



# Fundamental Algorithms for System Modeling, Analysis, and Optimization



Stavros Tripakis

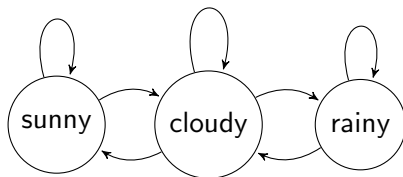
UC Berkeley  
EECS 144/244  
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Stochastic Systems

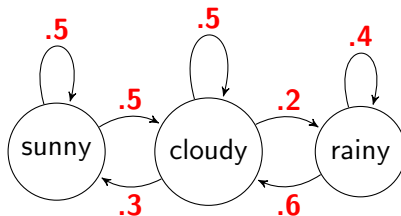
# Probabilistic systems

From non-deterministic systems:



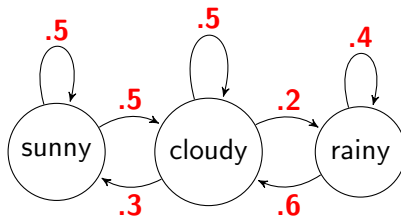
# Probabilistic systems

To probabilistic systems:



# Markov chains

A finite Markov chain:



$$P(\text{next state is "sunny"} \mid \text{current state is "sunny"}) = 0.5$$

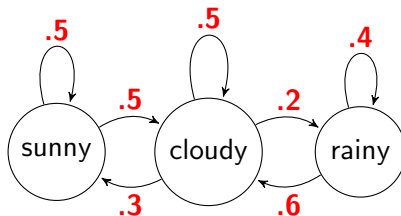
$$P(\text{next state is "sunny"} \mid \text{current state is "cloudy"}) = 0.3$$

$$P(\text{next state is "sunny"} \mid \text{current state is "rainy"}) = 0$$

...

# Markov chains

A finite Markov chain:



$$P(\text{next state is "sunny"} \mid \text{current state is "sunny"}) = 0.5$$

$$P(\text{next state is "sunny"} \mid \text{current state is "cloudy"}) = 0.3$$

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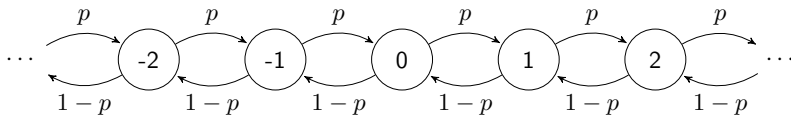
...

The *Markov property*: only current state matters:

$$P(s_{k+1} = v_{k+1} \mid s_k = v_k, s_{k-1} = v_{k-1}, \dots, s_0 = v_0) = P(s_{k+1} = v_{k+1} \mid s_k = v_k)$$

# Markov chains

Markov chains can be infinite:



# Markov chains

A finite Markov chain with  $n$  states can be represented by a square  $n \times n$  *probability matrix*  $\mathbf{P}$ :

$$\mathbf{P} = \begin{bmatrix} p_{11} & p_{12} & \cdots & p_{1n} \\ p_{21} & p_{22} & \cdots & p_{2n} \\ \vdots & & \ddots & \vdots \\ p_{n1} & p_{n2} & \cdots & p_{nn} \end{bmatrix}$$

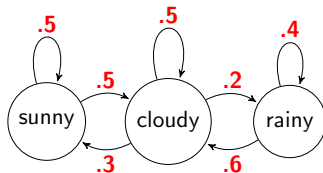
where

$$p_{ij} = P(\text{next state is } j \mid \text{current state is } i)$$

To be a valid probability matrix,  $\mathbf{P}$  must satisfy:

$$\forall i, j : p_{ij} \geq 0 \quad \text{and} \quad \forall i : \sum_{j=1}^n p_{ij} = 1$$

# Markov chains



$$\mathbf{P} = \begin{bmatrix} 0.5 & 0.5 & 0 \\ 0.3 & 0.5 & 0.2 \\ 0 & 0.6 & 0.4 \end{bmatrix}$$

# Transforming a process into a Markov chain

## Quiz:

Suppose that whether or not it rains today depends on the previous weather conditions during the last two days. Specifically:

- ▶ If it has rained for the past two days, then it will rain tomorrow with probability 0.7.
- ▶ If it rained today but not yesterday, then it will rain tomorrow with probability 0.5.
- ▶ If it rained yesterday but not today, then it will rain tomorrow with probability 0.4.
- ▶ If it has not rained in the past two days, then it will rain tomorrow with probability 0.2.

Is this process Markovian? If so build a Markov chain that models the process.

# Transforming a process into a Markov chain

## Quiz:

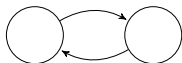
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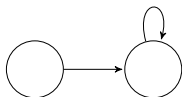
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# Discrete systems vs. Markov chains

Some discrete systems are Markov chains ...



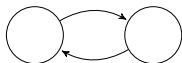
$$\mathbf{P} = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}$$



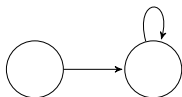
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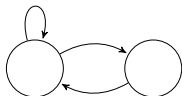


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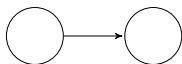


$$\mathbf{P} = \begin{bmatrix} 0 & 1 \\ 0 & 1 \end{bmatrix}$$

... but not all:



$$\mathbf{P} = \begin{bmatrix} ? & ? \\ 1 & 0 \end{bmatrix}$$

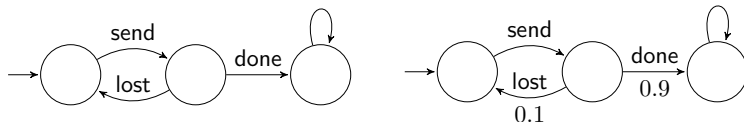


$$\mathbf{P} = \begin{bmatrix} 0 & 1 \\ \mathbf{0} & \mathbf{0} \end{bmatrix}$$

# Discrete systems vs. Markov chains

In the other direction, Markov chains are extensions of discrete systems:

MCs contain more information (next-state probabilities).



# COMPOSITION OF MARKOV CHAINS

—

# MARKOV DECISION PROCESSES

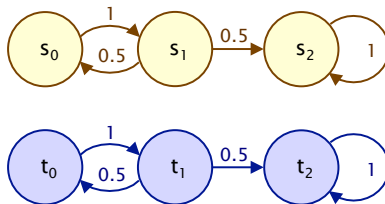
# Composition of Markov chains

Suppose we want to compose the following two Markov chains:

PRISM code:

```
module M1
  s : [0..2] init 0;
  [] s=0 -> (s'=1);
  [] s=1 -> 0.5:(s'=0) + 0.5:(s'=2);
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endmodule
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module M2 = M1 [ s=t ] endmodule
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Several figures due to Dave Parker.

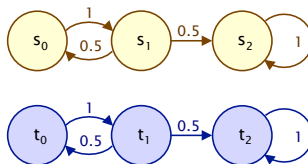
# Synchronous composition of Markov chains

What does the **synchronous** composition of these processes look like?

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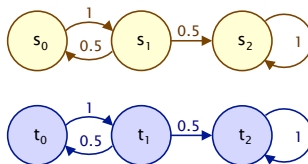
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$$P((s_1, t_1) \mid (s_0, t_0)) =$$

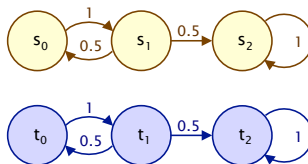
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$$P((s_1, t_1) \mid (s_0, t_0)) = 1 \cdot 1 = 1$$

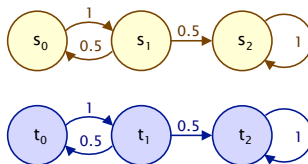
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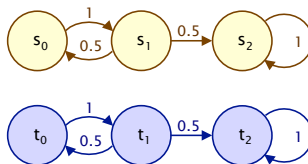
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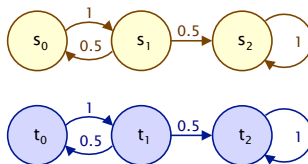
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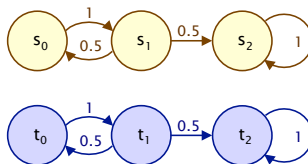
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Is the synchronous composition of two Markov chains a Markov chain?

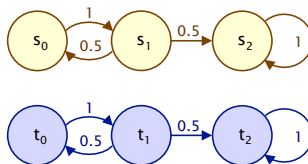
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Is the synchronous composition of two Markov chains a Markov chain?  
Yes!

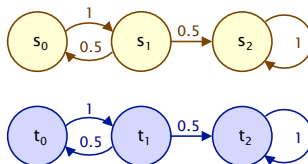
# Asynchronous composition of Markov chains

What would the **asynchronous** composition of these two processes look like?

PRISM code:

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  [] s=1 -> 0.5:(s'=0) + 0.5:(s'=2);
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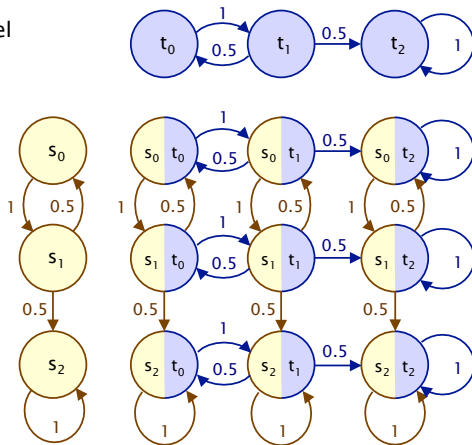
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# Asynchronous composition of Markov chains

**Asynchronous** parallel  
composition of two  
3-state DTMCs

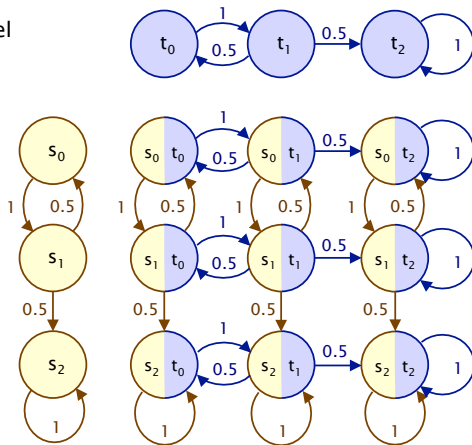
Action labels  
omitted here



# Asynchronous composition of Markov chains

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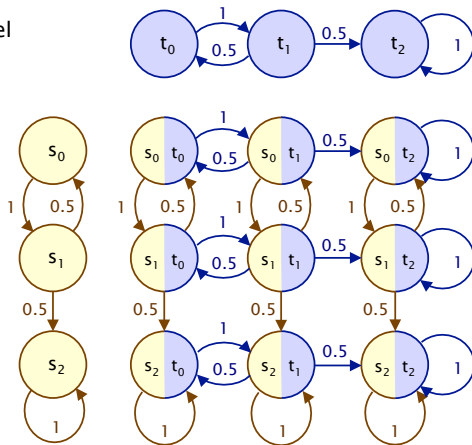


Is the asynchronous composition of two Markov chains a Markov chain?

# Asynchronous composition of Markov chains

**Asynchronous** parallel  
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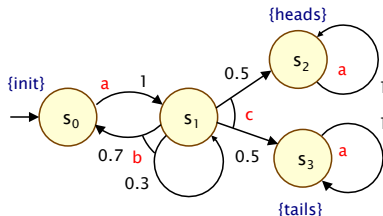
Action labels  
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Is the asynchronous composition of two Markov chains a Markov chain? No! It is a **Markov Decision Process**.

# Markov Decision Processes (MDPs)

Combine non-deterministic and probabilistic choice.

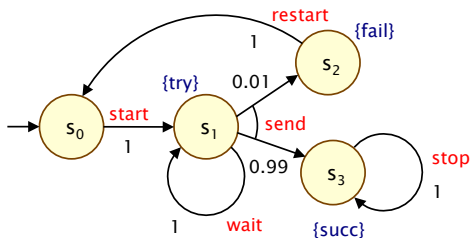


Intuitive semantics:

- ▶ First choose action non-deterministically among possible actions.
  - ▶ In state  $s_0$ , only one possible action,  $a$ .
  - ▶ In state  $s_1$ , two possible actions,  $b$  and  $c$ .
- ▶ Then, given chosen action, throw a dice and pick successor state w.r.t. the specified probability distribution for that action.

# Markov Decision Processes (MDPs)

Non-determinism has multiple uses, as in discrete systems.  
E.g., useful to model abstraction:



At state  $s_1$ , if channel is ready then attempt to send, otherwise wait.

Details of when channel is ready are not modeled.

# Model-checking MDPs

Tools such as PRISM answer queries like:

- **Byzantine agreement protocol**
  - $P_{\min=?} [ F (\text{agreement} \wedge \text{rounds} \leq 2) ]$
  - “what is the minimum probability that agreement is reached within two rounds?”
- **CSMA/CD communication protocol**
  - $P_{\max=?} [ F \text{ collisions} = k ]$
  - “what is the maximum probability of k collisions?”
- **Self-stabilisation protocols**
  - $P_{\min=?} [ F^{\leq t} \text{ stable} ]$
  - “what is the minimum probability of reaching a stable state within k steps?”

See PRISM web site and literature for details:

<http://www.prismmodelchecker.org/>

# ANALYSIS OF MARKOV CHAINS

Interesting questions:

- ▶ After  $k$  steps, what is the likelihood that the system is at state  $i$ ?
- ▶ In the long run, how much time does the system spend at state  $i$ ? (i.e., how often is  $i$  visited?)
- ▶ What is the probability that the system will ever reach a given state (or group of states)?
- ▶ What is the expected time until the system reaches a given state (or group of states)?

# Computing state probability vectors

Let  $\mathbf{x} = [p_1 \ p_2 \ \cdots \ p_n]$  be a *state probability vector*, where

$$p_i = P(\text{current state is } i)$$

Of course we must have:  $\sum_{i=1}^n p_i = 1$ .

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Let  $\mathbf{x}' = [p'_1 \ p'_2 \ \cdots \ p'_n]$  be the **next state** probability vector.

Then, for  $i = 1, \dots, n$ :

$$p'_i = p_1 \cdot p_{1i} + p_2 \cdot p_{2i} + \cdots + p_n \cdot p_{ni}$$

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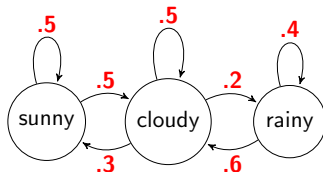
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So:

$$\mathbf{x}' = \mathbf{x} \cdot \mathbf{P}$$

# Computing state probability vectors

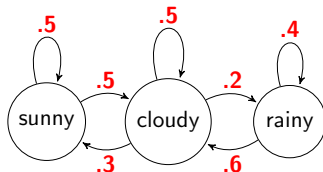
Example:



$$\begin{bmatrix} 1 & 0 & 0 \end{bmatrix} \cdot \begin{bmatrix} 0.5 & 0.5 & 0 \\ 0.3 & 0.5 & 0.2 \\ 0 & 0.6 & 0.4 \end{bmatrix} = \begin{bmatrix} 0.5 & 0.5 & 0 \end{bmatrix}$$

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# Computing state probability vectors

The probabilities for the **next state** are given by

$$\mathbf{x}' = \mathbf{x} \cdot \mathbf{P}$$

In general:

$$\mathbf{x}_{k+1} = \mathbf{x}_k \cdot \mathbf{P} = (\mathbf{x}_{k-1} \cdot \mathbf{P}) \cdot \mathbf{P} = \dots = \mathbf{x}_0 \cdot \underbrace{\mathbf{P} \cdot \mathbf{P} \dots \mathbf{P}}_{k+1 \text{ times}}$$

$$\mathbf{x}_{k+1} = \mathbf{x}_0 \cdot \mathbf{P}^{k+1}$$

- After  $k$  steps, what is the likelihood that the system is at state  $i$ ?

$$\mathbf{x}_k = \mathbf{x}_0 \cdot \mathbf{P}^k$$

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# Markov chains and graphs

A Markov chain has a graph structure.

We can partly answer the questions simply by studying this structure, **completely ignoring the probability numbers**.

## $n$ -step transition probabilities

Let  $\mathbf{P}_{ij}^n$  be the  $(i, j)$  element of  $\mathbf{P}^n$ .

What does  $\mathbf{P}_{ij}^n$  represent?

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What does  $\mathbf{P}_{ij}^n$  represent?

$\mathbf{P}_{ij}^n = P(s_{k+n} = j \mid s_k = i)$ : probability that, starting from state  $i$ , after  $n$  steps the state will be  $j$ .

# Classification of states in Markov chains

State  $i$  is *absorbing* if  $\mathbf{P}_{ii} = 1$ . This implies  $\mathbf{P}_{ij} = 0$  for all  $j \neq i$ .

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Two states  $i$  and  $j$  *communicate*, written  $i \leftrightarrow j$ , if  $i$  is accessible from  $j$  and vice-versa.

A set of states that communicate is called a *class*.

# Classification of states in Markov chains

State  $i$  is *absorbing* if  $\mathbf{P}_{ii} = 1$ . This implies  $\mathbf{P}_{ij} = 0$  for all  $j \neq i$ .  
This says that  $i$  only has a self-loop transition.

State  $j$  is *accessible* from state  $i$  if  $\mathbf{P}_{ij}^n > 0$  for some  $n$ .  
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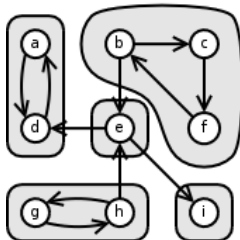
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A set of states that communicate is called a *class*.  
A class is a strongly-connected component.

# Strongly-connected components

In a directed graph  $G = (V, \rightarrow)$ , a *strongly-connected component* (SCC) is a subset of nodes  $C \subseteq V$ , such that every node in  $C$  is reachable from every other node in  $C$ .

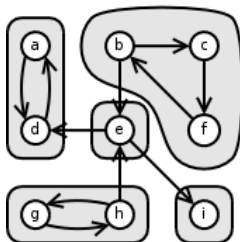
$C$  is called *maximal* if we cannot add more nodes to  $C$  and still preserve its SCC property, i.e.,  $\nexists C' \supset C$  s.t.  $C'$  is also a SCC.



# The acyclic graph of maximal SCCs

The set of all maximal SCCs of a graph defines a new graph, where nodes are maximal SCCs,  $C_1, C_2, \dots, C_m$ .

An edge  $C_i \rightsquigarrow C_j$  exists iff  $C_i \neq C_j$  and there is a node in  $C_i$  that has a successor node in  $C_j$ .



This graph of SCCs is by definition acyclic: why?

# Irreducible Markov chains

The Markov chain is *irreducible* if it has only one class, i.e., all states communicate with each other.

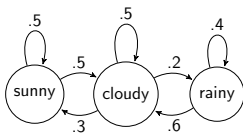
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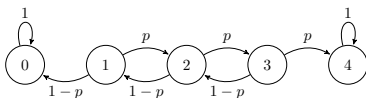
This says that the whole Markov chain is a SCC.

# Examples

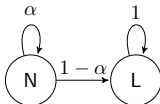
Weather model (irreducible):



Gambling model (reducible):

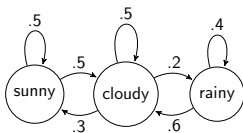


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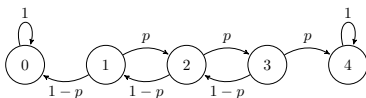


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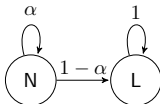
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Learning model (reducible):



States 0, 4, and L are absorbing states.

# Recurrent and transient states

Let  $i, j$  be two states. Define:

$f_{ij}^n$  = probability that, starting in  $i$ , the first transition into  $j$  happens after  $n$  steps

$$f_{ij} = \sum_{n=1}^{\infty} f_{ij}^n = \text{probability of reaching } j \text{ from } i \text{ in any \# steps}$$

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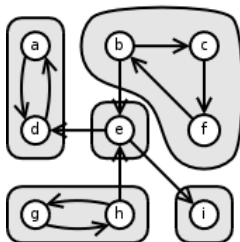
State  $i$  is:

- ▶ *Recurrent* if  $f_{ii} = 1$ .
- ▶ *Transient* if  $f_{ii} < 1$ .

# Structural characterization of transient and recurrent states

A SCC  $C$  is called *terminal* if there is no  $C'$  such that  $C \rightsquigarrow C'$ .

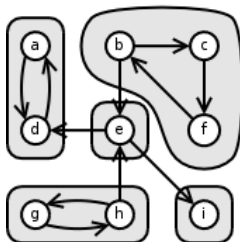
Otherwise  $C$  is called *transient*.



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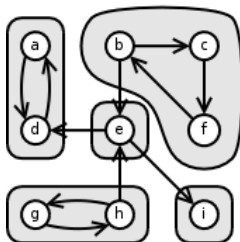
Terminal SCCs:  $\{a, d\}$  and  $\{i\}$ .

Transient SCCs:  $\{b, c, f\}$ ,  $\{e\}$ , and  $\{g, h\}$ .

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Terminal SCCs:  $\{a, d\}$  and  $\{i\}$ .

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- ▶ Recurrent states = states belonging to terminal SCCs.
- ▶ Transient states = states belonging to transient SCCs.

# Structural characterization of transient and recurrent states

If  $i$  is recurrent and  $i \leftrightarrow j$  then  $j$  is also recurrent.

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Because  $M$  is finite and deadlock-free.

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In a finite Markov chain  $M$ , some states will be recurrent. Why?  
Because  $M$  is finite and deadlock-free.

Does this hold also for infinite Markov chains?

Viewed as a graph, every finite Markov chain  $M$  has at least one terminal maximal SCC.

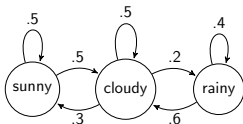
## Theorem

*In the long run the amount of time that  $M$  spends in transient SCCs is 0.*

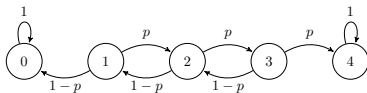
Therefore, the probability that after some time  $M$  will reach a terminal SCC and remain forever there is 1.

# Examples

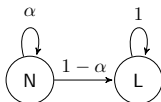
Weather model: all states visited infinitely often.



Gambling model: eventually system enters either 0 or 4 and then stays there forever.

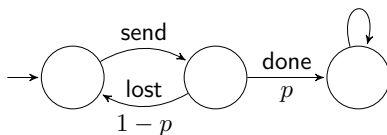


Learning model: eventually system enters L and never leaves.



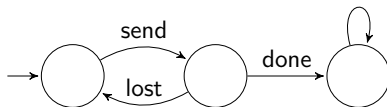
# Probabilities vs. Nondeterminism and Fairness

In a **probabilistic** system, the behavior where the message keeps getting lost after being sent has probability 0, **independently of the value of  $p$**  (provided  $p > 0$ ):



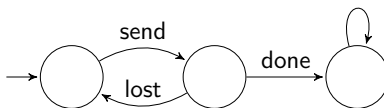
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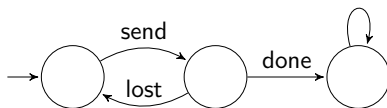


However, we can add fairness constraints to ensure that it does not, e.g.,:

$$\underbrace{((\Box \Diamond \text{send}) \rightarrow (\Diamond \text{done}))}_{\text{fairness constraint}} \rightarrow (\text{what we want})$$

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If this is all we need, probabilities are an overkill.

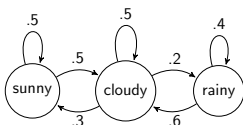
# Analysis of Markov chains

Interesting questions:

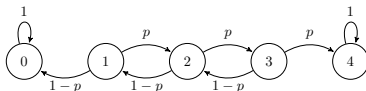
- ✓ After  $k$  steps, what is the likelihood that the system is at state  $i$ ?
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- ✓ (partly) What is the probability that the system will ever reach a given state (or group of states)?
- ▶ What is the expected time until the system reaches a given state (or group of states)?

# Examples

Weather model: all states visited infinitely often.

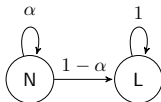


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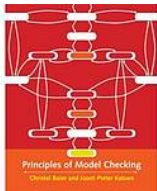
When does the gambler go bankrupt?

Learning model: eventually system enters L and never leaves.



# REACHABILITY PROBABILITIES IN MARKOV CHAINS

Material taken mainly from [Baier and Katoen(2008)], Chapter 10.



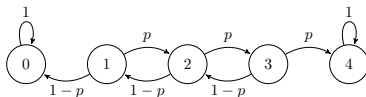
# Reachability question for Markov chains

Let  $M$  be a Markov chain and  $B$  a set of states of  $M$ .

**Reachability question for Markov chains:** what is the probability of reaching  $B$ ?

Note: this is not a yes/no question, as in standard model-checking. Here, we want to compute a probability  $p \in [0, 1]$ .

## Example: gambler's ruin



What is the probability that the gambler wins?

Model in PRISM:

```
2 dtmc // this model is a Markov chain
3
4 const double p = 0.7;
5
6 module M
7     s : [0..4] init 2;
8
9     [] s=0 -> (s'=0);
10    [] s=1 -> p : (s'=2) + (1-p) : (s'=0);
11    [] s=2 -> p : (s'=3) + (1-p) : (s'=1);
12    [] s=3 -> p : (s'=4) + (1-p) : (s'=2);
13    [] s=4 -> (s'=4);
14 endmodule
```

PCTL formula in PRISM:

$P=? [ F s=4 ]$

PRISM answers:

$p$	answer
0	0
0.5	0.499999...
0.7	0.844827...
1	1

# Computing reachability probabilities

Let  $x_i$  be the probability that the target set  $B$  is reached starting from state  $i$ .

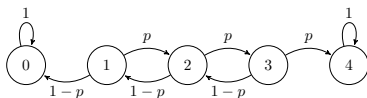
Then:

- ▶ If  $i \in B$  then  $x_i = 1$ .
- ▶ If  $i$  cannot reach  $B$  in the graph sense, then  $x_i = 0$ .
- ▶ Otherwise

$$x_i = \sum_j p_{ij} \cdot x_j$$

This forms a set of linear equations. For finite chains it is finite and is guaranteed to have a unique solution.

## Example: gambler's ruin



What is the probability that the gambler wins?

$$x_0 = 0$$

$$x_4 = 1$$

$$x_1 = p \cdot x_2 + (1 - p) \cdot x_0 = p \cdot x_2$$

$$x_2 = p \cdot x_3 + (1 - p) \cdot x_1$$

$$x_3 = p \cdot x_4 + (1 - p) \cdot x_2 = p + (1 - p) \cdot x_2$$

For  $p = \frac{1}{2}$ ,  $x_2 = \frac{1}{2}$ . For  $p = 0.7$ ,  $x_2 = 0.8448$ .

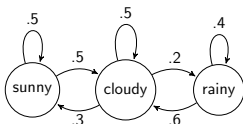
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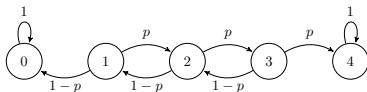
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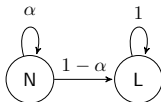


How much time does it rain on the average?

Gambling model: eventually system enters either 0 or 4 and then stays there forever.



Learning model: eventually system enters L and never leaves.



# STATIONARY DISTRIBUTION

## Periodic states

State  $i$  has period  $d \in \mathbb{N}$  if  $\mathbf{P}_{ii}^n = 0$  whenever  $n$  is not divisible by  $d$ , and  $d$  is the largest such number.

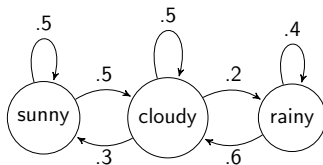
If  $d = 1$  (i.e.,  $\mathbf{P}_{ii}^n > 0$  for all  $n$ ) then state  $i$  is called *aperiodic*.

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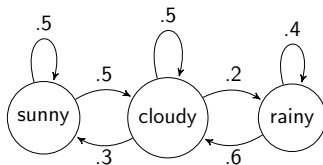


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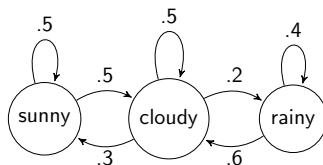
All states are aperiodic.

# Periodic states

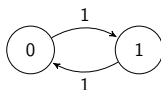
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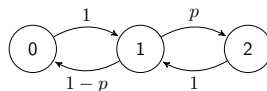
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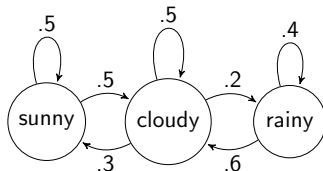
Both states have period 2.



All states have period 2.

# Ergodic states

An aperiodic recurrent state is called *ergodic*.



All states are ergodic.

# Stationary distribution

## Theorem

*Let  $M$  be a finite, irreducible Markov chain where all states of  $M$  are ergodic. Then the limit*

$$\lim_{k \rightarrow \infty} \mathbf{P}_{ij}^k$$

*exists and is independent of  $i$  (i.e.,  $\forall i, i' : \lim_{k \rightarrow \infty} \mathbf{P}_{ij}^k = \lim_{k \rightarrow \infty} \mathbf{P}_{i'j}^k$ ).*

*Furthermore, letting*

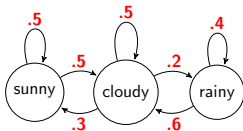
$$\pi_j = \lim_{k \rightarrow \infty} \mathbf{P}_{ij}^k$$

*then  $\pi = [\pi_1 \ \pi_2 \ \cdots \ \pi_n]$  is the unique non-negative solution of*

$$\pi = \pi \cdot \mathbf{P}$$

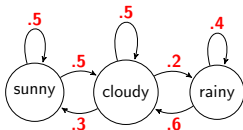
$\pi$  is called the *stationary distribution*.

## Stationary distribution: example



$$\mathbf{P} = \begin{bmatrix} 0.5 & 0.5 & 0 \\ 0.3 & 0.5 & 0.2 \\ 0 & 0.6 & 0.4 \end{bmatrix}$$

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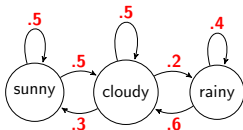


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$$\forall k > 15 : \mathbf{P}^k = \begin{bmatrix} 0.3103 & 0.5172 & 0.1724 \\ 0.3103 & 0.5172 & 0.1724 \\ 0.3103 & 0.5172 & 0.1724 \end{bmatrix}$$

What does this imply?

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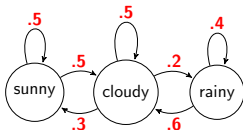


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for any probability vector  $\mathbf{x}_0$ ,  $\forall k > 15 : \mathbf{x}_0 \cdot \mathbf{P}^k = \pi = [0.3103 \ 0.5172 \ 0.1724]$

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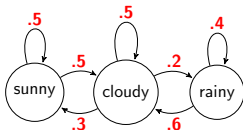
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So how much time does it rain on the average?

If the chain is not ergodic, the limit may not exist, e.g.,

$$\mathbf{P} = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}$$

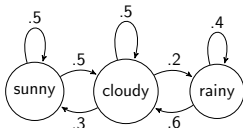
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Interesting questions:

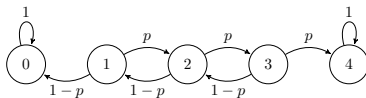
- ✓ After  $k$  steps, what is the likelihood that the system is at state  $i$ ?
- ✓ In the long run, how much time does the system spend at state  $i$ ? (i.e., how often is  $i$  visited?)
- ✓ What is the probability that the system will ever reach a given state (or group of states)?
- ▶ What is the expected time until the system reaches a given state (or group of states)?

# Examples

Weather model:

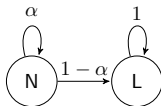


Gambling model:



How much time does an average game last?

Learning model:



How long until we learn something?

# Transient analysis

Order the states of a Markov chain  $M$  so that  $\{1, 2, \dots, t\}$  is the set of transient states.

Let

$$\mathbf{P}_T = \begin{bmatrix} p_{11} & p_{12} & \cdots & p_{1t} \\ p_{21} & p_{22} & \cdots & p_{2t} \\ \vdots & & \ddots & \vdots \\ p_{t1} & p_{t2} & \cdots & p_{tt} \end{bmatrix}$$

Observation: some rows of  $\mathbf{P}_T$  sum to  $< 1$  (otherwise this would be a SCC).

# Transient analysis

Let

$q_{ij}$  = mean time spent in  $j$ , given that the system starts in  $i$

Then

$$q_{ij} = \delta_{ij} + \sum_k p_{ik} \cdot q_{kj}$$

where

$$\delta_{ij} = \begin{cases} 1 & \text{if } i = j \\ 0 & \text{otherwise.} \end{cases}$$

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Because  $q_{kj} = 0$  when  $k$  is recurrent (cannot move from recurrent state to transient state).

where

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# Transient analysis

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$$\mathbf{Q} = \begin{bmatrix} q_{11} & q_{12} & \cdots & q_{1t} \\ q_{21} & q_{22} & \cdots & q_{2t} \\ \vdots & & \ddots & \vdots \\ q_{t1} & q_{t2} & \cdots & q_{tt} \end{bmatrix}$$

Then

$$\mathbf{Q} = \mathbf{I} + \mathbf{P}_T \cdot \mathbf{Q}$$

where  $\mathbf{I}$  is the identity matrix of size  $t$ .

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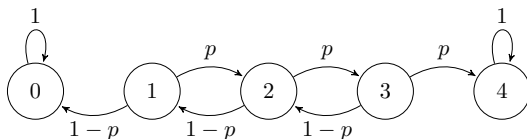
$$\mathbf{Q} = \mathbf{I} + \mathbf{P}_T \cdot \mathbf{Q}$$

where  $\mathbf{I}$  is the identity matrix of size  $t$ .

It can be shown that  $\mathbf{I} - \mathbf{P}_T$  is invertible. Therefore:

$$\mathbf{Q} = (\mathbf{I} - \mathbf{P}_T)^{-1}$$

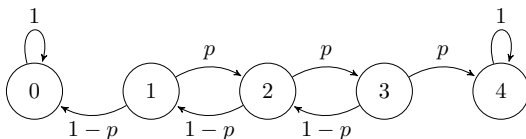
## Example: gambler's ruin



$$\mathbf{Q} = (\mathbf{I} - \mathbf{P}_T)^{-1} = \begin{bmatrix} 1 & -p & 0 \\ p-1 & 1 & -p \\ 0 & p-1 & 1 \end{bmatrix}^{-1} = \dots$$

$$\mathbf{Q}_{p=0} = \begin{bmatrix} 1 & 0 & 0 \\ 1 & 1 & 0 \\ 1 & 1 & 1 \end{bmatrix} \quad \mathbf{Q}_{p=1} = \begin{bmatrix} 1 & 1 & 1 \\ 0 & 1 & 1 \\ 0 & 0 & 1 \end{bmatrix} \quad \mathbf{Q}_{p=\frac{1}{2}} = \begin{bmatrix} 1.5 & 1 & 0.5 \\ 1 & 2 & 1 \\ 0.5 & 1 & 1.5 \end{bmatrix}$$

## Example: gambler's ruin

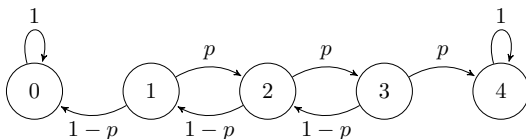


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What is the average playing time with  $p = \frac{1}{2}$ ?  
3 if I start with \$1 or \$3. 4 if I start with \$2.

# Analysis of Markov chains

Interesting questions:

- ✓ After  $k$  steps, what is the likelihood that the system is at state  $i$ ?
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<http://www.prismmodelchecker.org/lectures/>.



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