Flicker: An Execution Infrastructure for TCB Minimization

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Trusted Computing Base (TCB)



Flicker's Properties

- Isolate security-sensitive code execution from all other code and devices
- Attest to security-sensitive code and its arguments and nothing else
- Convince a remote party that securitysensitive code was protected
- Add < 250 LoC to the software TCB



Outline

- Introduction
- Background
 - Trusted Platform Module (TPM)
 - Late Launch
- Flicker Architecture and Extensions
- Flicker Applications
- Performance Evaluation
- Related Work and Conclusions

TPM Background

- The Trusted Platform Module (TPM) is a dedicated security chip
- Can provide an *attestation* to remote parties
 - Platform Configuration Registers (PCRs) summarize the computer's software state
 - PCR_Extend(N, V): $PCR_N = SHA-1(PCR_N | V)$
 - TPM provides a signature over PCR values
 - A subset of dynamic PCRs can be reset without a reboot



Late Launch Background

- Supported by new commodity CPUs
 - SVM for AMD
 - TXT (formerly LaGrande) for Intel
- Designed to launch a VMM without a reboot
 - Hardware-based protections ensure launch integrity
- New CPU instruction (SKINIT/SENTER) accepts a memory region as input and atomically:
 - Resets dynamic PCRs
 - Disables interrupts
 - Extends a measurement of the region into PCR 17
 - Begins executing at the start of the memory region





Architecture Overview

- Core technique
 - Pause current execution environment (untrusted OS)
 - Execute security-sensitive code with hardwareenforced isolation
 - Resume previous execution
- Extensions
 - Attest only to code execution and protection
 - Preserve state securely across invocations
 - Establish secure communication with remote parties

Execution Flow



Attestation





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Attestation



Context Switch with Sealed Storage

- Seal data under combination of code, inputs, outputs
- Data unavailable to other code



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Developing With Flicker

- Sensitive code linked against the Flicker library
- Customized linker script lays out binary
- Application interacts with Flicker via a Flicker kernel module



Default Functionality

- Shim can execute arbitrary x86 code but provides very limited functionality
- Fortunately, many security-sensitive functions do not require much

- E.g., key generation, encryption/decryption, FFT

- Functionality can be added to support a particular security-sensitive operation
- We have partially automated the extraction of support code for security-sensitive code

Existing Flicker Modules

- OS Protection
- Crypto
- Memory Alloc. ٠
- Secure Channel
- TPM Driver
- TPM Utilities

Crypto ops (RSA, SHA-1, etc.)

Memory protection, ring 3 execution

Malloc/free/realloc

- Secure remote communication
- Communicate with TPM
- Perform TPM ops

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Application: Rootkit Detector

- Administrator can check the integrity of remote hosts
 - E.g., only allow uncompromised laptops to connect to the corporate VPN



Application: SSH Passwords



Other Applications Implemented

- Enhanced Certificate Authority (CA)
 - Private signing key isolated from entire system
- Verifiable distributed computing
 - Verifiably perform a computational task on a remote computer
 - Ex: SETI@Home, Folding@Home, distcc





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Generic Context-Switch Overhead

- Each Flicker context switch requires:
 SKINIT
 - TPM-based protection of application state

Results



Rootkit Detection Performance



Running detector every 30 seconds has negligible impact on system throughput

SSH Performance

- Setup time (217 ms) dominated by key generation (185 ms)
- Password verification (937 ms) dominated by TPM Unseal (905 ms)

Adds < 2 seconds of delay to client login

Optimizing Flicker's Performance

- Non-volatile storage
 - Access control based on PCRs
 - Read in 20ms, Write in 200 ms
 - Store a symmetric key for "sealing" and "unsealing" state

Reduces context-switch overhead by an order of magnitude

Hardware Performance Improvements

[ASPLOS 2008]

- Late launch cost only incurred when Flicker session launches
- TPM (Un)Seal only used for long-term storage
- Multicore systems remain interactive
- Context switch overheads (common case) resemble VM switches today (~0.5 µs)

Ongoing Work

- Creating a trusted path to the user
- Porting implementation to Intel
- Improving automatic privilege separation

Related Work

- Secure coprocessors
 - Dyad [Yee 1994], IBM 4758 [JiSmiMi 2001]
- System-wide attestation
 - Secure Boot [ArFaSm 1997], IMA [SaZhJaDo 2004], Enforcer [MaSmWiStBa 2004]
- VMM-based isolation
 - BIND [ShPeDo2005], AppCores [SiPuHaHe 2006], Proxos [TaLiLi 2006]
- "Traditional" uses of late launch
 - Trustworthy Kiosks [GaCáBeSaDoZh 2006], OSLO [Kauer 2007],

Conclusions

- Flicker greatly reduces an application's TCB
- Isolate security-sensitive code execution
- Provide fine-grained attestations
- Allow application writers to focus on the security of their own code