Flexible, fine-grained distributed access control

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We're all ears

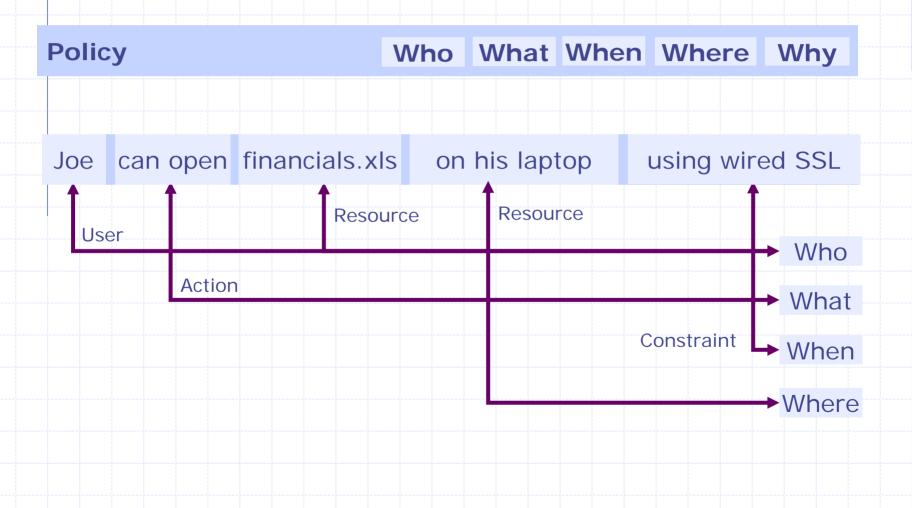


What policy concepts are important in healthcare?
What kind of systems should understand or enforce these policies?
How can tech geeks be useful?
What's all this talk about Brazilian skiing?

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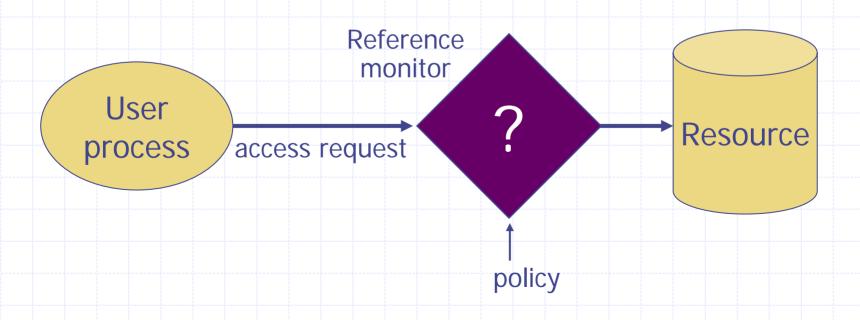
Enterprise Access Control



Traditional mechanisms

Assumptions

- System knows who the user is
 - User has entered a name and password, or other info
- Access requests pass through gatekeeper
 - System must not allow monitor to be bypassed



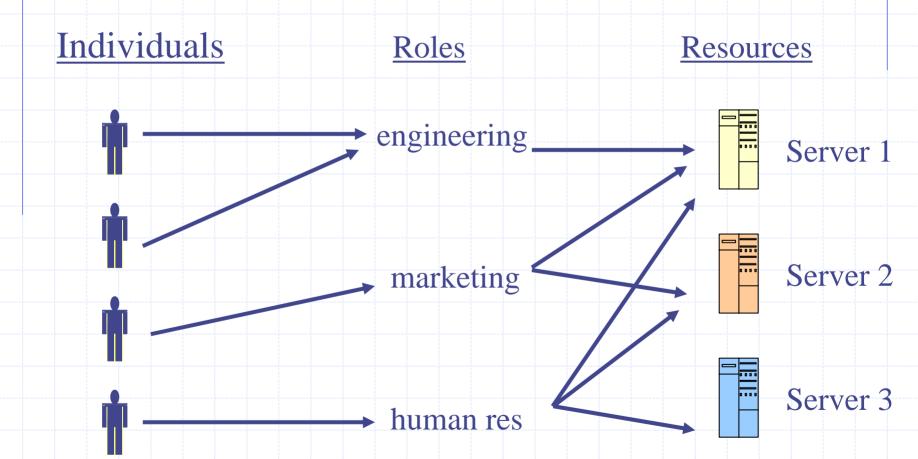
Access control matrix [Lampson]

Objects

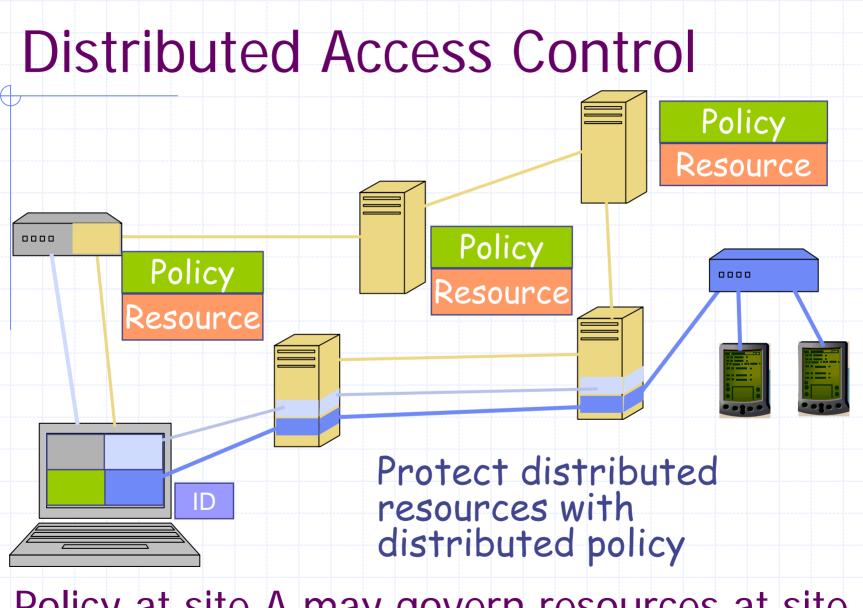
		File 1	File 2	File 3		File n
Subjects {	User 1	read	write	-	_	read
	User 2	write	write	write	-	-
	User 3		-		read	read
	User m	read	write	read	write	read

Access control list (ACL): column of matrix, often stored at resource

Role-Based Access Control



Leverage: user's change more frequently than roles



Policy at site A may govern resources at site B

Decentralized Policy Example



Alice





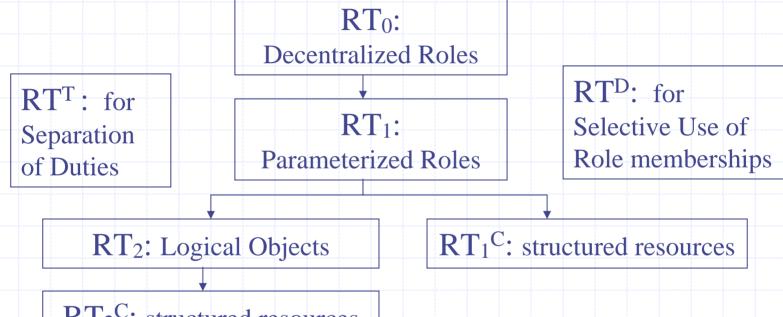
Alice is a student



StateU is a university

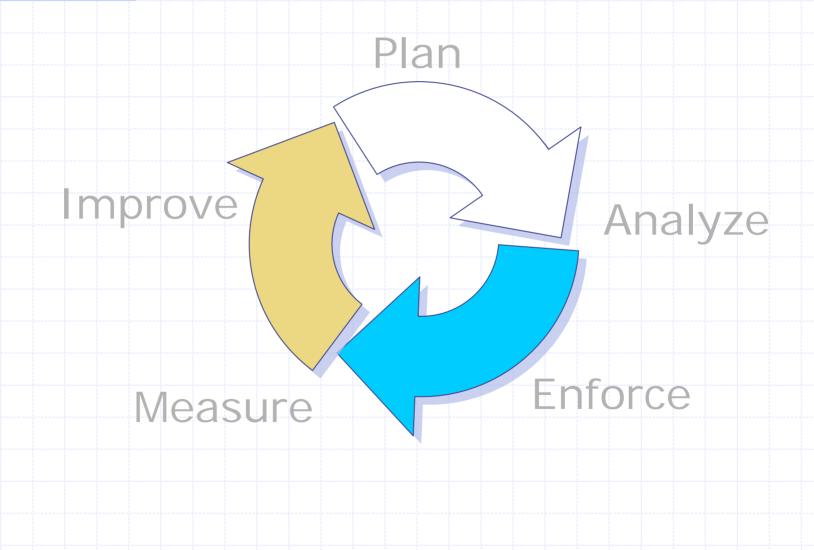


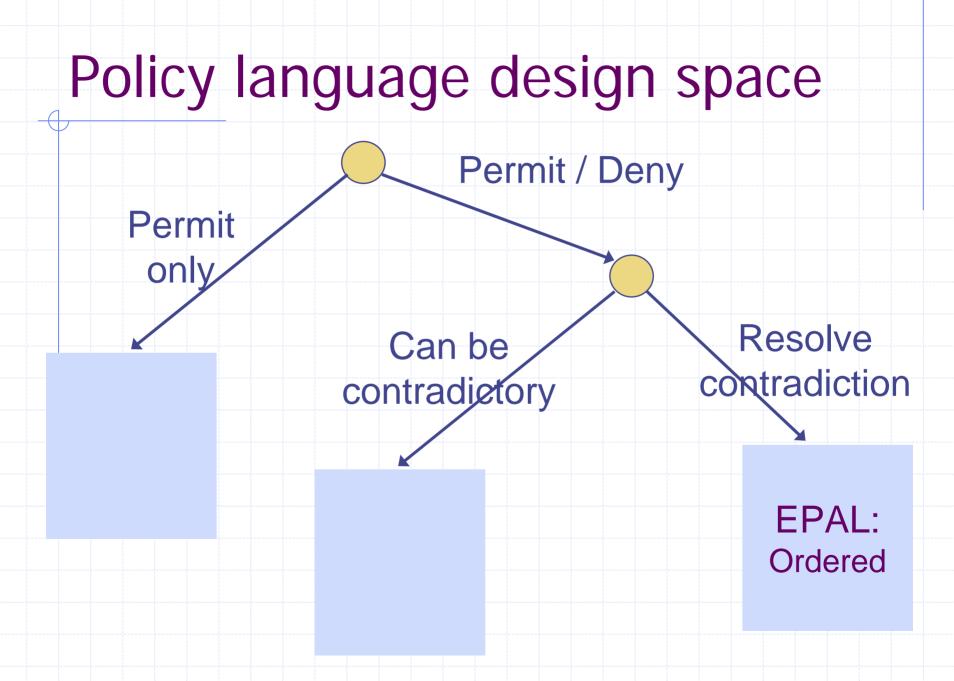
Role-based Trust-management (RT)



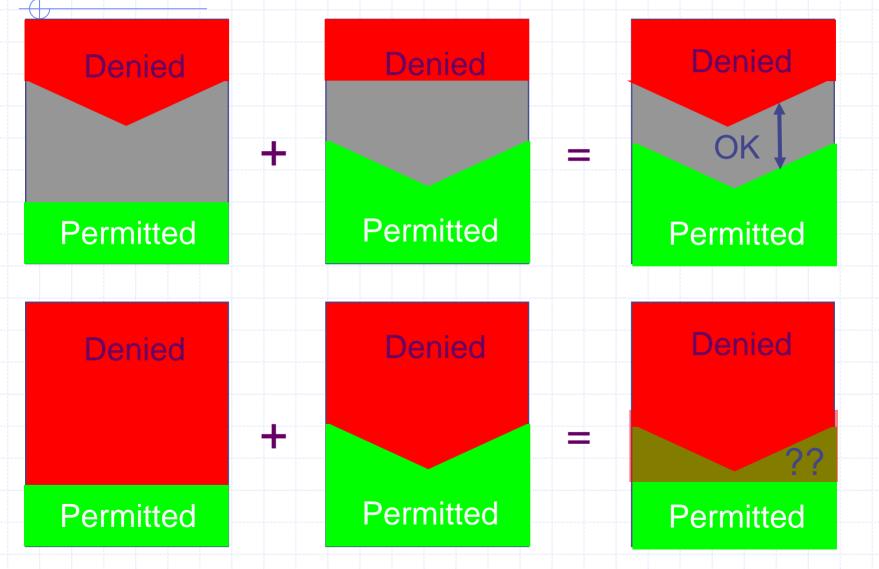
 RT_2^C : structured resources

Policy Management Lifecycle





Policy Combination



Contextual Integrity

Framework for privacy: Concept of contextual integrity Formalization in Linear Temporal Logic Application to privacy laws: HIPAA, GLBA, COPPA Related Work RBAC, XACML, P3P, EPAL

Overview of Contextual Integrity

Transfer of information between agents "Alice give Bob information about Charlie" Categorization Agents grouped into roles Information categorized by types Basic policy statements Manager may read employee's performance data If m is a manager, e is an employee, d is performance data about e, and m is e's manager then m may read d

Formalization in Temporal Logic

Syntax of logic

$$\begin{split} \varphi ::= \operatorname{send}(p_1, p_2, m) \mid \operatorname{contains}(m, q, t) \mid \operatorname{inrole}(p, r) \mid \operatorname{incontext}(p, c) \mid t \in t' \mid \\ \varphi \wedge \varphi \mid \neg \varphi \mid \varphi \mathcal{U}\varphi \mid \varphi \mathcal{S}\varphi \mid \bigcirc \varphi \mid \exists x : \tau.\varphi \end{split}$$

Formula representing contextual norms

 $\sigma \models \Box \forall p_1, p_2, q : P. \forall m : M. \forall t : T. \operatorname{incontext}(p_1, c) \land$

 $\operatorname{send}(p_1, p_2, m) \wedge \operatorname{contains}(m, q, t) \rightarrow \bigvee$

$$\varphi^+ \wedge \bigwedge$$

 $\varphi^+ \in \operatorname{norms}^+(c)$

$$\varphi^{-} \in \operatorname{norms}^{-}(e)$$

where norms have specific forms

positive norm: $\operatorname{inrole}(p_1, \hat{r}_1) \wedge \operatorname{inrole}(p_2, \hat{r}_2) \wedge \operatorname{inrole}(q, \hat{r}) \wedge (t \in \hat{t}) \wedge \theta \wedge \psi$ negative norm: $\operatorname{inrole}(p_1, \hat{r}_1) \wedge \operatorname{inrole}(p_2, \hat{r}_2) \wedge \operatorname{inrole}(q, \hat{r}) \wedge (t \in \hat{t}) \wedge \theta \to \psi$

Policy Operations and Relations

Standard automated LTL tools are applicable

- Policy consistency: LTL satisfiability
- Refinement: logical implication
- Combination: conjunction and disjunction
- Strong compliance: satisfiability
- Weak compliance: computable efficiently using concepts from LTL runtime verification

Application: HIPAA

Privacy Rule

- Covered entities (e.g. hospitals) can give protected health information about patients to health care providers
 - Sender role: Covered entity
 - Recipient role: Health care provider
 - Subject role: Patient
 - Information type: Protected health information

 $\operatorname{inrole}(p_1, \operatorname{covered-entity}) \land \operatorname{inrole}(p_2, \operatorname{provider}) \land \operatorname{inrole}(q, \operatorname{patient}) \land (t \in phi)$

Application: GLBA

Privacy Rule

- Financial institutions must notify consumers if they share their non-public personal information with non-affiliated companies, but the notification may occur either before or after the information sharing occurs.
 - Sender role: Financial institution
 - Recipient role: Non-affiliated company
 - Subject role: Consumer
 - Information type: Non-public personal information
 - Temporal condition: Notify data subject

 $\operatorname{inrole}(p_1, institution) \land \operatorname{inrole}(p_2, non-affiliate) \land \operatorname{inrole}(q, consumer) \land (t \in npi) \rightarrow (t \in npi) \land ($

 \diamondsuit send $(p_1, q, privacy-notice) \lor \diamondsuit$ send $(p_1, q, privacy-notice)$

Comparison

Role-based access control

- No subject of data, attributes, temporal conditions
- XACML
 - Attributes handled incorrectly (inheritance)
 - Combination occurs functionally, not logically
- EPAL
 - Obligations treated as uninterpreted symbols
 - Can only enforce week compliance
- P3P
 - Contains only simple opt-in / opt-out conditions

