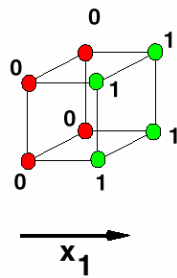
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Literals: What is a literal?

Literals

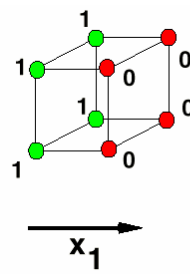
A literal is a variable or its negation y, \bar{y}

It represents a **logic function**



$$f = x_1$$

Green – ON-set
Red – OFF-set



$$f = \bar{x}_1$$

Boolean Formulas -- Syntax

Boolean functions can be **represented** by **formulas** defined as catenations of

- parentheses - (,)
- literals - $x, y, z, \bar{x}, \bar{y}, \bar{z}$
- Boolean operators - + (OR), \times (AND)
- complementation - e.g. $\overline{x + y}$

Examples:

$$f = x_1 \times \bar{x}_2 + \bar{x}_1 \times x_2$$

$$= (x_1 + x_2) \times (\bar{x}_1 + \bar{x}_2)$$

$$h = \frac{a + b \times c}{\bar{a} \times (\bar{b} + \bar{c})}$$

We will usually replace \times by catenation, e.g. $a \times b \rightarrow ab$.

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“Semantic” Description of Boolean Function

EXAMPLE: Truth table form of an incompletely specified function

$$\text{ff: } B^3 \longrightarrow Y^2$$

X_1	X_2	X_3	Y_1	Y_2
0	0	0	1	1
0	0	1	1	0
0	1	0	0	1
0	1	1	0	1
1	0	0	1	0
1	0	1	1	2
1	1	0	1	1
1	1	1	2	1

$$Y_1: \text{ON-SET}_1 = \{000, 001, 100, 101, 110\}$$

$$\text{OFF-SET}_1 = \{010, 011\}$$

$$\text{DC-SET}_1 = \{111\}$$

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Operations on Logic Functions

(1) Complement: $f \longrightarrow \bar{f}$
interchange ON and OFF-SETS

(2) Product (or intersection or logical AND)
 $h = f \cdot g$ (what happens to ON/OFF sets?)

(3) Sum (or union or logical OR):
 $h = f + g$ (ON/OFF sets?)

CNF and DNF

CNF: Conjunctive Normal Form (product of sums: POS)

DNF: Disjunctive Normal Form (sum of products: SOP)

CNF \rightarrow DNF: what is the worst-case blow up?

How about DNF \rightarrow CNF?

Cube

A cube is a **conjunction** (AND) of literals

Examples: (set of variables = {a,b,c,d})

ab

ab \bar{d}

a \bar{b} c \bar{d}

A cube is a logic function (also view as set)

2-level Minimization: Minimizing SOP (DNF)

$$F1 = \bar{A}\bar{B} + \bar{A}BD + \bar{A}B\bar{C}\bar{D} + ABC\bar{D} + A\bar{B} + ABD$$

Inputs	Outputs
00--	1
01-1	1
0100	1
1110	1
10--	1
11-1	1

$$F1 = \bar{B} + D + \bar{A}\bar{C} + AC$$

-0--	1
---1	1
0-0-	1
1-1-	1

↑
minimum representation
(number of cubes, literals)

Implicants

An *implicant* of f is a *cube* p that does not intersect the OFF-SET of f

$$p \subseteq f_{ON} \cup f_{DC}$$

Prime Implicants

An *implicant* of f is a *cube* p that does not intersect the OFF-SET of f

$$p \subseteq f_{ON} \cup f_{DC}$$

A *prime implicant* of f is an implicant p such that

- (1) No other implicant q contains it
(i.e. $p \not\subseteq q$)
- (2) $p \not\subseteq f_{DC}$

A *minterm* is a fully specified implicant
e.g., $011, 111$ (not $01-$)

Examples of Implicants/Primes

X_1	X_2	X_3	Y_1
0	0	0	1
0	0	1	1
0	1	0	0
0	1	1	0
1	0	0	1
1	0	1	1
1	1	0	1
1	1	1	2

000, 00- are implicants, but not primes (-0-)

1-1

0-0

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Prime and Irredundant Covers

A cover is a set of cubes C such that

$$C \supseteq f_{ON} \quad \text{and} \quad C \subseteq f_{ON} \cup f_{DC}$$

All of the ON-set is covered by C

C is contained in the ON-set and Don't Care Set

A *prime* cover is a cover whose cubes are all prime implicants

An *irredundant* cover is a cover C such that removing any cube from C results in a set of cubes that no longer covers the function (ON-set)

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Example Covers

X_1	X_2	X_3	Y_1	
0	0	0	1	00-
0	0	1	1	10- is a cover.
0	1	0	0	11-
0	1	1	0	
1	0	0	1	
1	0	1	1	
1	1	0	1	
1	1	1	2	

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Example Covers

X_1	X_2	X_3	Y_1	
0	0	0	1	00-
0	0	1	1	10- is a cover. Is it prime?
0	1	0	0	11- Is it irredundant?
0	1	1	0	
1	0	0	1	
1	0	1	1	
1	1	0	1	
1	1	1	2	

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Minimum covers

Defn: A *minimum cover* is a cover of minimum cardinality

Theorem: There exists a minimum cover that is a prime and irredundant cover.

Why?

Minimum covers

Defn: A *minimum cover* is a cover of minimum cardinality

Theorem: There exists a minimum cover that is a prime and irredundant cover.

Given any cover C

- (a) if redundant, not minimum
- (b) if any cube q is not prime, replace q with prime $p \supseteq q$ and continue until all cubes prime; it is a minimum prime cover

Example Covers

X_1	X_2	X_3	Y_1
0	0	0	1
0	0	1	1
0	1	0	0
0	1	1	0
1	0	0	1
1	0	1	1
1	1	0	1
1	1	1	2

00-
 10- is a cover. Is it prime?
 11- Is it irredundant?

What is a minimum prime and
 irredundant cover for the function?

Example Covers

X_1	X_2	X_3	Y_1
0	0	0	1
0	0	1	1
0	1	0	0
0	1	1	0
1	0	0	1
1	0	1	1
1	1	0	1
1	1	1	2

00-
 10- is a cover. Is it prime?
 11- Is it irredundant?

-0-
 11- is a cover. Is it prime?
 Is it irredundant?
 Is it minimum?

Example Covers

X_1	X_2	X_3	Y_1	
0	0	0	1	00-
0	0	1	1	10- is a cover. Is it prime?
0	1	0	0	11- Is it irredundant?
0	1	1	0	
1	0	0	1	-0-
1	0	1	1	11- is a cover. Is it prime?
1	1	0	1	Is it irredundant?
1	1	1	2	Is it minimum?

What about

-0-

1--

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The Quine-McCluskey Method: Exact Minimization

Step 1: List all minterms in ON-SET and DC-SET

Step 2: Use a prescribed sequence of steps to find all the prime implicants of the function

Step 3: Construct the prime implicant table

Step 4: Find a minimum set of prime implicants that cover all the minterms

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Espresso Algorithm: Heuristic Minimization

```
ESPRESSO (F, DC)  {  
  F is ON-SET, DC is Don't Care Set  
  1. R = U - (F ∪ DC)      U is universe cube  
  2. n = |F|  
  3. F = Reduce (F, DC); // reduce implicants in F  
    to non-prime cubes  
  4. F = Expand (F, R); // expand cubes to prime  
    implicants  
  5. F = Irredundant (F, DC); // extract minimal  
    cover of prime implicants  
  6. If |F| < n goto 2, else, post-process & exit  
}
```

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Multi-level Logic Optimization

2-level optimization is a 'solved' problem:
Espresso is considered the last word on the topic

But most circuits are not two-level!

Need techniques to optimize size of multi-level circuits

- Size measured in terms of number of literals, depth of the circuit, etc.

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