

## Is Truly Real-Time Computing Becoming Unachievable?

### Edward A. Lee

Robert S. Pepper Distinguished Professor and Chair of EECS, UC Berkeley

Keynote Talk

Real-Time and Embedded Technology and Applications Symposium (RTAS)

Bellevue, WA April 3 - April 6, 2007

## *The Vision*: Reliable and Evolvable Networked Time-Sensitive Systems, Integrated with Physical Processes

Orchestrated networked resources built with sound design principles on suitable abstractions.



*The Vision*: Reliable and Evolvable Networked Time-Sensitive Systems, Integrated with Physical Processes

Orchestrated networked resources built with sound design principles on suitable abstractions.



## • • • A Fact About Programs

Correct execution of a program in C, C#, Java, Haskell, etc. has nothing to do with how long it takes to do anything. All our computation and networking abstractions are built on this premise.



Timing of programs is not repeatable, except at very coarse granularity.

Programmers have to step *outside* the programming abstractions to specify timing behavior.

# Techniques that Exploit this Fact

- Programming languages
- Virtual memory
- Caches
- Dynamic dispatch
- Speculative execution
- Power management (voltage scaling)
- Memory management (garbage collection)
- Just-in-time (JIT) compilation
- Multitasking (threads and processes)
- Component technologies (OO design)
- Networking (TCP)
- o ...

# • • A Story



In "fly by wire" aircraft, certification of the software is extremely expensive. Regrettably, it is not the software that is certified but the entire system. If a manufacturer expects to produce a plane for 50 years, it needs a 50-year stockpile of fly-by-wire components that are all made from the same mask set on the same production line. Even a slight change or "improvement" might affect timing and require the software to be recertified.

## Abstraction Layers



The purpose for an abstraction is to hide details of the implementation below and provide a platform for design from above.

## Abstraction Layers



Every abstraction layer has failed for real-time programs.

The design *is* the implementation.



How about "raising the level of abstraction" to solve these problems?

# But these higher abstractions rely on an increasingly problematic fiction: WCET

A war story:

Ferdinand et al. determine the WCET of astonishingly simple avionics code from Airbus running on a Motorola ColdFire 5307, a pipelined CPU with a unified code and data cache. Despite the software consisting of a fixed set of non-interacting tasks containing only simple control structures, their solution required detailed modeling of the seven-stage pipeline and its precise interaction with the cache, generating a large integer linear programming problem. The technique successfully computes WCET, but only with many caveats that are increasingly rare in software.

Fundamentally, the ISA of the processor has failed to provide an adequate abstraction.

C. Ferdinand et al., "Reliable and precise WCET determination for a real-life processor." EMSOFT 2001.



## The Key Problem

Electronics technology delivers highly and precise timing...

... and the overlaying software abstractions discard it.

# Real-Time and Concurrency are Integrally Intertwined

Threads and objects dominate concurrent software.

- *Threads*: Sequential computation with shared memory.
- Objects: Collections of state variables with procedures for observing and manipulating that state.

Even distributed objects create the illusion of shared memory through proxies.

- The components (objects) are (typically) not active.
- Threads weave through objects in unstructured ways.
- This is the source of many software problems.



## Nontrivial software written with threads, and locks are incomprehensible to humans.

## Is Concurrency Hard?



# • • • It is Threads that are Hard!

Threads are sequential processes that share memory. From the perspective of any thread, the entire state of the universe can change between any two atomic actions (itself an ill-defined concept).

Imagine if the physical world did that...



Threads are wildly nondeterministic.

The programmer's job is to prune away the nondeterminism by imposing constraints on execution order (e.g., mutexes) and limiting shared data accesses (e.g., OO design).

# • • We Can Incrementally Improve Threads

- Object Oriented programming
- Coding rules (Acquire locks in the same order...)
- Libraries (Stapl, Java 5.0, ...)
- Patterns (MapReduce, ...)
- Transactions (Databases, ...)
- Formal verification (Blast, thread checkers, ...)
- Enhanced languages (Split-C, Cilk, Guava, ...)
- Enhanced mechanisms (Promises, futures, ...)

# But is it enough to refine a mechanism with flawed foundations?

## Do Threads Provide a Sound Foundation?

If the foundation is bad, then we either tolerate *brittle designs* that are difficult to make work, or we have to rebuild from the foundations.

Note that this whole enterprise is held up by threads



# • • • What are Brittle Designs?

#### Small changes have big consequences...

Patrick Lardieri, *Lockheed Martin ATL*, about a vehicle management system in the JSF program:

"Changing the instruction memory layout of the Flight Control Systems Control Law process to optimize 'Built in Test' processing led to an unexpected performance change - System went from meeting real-time requirements to missing most deadlines due to a change that was expected to have no impact on system performance."

National Workshop on High-Confidence Software Platforms for Cyber-Physical Systems (HCSP-CPS) Arlington, VA November 30 –December 1, 2006

# The Current State of Affairs

We build real-time software on abstractions where time is irrelevant using concurrency models that are incomprehensible.



# Just think what we could do with the right abstractions!



### Reintroduce time into the core abstractions:

• Bottom up: Make timing repeatable.

• Top down: Timed, concurrent components.

# Bottom Up: Make Timing Repeatable

We need a major historical event like this one:

In 1980, Patterson and Ditzel did not invent reduced instruction set computers (RISC machines).

See D. A. Patterson and D. R. Ditzel. "The case for the reduced instruction set computer." *ACM SIGARCH Computer Architecture News*, 8(6):25–33, Oct. 1980.

# It is Time for Another Major Historical Event

In 2007, Edwards and Lee did not invent precision-timed computers (PRET machines).

See S. Edwards and E. A. Lee, "**The Case for the Precision Timed (PRET) Machine**," to appear in the *Wild and Crazy Ideas* Track of the *Design Automation Conference* (DAC), June 2007.

see: http://www.eecs.berkeley.edu/Pubs/TechRpts/2006/EECS-2006-149.html

## Can Hardware Deliver on Timed Semantics?

PRET (Precision Timing) Machines Make temporal behavior as important as logical function.

Timing precision is easy to achieve if you are willing to forgo performance. Let's not do that. Challenges:

- Memory hierarchy (scratchpads?)
- Deep pipelines (interleaving?)
- ISAs with timing (deadline instructions?)
- Predictable memory management (Metronome?)
- Languages with timing (discrete events? Giotto?)
- Predictable concurrency (synchronous languages?)
- Composable timed components (actor-oriented?)
- Precision networks (TTA? Time synchronization?)
- Dynamic adaptibility (admission control?)

## Making PRET Machines Practical

- Start with hardware designs on FPGAs
- Use soft cores
- Add precision-timing instructions
- Scale up from there



BEE 2, FPGA system, from BWRC



Ramp blue experimental platform, from the RAMP project.

e.g. Stephen Edwards (Columbia) has achieved software designs with ~40ns timing precision on simple soft cores. Source code is smaller and simpler than VHDL specification of comparable hardware.



### Reintroduce time into the core abstractions:

### • Bottom up: Make timing repeatable.

• Top down: Timed, concurrent components.

# New Component Technology is more Palatable than New Languages

### • It leverages:

- Language familiarity
- Component libraries
- Legacy subsystems
- Design tools
- The simplicity of sequential reasoning
- It allows for innovation in
  - Distributed time-sensitive system design
  - Hybrid systems design
  - Service-oriented architectures
- Software is intrinsically concurrent
  - Better use of multicore machines
  - Better use of networked systems
  - Better potential for robust design



#### The established: Object-oriented:



The alternative: Actor oriented:



#### The First (?) Actor-Oriented Programming Language The On-Line Graphical Specification of Computer Procedures W. R. Sutherland, Ph.D. Thesis, MIT, 1966



MIT Lincoln Labs TX-2 Computer



Bert Sutherland with a light pen



Bert Sutherland used the first acknowledged objectoriented framework (Sketchpad, created by his brother, Ivan Sutherland) to create the first actor-oriented programming language (which had a visual syntax).

Partially constructed actor-oriented model with a class definition (top) and instance (below).

# Examples of Actor-Oriented Systems

- o CORBA event service (distributed push-pull)
- o ROOM and UML-2 (dataflow, Rational, IBM)
- o VHDL, Verilog (discrete events, Cadence, Synopsys, ...)
- o LabVIEW (structured dataflow, National Instruments)
- o Modelica (continuous-time, constraint-based, Linkoping)
- o OPNET (discrete events, Opnet Technologies)
- o SDL (process networks)
- o Occam (rendezvous)
- o Simulink (Continuous-time, The MathWorks)
- o SPW (synchronous dataflow, Cadence, CoWare)
- o ...

The semantics of these differ considerably, with significantly different approaches to concurrency.

Most of these are domain specific.

Many of these have visual syntaxes.

# • • Challenges

#### The technology is immature:

- Commercial actor-oriented systems are domain-specific
- Development tools are limited
- Little language support in C++, C#, Java
- Modularity mechanisms are underdeveloped
- Type systems are primitive
- Compilers (called "code generators") are underdeveloped
- Formal methods are underdeveloped
- Libraries are underdeveloped

We are addressing these problems.

# Enter Ptolemy II: Our Laboratory for Experiments with Models of Computation



# Ptolemy II: Functionality of Components is Given in C or Java (which can wrap C, C++, Perl, Python, MATLAB, Web services, Grid services, ...)





## Using DE Semantics in Distributed Real-Time Systems

- DE is usually a simulation technology.
- Distributing DE is done for acceleration.
- Hardware design languages (e.g. VHDL) use DE where time stamps are literally interpreted as real time.

- We are using DE for distributed real-time software, binding time stamps to real time only where necessary.
- *PTIDES*: Programming Temporally Integrated Distributed Embedded Systems (work done with Yang Zhao and Jie Liu).

### Consider a scenario:



*Assumption*: Clocks on the distributed platforms are synchronized to some known precision (e.g. NTP, IEEE 1588)



### Bind model time to real time at the sensors:



Bind model time to real time at the actuators:



### Schedulability is not violating these timing inequalities.



Conservative distributed DE (Chandy & Misra) would block actuation unnecessarily.



Optimistic distributed DE (Jefferson) would require being able to roll back the physical world.



# PTIDES uses static causality analysis to determine when events can be safely processed.



# The execution model prevents remote processes from blocking local ones, and does not require backtracking.



### However, this program is not schedulable!



### Remote events also trigger real-time violations.

Schedulability analysis tells us the program is flawed.



The program can be fixed with actors that increment the time stamps (model-time delays).



### This relaxes scheduling constraints...



... and we can derive conditions for schedulability...



... and being explicit about time delays means that we can analyze control system dynamics...



See "A Programming Model for Time-Synchronized Distributed Real-Time Systems", Yang Zhao, Jie Liu, and Edward A. Lee, RTAS '07, to be presented Friday.



## Is Truly Real-Time Computing Becoming Unachievable?

Yes! But the problem is solvable:

Actor-oriented component architectures implemented in *coordination languages* that complement rather than replace existing languages (e.g. PTIDES).

and

PRET machines that deliver repeatable timing with efficient pipelining, memory hierarchy, and networking

See the Ptolemy Project for ongoing research addressing these problems: http://ptolemy.org