EECS 144/244: Fundamental Algorithms for System Modeling, Analysis, and Optimization

Probabilistic Systems

Lecture: Markov Chains and Markov Decision Processes

Stavros Tripakis

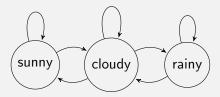
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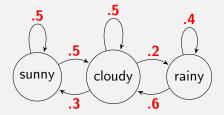
Probabilistic systems

From non-deterministic systems:

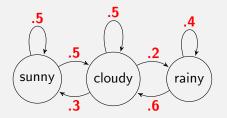


Probabilistic systems

To probabilistic systems:

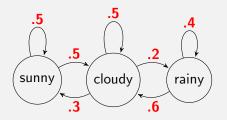


A finite Markov chain:



```
P({\sf next\ state\ is\ "sunny"}\ |\ {\sf current\ state\ is\ "sunny"}) \quad = \quad 0.5
P(\text{next state is "sunny"} \mid \text{current state is "cloudy"}) = 0.3
P(\text{next state is "sunny"} \mid \text{current state is "rainy"}) = 0
```

A finite Markov chain:

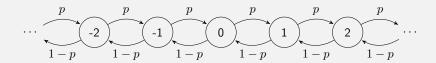


$$\begin{array}{lll} P(\mbox{next state is "sunny"} \mid \mbox{current state is "sunny"}) &=& 0.5 \\ P(\mbox{next state is "sunny"} \mid \mbox{current state is "cloudy"}) &=& 0.3 \\ P(\mbox{next state is "sunny"} \mid \mbox{current state is "rainy"}) &=& 0 \\ \dots \end{array}$$

The Markov property: only current state matters:

$$P(s_{k+1} = v_{k+1} \mid s_k = v_k, s_{k-1} = v_{k-1}, ..., s_0 = v_0) = P(s_{k+1} = v_{k+1} \mid s_k = v_k)$$

An infinite Markov chain:



A finite Markov chain with n states can be represented by a square $n \times n$ probability matrix \mathbf{P} :

$$\mathbf{P} = \begin{bmatrix} p_{11} & p_{12} & \cdots & p_{1n} \\ p_{21} & p_{22} & \cdots & p_{2n} \\ \vdots & & \ddots & \vdots \\ p_{n1} & p_{n2} & \cdots & p_{nn} \end{bmatrix}$$

where

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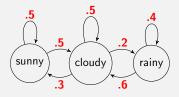
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To be a valid probability matrix, \mathbf{P} must satisfy:

$$\forall i, j : p_{ij} \ge 0$$
 and $\forall i : \sum_{j=1}^{n} p_{ij} = 1$



$$\mathbf{P} = \left[\begin{array}{ccc} 0.5 & 0.5 & 0 \\ 0.3 & 0.5 & 0.2 \\ 0 & 0.6 & 0.4 \end{array} \right]$$

Transforming a process into a Markov chain

Homework:

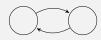
Suppose that whether or not it rains today depends on the previous weather conditions during the last two days. Specifically:

- ▶ If it has rained for the past two days, then it will rain tomorrow with probability 0.7.
- ▶ If it rained today but not yesterday, then it will rain tomorrow with probability 0.5.
- ▶ If it rained yesterday but not today, then it will rain tomorrow with probability 0.4.
- ▶ If it has not rained in the past two days, then it will rain tomorrow with probability 0.2.

Is this process Markovian? If so build a Markov chain that models the process.

Discrete systems vs. Markov chains

Some discrete systems are Markov chains ...

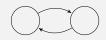


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... but not all:



$$\mathbf{P} = \left[\begin{array}{cc} ? & ? \\ 1 & 0 \end{array} \right]$$

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Discrete systems vs. Markov chains

In the other direction, Markov chains are extensions of discrete systems:

MCs contain more information (next-state probabilities).

ANALYSIS OF MARKOV CHAINS

Interesting questions:

- ► After *k* steps, what is the likelihood that the system is at state *i*?
- ▶ In the long run, how much time does the system spend at state *i*? (i.e., how often is *i* visited?)
- ► What is the probability that the system will ever reach a given state (or group of states)?
- ► What is the expected time until the system reaches a given state (or group of states)?

Let
$$\mathbf{x} = [p_1 \; p_2 \; \cdots \; p_n]$$
 be a state probability vector, where
$$p_i = P(\text{current state is } i)$$

Of course we must have: $\sum_{i=1}^{n} p_i = 1$.

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Let $\mathbf{x}' = [p_1' \ p_2' \ \cdots \ p_n']$ be the **next state** probability vector. Then, for i = 1, ..., n:

$$p_i' = p_1 \cdot p_{1i} + p_2 \cdot p_{2i} + \dots + p_n \cdot p_{ni}$$

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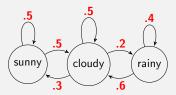
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So:

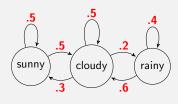
$$\mathbf{x}' = \mathbf{x} \cdot \mathbf{P}$$

Example:



$$\begin{bmatrix} 1 & 0 & 0 \end{bmatrix} \cdot \begin{bmatrix} 0.5 & 0.5 & 0 \\ 0.3 & 0.5 & 0.2 \\ 0 & 0.6 & 0.4 \end{bmatrix} = \begin{bmatrix} 0.5 & 0.5 & 0 \end{bmatrix}$$

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The probabilities for the next state are given by

$$\mathbf{x}' = \mathbf{x} \cdot \mathbf{P}$$

In general:

$$\begin{array}{rcl} \mathbf{x}_{k+1} & = & \mathbf{x}_k \cdot \mathbf{P} = & (\mathbf{x}_{k-1} \cdot \mathbf{P}) \cdot \mathbf{P} = & \cdots = & \mathbf{x}_0 \cdot \underbrace{\mathbf{P} \cdot \mathbf{P} \cdot \cdots \mathbf{P}}_{k+1 \text{ times}} \\ \\ \mathbf{x}_{k+1} & = & \mathbf{x}_0 \cdot \mathbf{P}^{k+1} \end{array}$$

Stavros Tripakis: EECS 144/244 - Probabilistic Systems

Interesting questions:

- ► After *k* steps, what is the likelihood that the system is at state *i*?
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Which of these questions do we have an answer to already?

► After *k* steps, what is the likelihood that the system is at state *i*?

$$\mathbf{x}_k = \mathbf{x}_0 \cdot \mathbf{P}^k$$

Interesting questions:

- \checkmark After k steps, what is the likelihood that the system is at state i?
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Let's see how to answer the rest.

Markov chains and graphs

A Markov chain has a graph structure.

We can partly answer the questions simply by studying this structure, **completely ignoring the probability numbers**.

n-step transition probabilities

Let \mathbf{P}_{ij}^n be the (i,j) element of \mathbf{P}^n .

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n-step transition probabilities

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What does \mathbf{P}_{ij}^n represent?

 $\mathbf{P}_{ij}^n = P(s_{k+n} = j \mid s_k = i)$: probability that, starting from state i, after n steps the state will be j.

State i is absorbing if $P_{ii} = 1$. This implies $P_{ij} = 0$ for all $j \neq i$.

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Two states i and j communicate, written $i \leftrightarrow j$, if i is accessible from j and vice-versa.

A set of states that communicate is called a class.

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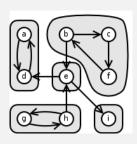
Two states i and j communicate, written $i \leftrightarrow j$, if i is accessible from j and vice-versa.

A set of states that communicate is called a *class*. A class is a strongly-connected component.

Strongly-connected components

In a directed graph $G=(V,\to)$, a strongly-connected component (SCC) is a subset of nodes $C\subseteq V$, such that every node in C is reachable from every other node in C.

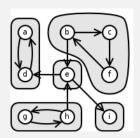
C is called *maximal* if we cannot add more nodes to C and still preserve its SCC property, i.e., $\not\exists C' \supset C$ s.t. C' is also a SCC.



The acyclic graph of maximal SCCs

The set of all maximal SCCs of a graph defines a new graph, where nodes are maximal SCCs, $C_1, C_2, ..., C_m$.

An edge $C_i \sim C_j$ exists iff $C_i \neq C_j$ and there is a node in C_i that has a successor node in C_j .



This graph of SCCs is by definition acyclic: why?

Irreducible Markov chains

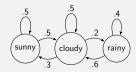
The Markov chain is *irreducible* if it has only one class, i.e., all states communicate with each other.

Irreducible Markov chains

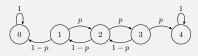
The Markov chain is *irreducible* if it has only one class, i.e., all states communicate with each other.

This says that the whole Markov chain is a SCC.

Weather model (irreducible):



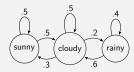
Gambling model (reducible):



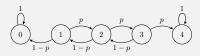
Learning model (reducible):



Weather model (irreducible):



Gambling model (reducible):



Learning model (reducible):



States 0, 4, and L are absorbing states.

Recurrent and transient states

Let i, j be two states. Define:

 $f_{ij}^n =$ probability that, starting in i, the first transition into j happens after n steps

$$f_{ij} = \sum_{n=1}^{\infty} f_{ij}^n = \text{probability of reaching } j \text{ from } i \text{ in any } \# \text{ steps}$$

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State i is:

- Recurrent if $f_{ii} = 1$.
- ▶ Transient if $f_{ii} < 1$.

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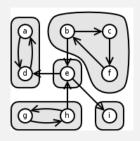
- Recurrent if $f_{ii} = 1$.
- ▶ Transient if $f_{ii} < 1$.

It follows that:

- ▶ If i is recurrent then system visits i infinitely often (with probability 1).
- ► If i is transient then system visits i only finitely many times (with probability 1).

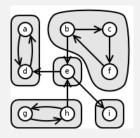
A SCC C is called *terminal* if there is no C' such that $C \leadsto C'$.

Otherwise C is called *transient*.



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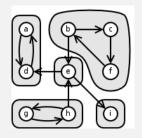


Terminal SCCs: $\{a, d\}$ and $\{i\}$.

Transient SCCs: $\{b, c, f\}$, $\{e\}$, and $\{g, h\}$.

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Terminal SCCs: $\{a, d\}$ and $\{i\}$.

Transient SCCs: $\{b,c,f\}$, $\{e\}$, and $\{g,h\}$.

- ▶ Recurrent states = states belonging to terminal SCCs.
- ► Transient states = states belonging to transient SCCs.

If i is recurrent and $i \leftrightarrow j$ then j is also recurrent.

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In a finite Markov chain M, some states will be recurrent. Why?

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Does this hold also for infinite Markov chains?

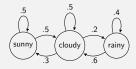
Viewed as a graph, every finite Markov chain M has at least one terminal maximal SCC.

Theorem

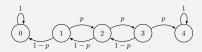
In the long run the amount of time that M spends in transient SCCs is 0.

Therefore, the probability that after some time M will reach a terminal SCC and remain forever there is 1.

Weather model: all states visited infinitely often.

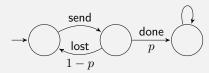


Gambling model: eventually system enters either 0 or 4 and then stays there forever.

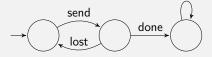




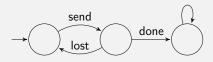
In a probabilistic system, the behavior where the message keeps getting lost after being sent has probability 0, **independently of the value of** p (provided p > 0):



In a non-deterministic system, it is possible that the message always gets lost:



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However, we can add fairness constraints to ensure that it does not, e.g.,:

$$\underbrace{\left((\Box \diamondsuit \mathsf{send}) \to (\diamondsuit \mathsf{done})\right)}_{\mathsf{fairness constraint}} \to (\mathsf{what we want})$$

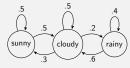
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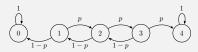
If this is all we need, probabilities are an overkill.

Weather model: all states visited infinitely often.



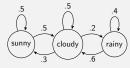
How much time does it rain on the average?

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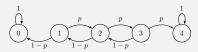


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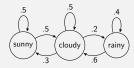
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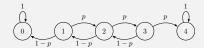


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How much time does it rain on the average?

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When does the gambler go bankrupt?



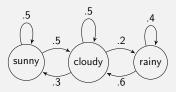
State i has period $d \in \mathbb{N}$ if $\mathbf{P}_{ii}^n = 0$ whenever n is not divisible by d, and d is the largest such number.

If d=1 (i.e., $\mathbf{P}_{ii}^n>0$ for all n) then state i is called *aperiodic*.

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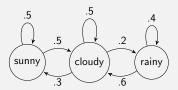
Examples:



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Examples:

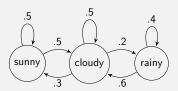


All states are aperiodic.

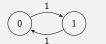
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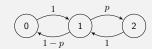
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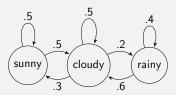




Both states have period 2. All states have period 2.

Ergodic states

An aperiodic recurrent state is called ergodic.



All states are ergodic.

Stationary distribution

Theorem

Let M be a finite, irreducible Markov chain where all states of M are ergodic. Then the limit

$$\lim_{k\to\infty}\mathbf{P}^k_{ij}$$

exists and is independent of i (i.e., $\forall i, i' : \lim_{k \to \infty} \mathbf{P}_{ij}^k = \lim_{k \to \infty} \mathbf{P}_{i'j}^k$).

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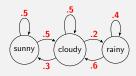
Furthermore, letting

$$\pi_j = \lim_{k \to \infty} \mathbf{P}_{ij}^k$$

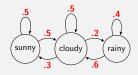
then $\pi = [\pi_1 \ \pi_2 \ \cdots \ \pi_n]$ is the unique non-negative solution of

$$\pi = \pi \cdot \mathbf{P}$$

 π is called the *stationary distribution*.



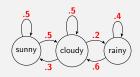
$$\mathbf{P} = \left[\begin{array}{ccc} 0.5 & 0.5 & 0 \\ 0.3 & 0.5 & 0.2 \\ 0 & 0.6 & 0.4 \end{array} \right]$$



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$$\mathbf{P} = \begin{bmatrix} 0.5 & 0.5 & 0 \\ 0.3 & 0.5 & 0.2 \\ 0 & 0.6 & 0.4 \end{bmatrix} \qquad \forall k > 15 : \mathbf{P}^k = \begin{bmatrix} 0.3103 & 0.5172 & 0.1724 \\ 0.3103 & 0.5172 & 0.1724 \\ 0.3103 & 0.5172 & 0.1724 \end{bmatrix}$$

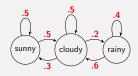
What does this imply?



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What does this imply?

for any probability vector $\mathbf{x}_0, \forall k > 15: \mathbf{x}_0 \cdot \mathbf{P}^k = \pi = [0.3103 \ 0.5172 \ 0.1724]$

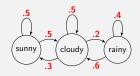


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i.e., stationary distribution π independent from the initial state distribution \mathbf{x}_0 .



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i.e., stationary distribution π independent from the initial state distribution \mathbf{x}_0 .

So how much time does it rain on the average?

Stationary distribution

If the chain is not ergodic, the limit may not exist, e.g.,

$$\mathbf{P} = \left[\begin{array}{cc} 0 & 1 \\ 1 & 0 \end{array} \right]$$

Analysis of Markov chains

Interesting questions:

- ✓ After k steps, what is the likelihood that the system is at state i?
- ▶ In the long run, how much time does the system spend at state *i*? (i.e., how often is *i* visited?)
- What is the probability that the system will ever reach a given state (or group of states)?
- ▶ What is the expected time until the system reaches a given state (or group of states)?

Which of these questions do we have an answer to already?

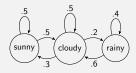
Analysis of Markov chains

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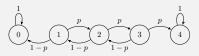
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Examples

Weather model:



Gambling model:



How much time does an average game last?

Learning model:



How long until we learn something?

Order the states of a Markov chain M so that $\{1,2,...,t\}$ is the set of transient states.

Let

$$\mathbf{P}_T = \begin{bmatrix} p_{11} & p_{12} & \cdots & p_{1t} \\ p_{21} & p_{22} & \cdots & p_{2t} \\ \vdots & & \ddots & \vdots \\ p_{t1} & p_{t2} & \cdots & p_{tt} \end{bmatrix}$$

Observation: some rows of \mathbf{P}_T sum to < 1 (otherwise this would be a SCC).

Let

 $q_{ij}=$ mean time spent in j, given that the system starts in i

Then

$$q_{ij} = \delta_{ij} + \sum_{k} p_{ik} \cdot q_{kj}$$

where

$$\delta_{ij} = \begin{cases} 1 & \text{if } i = j \\ 0 & \text{otherwise.} \end{cases}$$

Let

 $q_{ij}=$ mean time spent in j, given that the system starts in i

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$$q_{ij} = \delta_{ij} + \sum_{k} p_{ik} \cdot q_{kj}$$
$$= \delta_{jk} + \sum_{k=1}^{t} p_{ik} \cdot q_{kj} \qquad why?$$

where

$$\delta_{ij} = \begin{cases} 1 & \text{if } i = j \\ 0 & \text{otherwise.} \end{cases}$$

Let

 $q_{ij}=% \frac{1}{2}\left(ij\right) =0$ mean time spent in j, given that the system starts in i

Then

$$q_{ij} = \delta_{ij} + \sum_{k} p_{ik} \cdot q_{kj}$$
$$= \delta_{jk} + \sum_{k=1}^{t} p_{ik} \cdot q_{kj} \qquad why?$$

Because $q_{kj} = 0$ when k is recurrent (cannot move from recurrent state to transient state).

where

$$\delta_{ij} = \begin{cases} 1 & \text{if } i = j \\ 0 & \text{otherwise.} \end{cases}$$

Let

$$\mathbf{Q} = \begin{bmatrix} q_{11} & q_{12} & \cdots & q_{1t} \\ q_{21} & q_{22} & \cdots & q_{2t} \\ \vdots & & \ddots & \vdots \\ q_{t1} & q_{t2} & \cdots & q_{tt} \end{bmatrix}$$

Then

$$\mathbf{Q} = \mathbf{I} + \mathbf{P}_T \cdot \mathbf{Q}$$

where ${f I}$ is the identity matrix of size t.

Let

$$\mathbf{Q} = \begin{bmatrix} q_{11} & q_{12} & \cdots & q_{1t} \\ q_{21} & q_{22} & \cdots & q_{2t} \\ \vdots & & \ddots & \vdots \\ q_{t1} & q_{t2} & \cdots & q_{tt} \end{bmatrix}$$

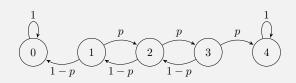
Then

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where I is the identity matrix of size t.

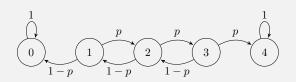
It can be shown that $I - P_T$ is invertible. Therefore:

$$\mathbf{Q} = (\mathbf{I} - \mathbf{P}_T)^{-1}$$



$$\mathbf{Q} = (\mathbf{I} - \mathbf{P}_T)^{-1} = \begin{bmatrix} 1 & -p & 0 \\ p - 1 & 1 & -p \\ 0 & p - 1 & 1 \end{bmatrix}^{-1} = \cdots$$

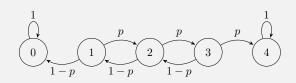
$$\mathbf{Q}_{p=0} = \begin{bmatrix} 1 & 0 & 0 \\ 1 & 1 & 0 \\ 1 & 1 & 1 \end{bmatrix} \quad \mathbf{Q}_{p=1} = \begin{bmatrix} 1 & 1 & 1 \\ 0 & 1 & 1 \\ 0 & 0 & 1 \end{bmatrix} \quad \mathbf{Q}_{p=\frac{1}{2}} = \begin{bmatrix} 1.5 & 1 & 0.5 \\ 1 & 2 & 1 \\ 0.5 & 1 & 1.5 \end{bmatrix}$$



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What is the average playing time with $p = \frac{1}{2}$?



$$\mathbf{Q} = (\mathbf{I} - \mathbf{P}_T)^{-1} = \begin{bmatrix} 1 & -p & 0 \\ p - 1 & 1 & -p \\ 0 & p - 1 & 1 \end{bmatrix}^{-1} = \cdots$$

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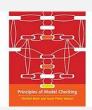
What is the average playing time with $p = \frac{1}{2}$? 3 if I start with \$1 or \$3. 4 if I start with \$2.

Analysis of Markov chains

Interesting questions:

- \checkmark After k steps, what is the likelihood that the system is at state i?
- ✓ In the long run, how much time does the system spend at state i? (i.e., how often is i visited?)
- ► What is the probability that the system will ever reach a given state (or group of states)?
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REACHABILITY PROBABILITIES IN MARKOV CHAINS



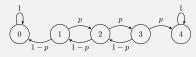
Material taken mainly from [BK08], Chapter 10.

Reachability question for Markov chains

Let M be a Markov chain and B a set of states of M.

Reachability question for Markov chains: what is the probability of reaching B?

Note: this is not a yes/no question, as in standard model-checking. Here, we want to compute a probability $p \in [0,1]$.



When does the gambler go bankrupt?

Model in PRISM:

```
dtmc // this model is a Markov chain

additional constants and chain

set = 0.7;

module M

s: [0..4] init 2;

g: [] s=0 -> (s'=0);

[] s=1 -> p: (s'=2) + (1-p): (s'=0);

[] s=2 -> p: (s'=3) + (1-p): (s'=1);

[] s=3 -> p: (s'=4) + (1-p): (s'=2);

[] s=4 -> (s'=4);

dtmc // this model is a Markov chain

additional constants and chain

additional constants and chain

dtmc // this model is a Markov chain

additional constants and chain

addit
```

PCTL formula in PRISM:

P=? [F s=4]

PRISM answers:

p	answer
0	0
0.5	$0.499999 \cdots$
0.7	$0.844827 \cdots$
1	1

Computing reachability probabilities

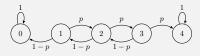
Let x_i be the probability that the target set B is reached starting from state i.

Then:

- ▶ If $i \in B$ then $x_i = 1$.
- ▶ If i cannot reach B in the graph sense, then $x_i = 0$.
- Otherwise

$$x_i = \sum_j p_{ij} \cdot x_j$$

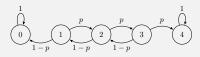
This forms a set of linear equations. For finite chains it is finite and is guaranteed to have a unique solution.



When does the gambler go bankrupt?

$$\begin{array}{rcl} x_0 & = & 0 \\ x_4 & = & 1 \\ x_1 & = & p \cdot x_2 + (1-p) \cdot x_0 = p \cdot x_2 \\ x_2 & = & p \cdot x_3 + (1-p) \cdot x_1 \\ x_3 & = & p \cdot x_4 + (1-p) \cdot x_2 = p + (1-p) \cdot x_2 \end{array}$$

For
$$p = \frac{1}{2}$$
, $x_2 = \frac{1}{2}$. For $p = 0.7$, $x_2 = 0.8448$.



When does the gambler go bankrupt?

$$x_0=0$$
 This is necessary for uniqueness of solution. Why?
$$x_4=1$$

$$x_1=p\cdot x_2+(1-p)\cdot x_0=p\cdot x_2$$

$$x_2=p\cdot x_3+(1-p)\cdot x_1$$

$$x_3=p\cdot x_4+(1-p)\cdot x_2=p+(1-p)\cdot x_2$$

For
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and MARKOV DECISION PROCESSES

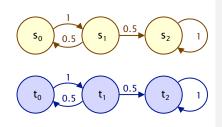
Composition of Markov chains

Suppose we want to compose the following two Markov chains:

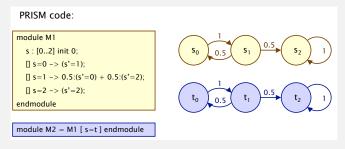
PRISM code:

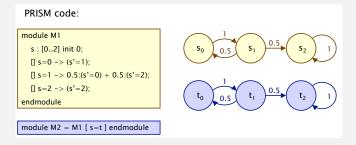
module M1
$$\begin{split} s: [0..2] & \text{ init } 0; \\ [] & s=0 \rightarrow (s'=1); \\ [] & s=1 \rightarrow 0.5: (s'=0) + 0.5: (s'=2); \\ [] & s=2 \rightarrow (s'=2); \\ endmodule \end{split}$$

module M2 = M1 [s=t] endmodule

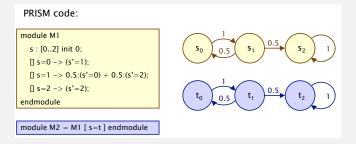


Several figures due to Dave Parker.

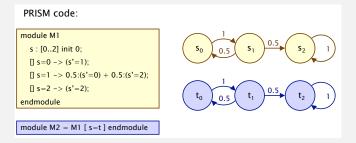




$$P((s_1, t_1) \mid (s_0, t_0)) =$$

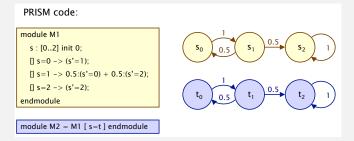


$$P((s_1, t_1) \mid (s_0, t_0)) = 1 \cdot 1 = 1$$



$$P((s_1, t_1) \mid (s_0, t_0)) = 1 \cdot 1 = 1$$

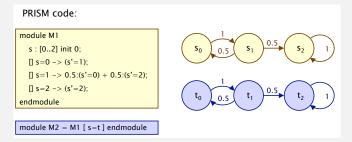
 $P((s_2, t_2) \mid (s_1, t_1)) =$



$$P((s_1, t_1) \mid (s_0, t_0)) = 1 \cdot 1 = 1$$

 $P((s_2, t_2) \mid (s_1, t_1)) = \frac{1}{2} \cdot \frac{1}{2} = \frac{1}{4}$

What does the synchronous composition of these processes look like?

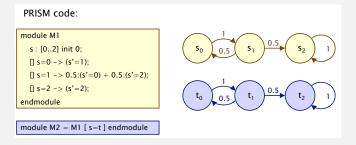


$$P((s_1, t_1) \mid (s_0, t_0)) = 1 \cdot 1 = 1$$

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Stavros Tripakis: EECS 144/244 - Probabilistic Systems

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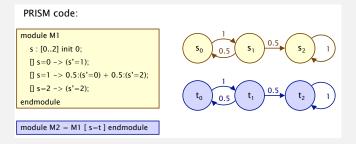


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Is the synchronous composition of two Markov chains a Markov chain?

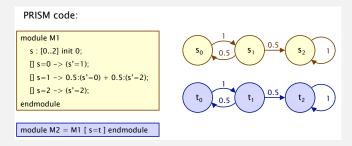
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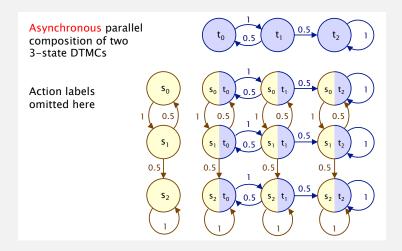


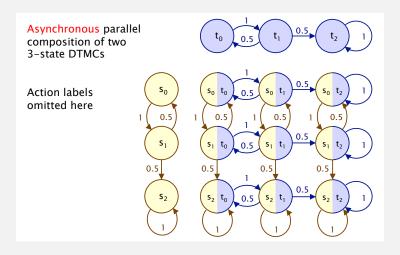
$$P((s_1, t_1) \mid (s_0, t_0)) = 1 \cdot 1 = 1$$

 $P((s_2, t_2) \mid (s_1, t_1)) = \frac{1}{2} \cdot \frac{1}{2} = \frac{1}{4}$

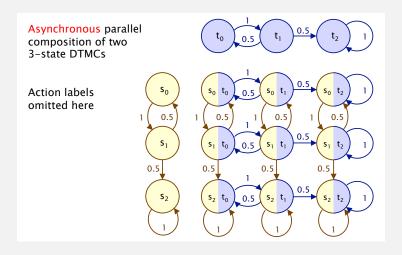
Is the synchronous composition of two Markov chains a Markov chain? Yes!







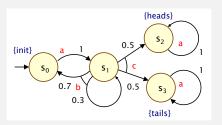
Is the asynchronous composition of two Markov chains a Markov chain?



Is the asynchronous composition of two Markov chains a Markov chain? No! It is a Markov Decision Process.

Markov Decision Processes (MDPs)

Combine non-deterministic and probabilistic choice.

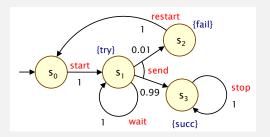


Intuitive semantics:

- First choose action non-deterministically among possible actions.
 - ▶ In state s_0 , only one possible action, a.
 - ▶ In state s_1 , two possible actions, b and c.
- ► Then, given chosen action, throw a dice a pick successor state w.r.t. the specified probability distribution for that action.

Markov Decision Processes (MDPs)

Non-determinism has multiple uses, as in discrete systems. E.g., useful to model abstraction:



At state s_1 , if channel is ready attempt to send, otherwise wait. Details of when channel is ready are not modeled.

Model-checking MDPs

Answers queries like:

- Byzantine agreement protocol
 - $P_{min=?}$ [F (agreement ∧ rounds ≤ 2)]
 - "what is the minimum probability that agreement is reached within two rounds?"
- · CSMA/CD communication protocol
 - P_{max=?} [F collisions=k]
 - "what is the maximum probability of k collisions?"
- Self-stabilisation protocols
 - $-P_{min=?}$ [$F^{\leq t}$ stable]
 - "what is the minimum probability of reaching a stable state within k steps?"

See PRISM web site and literature for details: http://www.prismmodelchecker.org/

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